CS61B Lecture #9: Interfaces and Abstract Classes

- HW #3 should be posted.
- Project #1 will be posted by the end of Thursday. Due 9 October.
- First test in class 17 October.
- Readings for today: §5.5 of A Java Reference.
Abstract Methods and Classes

• Instance method can be abstract: No body given; must be supplied in subtypes.

• One good use is in specifying a pure interface to a family of types:

```java
/** A drawable object. */
public abstract class Drawable { // "abstract" = "can’t say new Drawable"
    /** Expand THIS by a factor of SIZE */
    public abstract void scale (double size);
    /** Draw THIS on the standard output. */
    public abstract void draw ();
}
```

Now a Drawable is something that has at least the operations scale and draw on it. Can’t create a Drawable because it’s abstract—in particular, it has two methods without any implementation.

• BUT, we can write methods that operate on Drawables:

```java
void drawAll (Drawable[] thingsToDraw) {
    for (Drawable thing : thingsToDraw)
        thing.draw ();
}
```

• But draw has no implementation! How can this work?
Concrete Subclasses

- Can define kinds of Drawables that are non-abstract. To do so, must supply implementations for all methods:

  public class Rectangle extends Drawable {
      public Rectangle (double w, double h) { this.w = w; this.h = h; }
      public void scale (double size) { w *= size; h *= size; }
      public void draw () { draw a w × h rectangle }
      private double w,h;
  }

  Any Circle or Rectangle is a Drawable.

  public class Circle extends Drawable {
      public Circle (double rad) { this.rad = rad; }
      public void scale (double size) { rad *= size; }
      public void draw () { draw a circle with radius rad }
      double rad;
  }

- So, writing

  Drawable[] things = { new Rectangle (3, 4), new Circle (2) };
  drawAll (things);

  draws a 3 × 4 rectangle and a circle with radius 2.
Interfaces

- In generic use, an interface is a “point where interaction occurs between two systems, processes, subjects, etc.” (Concise Oxford Dictionary).

- In programming, often use the term to mean a description of this generic interaction, specifically, a description of the functions or variables by which two things interact.

- Java uses the term to refer to a slight variant of an abstract class that contains only abstract methods (and static constants).

- Idea is to treat Java interfaces as the public specifications of data types, and classes as their implementations:

  public interface Drawable {
    void scale (double size); // Automatically public abstract.
    void draw ();
  }

  public class Rectangle implements Drawable {
     ...
  }

- Interfaces are automatically abstract: can't say new Drawable();
can say new Rectangle(...).

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Multiple Inheritance

• Can extend one class, but implement any number of interfaces.

• Contrived Example:

```java
interface Readable {
    Object get ();
}

interface Writable {
    void put (Object x);
}

class Source implements Readable {
    public Object get () { ... }
}

class Sink implements Writable {
    public void put (Object x) { ... }
}

class Variable implements Readable, Writable {
    public Object get () { ... }
    public void put (Object x) { ... }
}
```

• The first argument of copy can be a Source or a Variable. The second can be a Sink or a Variable.
Review: Higher-Order Functions

• In Scheme, you had *higher-order functions* like this (adapted from *SICP*)

```scheme
(define (map proc items)
    (if (null? items)
        nil
        (cons (proc (car items)) (map proc (cdr items))))
)
```

and could write

```scheme
(map abs (list -10 2 -11 17))
```

```scheme
(10 2 11 17)
```

```scheme
(map (lambda (x) (* x x)) (list 1 2 3 4))
```

```scheme
(1 4 9 16)
```

• Java does not have these directly, but can use abstract classes or interfaces and subtyping to get the same effect (with more writing)
Map in Java

/** Function with one integer argument */

public interface IntUnaryFunction {
    int apply (int x);
}

IntList map (IntUnaryFunction proc, IntList items) {
    if (items == null)
        return null;
    else return new IntList (proc.apply (items.head),
                              map (proc, items.tail));
}

• It's the use of this function that's clumsy. First, define class for absolute value function; then create an instance:

class Abs implements IntUnaryFunction {
    public int apply (int x) { return Math.abs (x); }
}

map (new Abs (), some list);

• Or, we can write a lambda expression (sort of):

map (new IntUnaryFunction () {
    public int apply (int x) { return x*x; }
}, some list);

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Review: A Puzzle

class A {
    void f () { System.out.println ("A.f"); }
    void g () { f (); /* or this.f() */ }
    //static void g (A y) { y.f(); }
}

class B extends A {
    void f () { System.out.println ("B.f"); }
    // } } 
}

class C {
    static void main (String[] args) {
        B aB = new B ();
        h (aB);
    }
    static void h (A x) { x.g(); }
    //static void h (A x) { A.g(x); } // x.g(x) also legal here
}

1. What is printed?  Choices:

2. What if we made g static? a. A.f
3. What if we made f static? b. B.f
4. What if f were not defined in A? c. Some kind of error
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1. What is printed?

2. What if we made \texttt{g} static?

3. What if we made \texttt{f} static?

4. What if \texttt{f} were not defined in \texttt{A}?

\textbf{Choices:}

a. \texttt{A.f} \\
b. \texttt{B.f} \\
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    void f () { System.out.println ("A.f"); }
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    static void g (A y) { y.f(); }
}

class B extends A {
    void f () { System.out.println ("B.f"); }
    // ... }
}
class C {
    static void main (String[] args) {
        B aB = new B ();
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3. What if we made \texttt{f} static?
4. What if \texttt{f} were not defined in \texttt{A}?

\begin{itemize}
    \item \textbf{Choices:}
    \begin{enumerate}
        \item a. A.f
        \item b. B.f
        \item c. Some kind of error
    \end{enumerate}
\end{itemize}
class A {
    void g () { f (); /* or this.f() */ }
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    void f () {
        System.out.println ("B.f");
    }
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class C {
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Choices:
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Answer to Puzzle

1. Executing `java C` prints ___, because

   1. `C.main` calls `h` and passes it `aB`, whose dynamic type is `B`.
   2. `h` calls `x.g()`. Since `g` is inherited by `B`, we execute the code for `g` in class `A`.
   3. `g` calls `this.f()`. Now `this` contains the value of `h`’s argument, whose dynamic type is `B`. Therefore, we execute the definition of `f` that is in `B`.
   4. In calls to `f`, in other words, static type is ignored in figuring out what method to call.

