Today: Backtracking searches, game trees (DSIJ, Section 6.5)

Coming Up: Concurrency and synchronization (Data Structures, Chapter 10, and Assorted Materials On Java, Chapter 6; Graph Structures: DSIJ, Chapter 12.)
Searching by “Generate and Test”

• We’ve been considering the problem of searching a set of data stored in some kind of data structure: “Is \( x \in S \)?”

• But suppose we don’t have a set \( S \), but know how to recognize what we’re after if we find it: “Is there an \( x \) such that \( P(x) \)?”

• If we know how to enumerate all possible candidates, can use approach of Generate and Test: test all possibilities in turn.

• Can sometimes be more clever: avoid trying things that won’t work, for example.

• What happens if the set of possible candidates is infinite?
Backtracking Search

- Backtracking search is one way to enumerate all possibilities.
- Example: *Knight’s Tour*. Find all paths a knight can travel on a chess-board such that it touches every square exactly once and ends up one knight move from where it started.
- In the example below, the numbers indicate position numbers (knight starts at 0).
- Here, knight (N) is stuck; how to handle this?

```
  6  5
  4  7
10  2
 8  3  0
 N 9  1
```
**Append to PATH a sequence of knight moves starting at ROW, COL**
* that avoids all squares that have been hit already and
* that ends up one square away from ENDROW, ENDCOL. B[i][j] is
* true iff row i and column j have been hit on PATH so far.
* Returns true if it succeeds, else false (with no change to PATH).
* Call initially with PATH containing the starting square, and
* the starting square (only) marked in B. */

```java
boolean findPath (boolean[][] b, int row, int col,
                 int endRow, int endCol, List path) {
    if (path.size () == 64) return isKnightMove (row, col, endRow, endCol);
    for (r, c = all possible moves from (row, col)) {
        if (! b[r][c]) {
            b[r][c] = true; // Mark the square
            path.add (new Move (r, c));
            if (findPath (b, r, c, endRow, endCol, path)) return true;
            b[r][c] = false; // Backtrack out of the move.
            path.remove (path.size ()-1);
        }
    }
    return false;
}
```
Another Kind of Search: Best Move

- Consider the problem of finding the best move in a two-person game.
- One way: assign a value to each possible move and pick highest.
  - Example: number of our pieces - number of opponent’s pieces.
- But this is misleading. A move might give us more pieces, but set up a devastating response from the opponent.
- So, for each move, look at opponent’s possible moves, assume he picks the best one for him, and use that as the value.
- But what if you have a great response to his response?
- How do we organize this sensibly?
Game Trees, Minimax

• Think of the space of possible continuations of the game as a tree.
• Each node is a position, each edge a move.

Numbers are the values we guess for the positions (larger means better for me). Starred nodes would be chosen.

I always choose child (next position) with maximum value; opponent chooses minimum value (“Minimax algorithm”)
Alpha-Beta Pruning

• We can prune this tree as we search it.

```
-5
  /   |
-5   ≥5
  /   / |
-30 -5 5
```

```
-20
  /   |
-20 -30
  /   |
```

• At the '≥ 5' position, I know that the opponent will not choose to move here (since he already has a -5 move).

• At the '≤ -20' position, my opponent knows that I will never choose to move here (since I already have a -5 move).
Cutting off the Search

• If you could traverse game tree to the bottom, you’d be able to force a win (if it’s possible).

• Sometimes possible near the end of a game.

• Unfortunately, game trees tend to be either infinite or impossibly large.

• So, we choose a maximum depth, and use a heuristic value computed on the position alone (called a static valuation) as the value at that depth.

• Or we might use iterative deepening (kind of breadth-first search), and repeat the search at increasing depths until time is up.

• Much more sophisticated searches are possible, however (take CS188).
Some Pseudocode for Searching

/** A legal move for WHO that either has an estimated value >= CUTOFF
 * or that has the best estimated value for player WHO, starting from
 * position START, and looking up to DEPTH moves ahead. */
Move findBestMove (Player who, Position start, int depth, double cutoff)
{
    if (start is a won position for who) return WON_GAME; /* Value ∞ */
    else if (start is a lost position for who) return LOST_GAME; /* Value −∞ */
    else if (depth == 0) return guessBestMove (who, start, cutoff);

    Move bestSoFar = REALLY_BAD_MOVE;
    for (each legal move, M, for who from position start) {
        Position next = start.makeMove (M);
        Move response = findBestMove (who.opponent (), next, depth-1, -bestSoFar.value ());
        if (-response.value () > bestSoFar.value ()) {
            Set M's value to -response.value (); // Value for who = - Value for opponent
            bestSoFar = M;
            if (M.value () >= cutoff) break;
        }
    }
    return bestSoFar;
}
Static Evaluation

- This leaves static evaluation, which looks just at the next possible move:

  ```
  Move guessBestMove (Player who, Position start, double cutoff) {
      Move bestSoFar;
      bestSoFar = Move.REALLY_BAD_MOVE;
      for (each legal move, M, for who from position start) {
          Position next = start.makeMove (M);
          Set M’s value to heuristic guess of value to who of next;
          if (M.value () > bestSoFar.value ()) {
              bestSoFar = M;
              if (M.value () >= cutoff)
                  break;
          }
      }
      return bestSoFar;
  }
  ```