

61A Lecture 25

Announcements

Programming Languages

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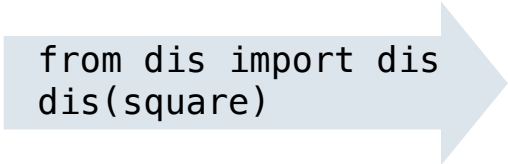
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from dis import dis  
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To create a new programming language, you either need a:

- **Specification:** A document describe the precise syntax and semantics of the language
- **Canonical Implementation:** An interpreter or compiler for the language

Parsing

Reading Scheme Lists

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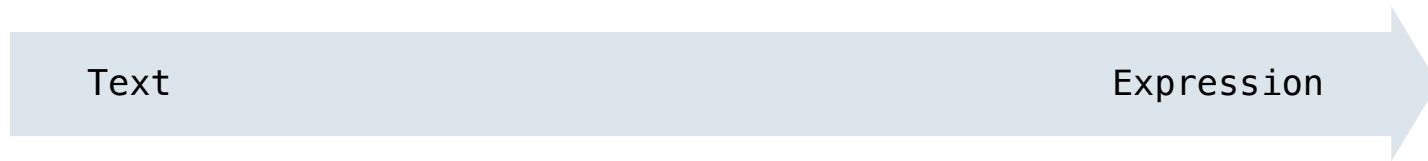
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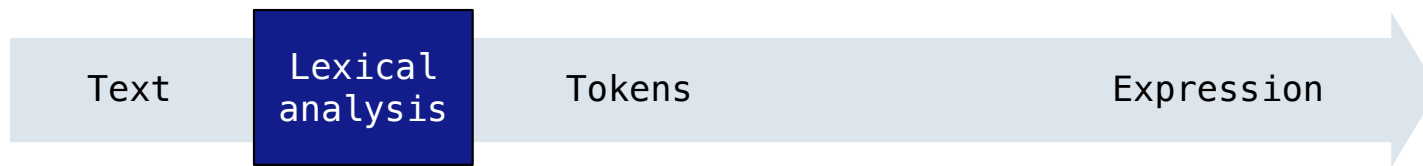
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| | | |
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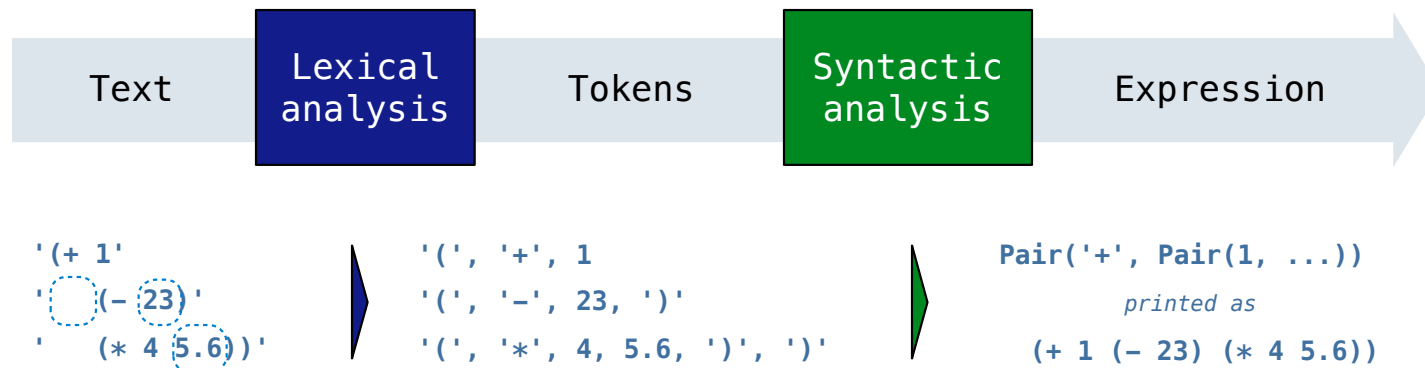
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(Demo)

Calculator

(Demo)

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...
TypeError: length attempted on improper list
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```

```
>>> s = Pair(1, Pair(2, Pair(3, nil)))
>>> print(s)
(1 2 3)
>>> len(s)
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>>> print(Pair(1, 2))
(1 . 2)
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Traceback (most recent call last):
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TypeError: length attempted on improper list
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Scheme expressions are represented as Scheme lists! Source code is data

The Pair Class

The Pair class represents Scheme pairs and lists. A list is a pair whose second element is either a list or nil.

```
class Pair:
    """A Pair has two instance attributes:
    first and second.

    For a Pair to be a well-formed list,
    second is either a well-formed list or nil.
    Some methods only apply to well-formed lists.
    """
    def __init__(self, first, second):
        self.first = first
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>>> s = Pair(1, Pair(2, Pair(3, nil)))
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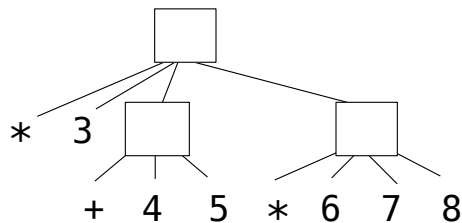
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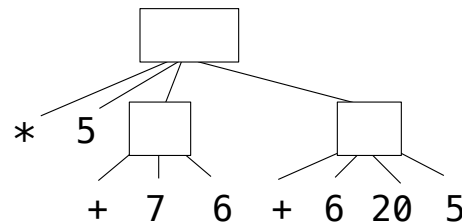
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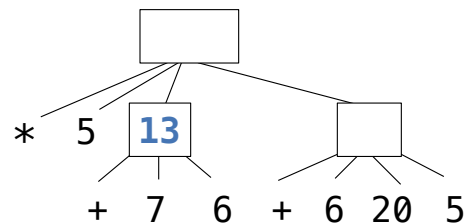
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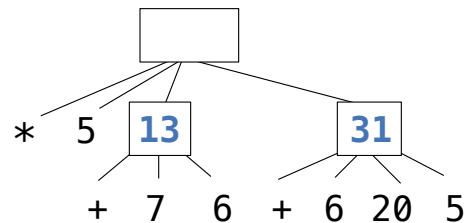
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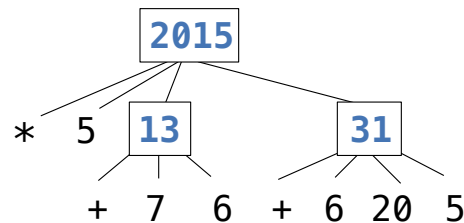
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def calc_apply(operator, args):
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...
...
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(Demo)

Interactive Interpreters

Read-Eval-Print Loop

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