2. If `g` were static, we see ___; selection of `f` still depends on dynamic type of `this`.

3. If `f` were static, would print ___ because then selection of `f` would depend on static type of `this`, which is `A`.

4. If `f` were not defined in `A`, we’d get ________________.
Answer to Puzzle

1. Executing `java C` prints `B.f`, because
   1. `C.main` calls `h` and passes it `aB`, whose dynamic type is `B`.
   2. `h` calls `x.g()`. Since `g` is inherited by `B`, we execute the code for `g` in class `A`.
   3. `g` calls `this.f()`. Now `this` contains the value of `h`'s argument, whose dynamic type is `B`. Therefore, we execute the definition of `f` that is in `B`.
   4. In calls to `f`, in other words, static type is ignored in figuring out what method to call.

2. If `g` were static, we see `B.f`; selection of `f` still depends on dynamic type of `this`.

3. If `f` were static, would print `A.f` because then selection of `f` would depend on static type of `this`, which is `A`.

4. If `f` were not defined in `A`, we'd get a compile-time error.
Example: Designing a Class

Problem: Want a class that represents histograms, like this one:

![Histogram Example](image)

Analysis: What do we need from it? At least:

- Specify buckets and limits.
- Accumulate counts of values.
- Retrieve counts of values.
- Retrieve numbers of buckets and other initial parameters.
**Specification Seen by Clients**

- The *clients* of a module (class, program, etc.) are the programs or methods that *use* that module’s exported definitions.

- In Java, intention is that exported definitions are designated `public`.

- Clients are intended to rely on *specifications*, not code.

- **Syntactic specification**: method and constructor headers—syntax needed to use.

- **Semantic specification**: what they do. No formal notation, so use comments.
  - Semantic specification is a *contract*.
  - Conditions client must satisfy (*preconditions*, marked “Pre:” in examples below).
  - Promised results (*postconditions*).
  - Design these to be all the *client needs*!
  - Exceptions communicate errors, specifically failure to meet pre-conditions.
/** A histogram of floating-point values */
public interface Histogram {
    /** The number of buckets in THIS. */
    int size();

    /** Lower bound of bucket #K. Pre: 0<=K<size(). */
    double low(int k);

    /** # of values in bucket #K. Pre: 0<=K<size(). */
    int count(int k);

    /** Add VAL to the histogram. */
    void add(double val);
}

void fillHistogram(Histogram H, Scanner in) {
    while (in.hasNextDouble())
        H.add(in.nextDouble());
}

void printHistogram(Histogram H) {
    for (int i = 0; i < H.size(); i += 1)
        System.out.printf(">=%5.2f | %4d%n", H.low(i), H.count(i));
}
public class FixedHistogram implements Histogram {
    private double low, high; /* From constructor*/
    private int[] count; /* Value counts */

    /** A new histogram with SIZE buckets recording values >= LOW and < HIGH. */
    public FixedHistogram (int size, double low, double high)
    {
        if (low >= high || size <= 0) throw new IllegalArgumentException ();
        this.low = low; this.high = high;
        this.count = new int[size];
    }

    public int size () { return count.length; }
    public double low (int k) { return low + k * (high-low)/count.length; }

    public int count (int k) { return count[k]; }

    public void add (double val) {
        int k = (int) ((val-low)/(high-low) * count.length);
        if (k >= 0 && k < count.length) count[k] += 1;
    }
}
Let's Make a Tiny Change

Don't require *a priori* bounds:

class FlexHistogram implements Histogram {
  /** A new histogram with SIZE buckets. */
  public FlexHistogram (int size) {
      ?
  }
  // What needs to change?
}

- How would you do this? Profoundly changes implementation.
- But *clients* (like printHistogram and fillHistogram) *still work with no changes.*
- Illustrates the power of *separation of concerns.*
Implementing the Tiny Change

• Pointless to pre-allocate the count array.
• Don’t know bounds, so must save arguments to add.
• Then recompute count array “lazily” when count(···) called.
• Invalidate count array whenever histogram changes.

```java
class FlexHistogram implements Histogram {
    private List<Double> values = ...; // Java library type (later)
    int size;
    private int[] count;

    public FlexHistogram (int size) { this.size = size; this.count = null; }

    public void add (double x) { count = null; values.add (x); }

    public int count (int k) {
        if (count == null) { compute count from values here. }
        return count[k];
    }
}
```
Advantages of Procedural Interface over Visible Fields

By using public method for \texttt{count} instead of making the array \texttt{count} visible, the “tiny change” is transparent to clients:

- If client had to write \texttt{myHist.count[k]}, would mean
  
  “The number of items currently in the \(k^{\text{th}}\) bucket of histogram \texttt{myHist} (and by the way, there is an array called \texttt{count} in \texttt{myHist} that always holds the up-to-date count).”

- Parenthetical comment \textit{useless} to the client.

- But if \texttt{count} array had been visible, after “tiny change,” every \texttt{use} of \texttt{count} in client program would have to change.

- So using a method for the public \texttt{count} decreases what client \texttt{has to know}, and (therefore) has to change.