# EECS 151/251A Homework 1 

Due Monday, Feb 4 ${ }^{\text {th }}$, 2019

## Problem 1: Moore's Law [20 pts]

Consider state-of-the-art processor chips from the 1970's, 1980's, 1990's, 2000's, and after 2010. Choose a processor from each period. (You may choose which every processor you like, but make sure they are spaced out by around 10 years. Use of Wikipedia or WikiChip is acceptable.)

1. For each look up the approximate number of transistor per chip. Plot the number of transistors per chip over time, with a log scale on the $y$-axis [ 5 pts ].
2. For your processor choice after 2010, approximate the ratio of transistors used for on-chip memory (caches) vs logic circuits. Assume 6 transistors per memory bit [5 pts].
3. List a major product innovation enabled by each of these processors [5 pts].
4. On another set of axis, plot the clock frequency of each over time (on a log plot) [5 pts].

## Solution:

| Chip | Year | Transistors | Clock <br> Freq. <br> Low <br> (MHz) | Clock <br> Freq. <br> High <br> (MHz) | Key Advance | Src |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & \text { Intel } \\ & 8080 \end{aligned}$ | 1974 | 6K | 2 | 3.125 | One of the first widespread microprocessors owing in part to its 40-pin package, faster transistors (NMOS), compatibility with TTL, and and increased memory bus size (with ability to access 64 KB of memory) | https://en.wikipe org/wiki/Intel_80 |
| Intel <br> 80386 | 1985 | 275K | 12 | 40 | First Implementation of 32-bit extension to 80286 . A common ancestor of modern x86 processors and is also known by the name i386. | https://en.wikiped org/wiki/Intel_803 |
| Intel <br> Pen- <br> tium <br> P55C | 1997 | 4.5M | 120 | 233 | First Implementation of MMX (SIMD) extension in Intel Processors. <br> Predecessor to SSE and AVX vector extensions | https://en.wikiped <br> org/wiki/P5_ <br> (microarchitecture) <br> \#MMX |
| Intel <br> Prescott | 2004 | 125M | 2400 | 3067 | First of the Pentium line to introduce SSE3. Major overhaul that hit the heat barrier | https://en.wikiped org/wiki/Pentium_4 https://en.wikiped org/wiki/List_of_ Intel_Pentium_4_ microprocessors\# Prescott_(90_nm) |
| Intel <br> Ivy <br> Bridge- <br> E | 2013 | 1.86B | 3400 | 3600 | Die shrinking of Sandy Bridge-E to 22 nm (finfet) | https://en.wikipedia org/wiki/Ivy_Bridg (microarchitecture) https://en.wikiped org/wiki/List_of _ Intel_Core_i7_ microprocessors\# \%22Ivy_BridgeE\%22_(22_nm) |


| Intel Ivy Bridge-E Details: |  |
| :--- | :--- |
| Num Transistors | 1860000000 |
| Num Cores | 6 |
| L1 Cache Per Core (Bytes) | 32768 |
| L2 Cache Per Core | 262144 |
| L3 Cache Shared (Bytes) | 15728640 |
| Cahce Total (Bytes) | 17498112 |
| Cache Total (Bits) | 139984896 |
| Transistors Per Bit (Est) | 6 |
| Cache Transistors (Est) | 839909376 |
| Logic Transistors (Est) | 1020090624 |
| Ratio Memory/Logic | 0.823367411 |



Figure 1: Transistors Trend


## Problem 2: Design Tradeoffs [5 pts]

Sketch a plot showing what you expect the design space of a processor looks like in terms of power (watts) and performance (frequency). Show several possible design points as dots in the design space including some on the Pereto Optimal frontier. Also, draw the Pareto Optimal frontier.

## Solution:



Want this to be low

Figure 3: Performance, Power Design Space

## Problem 3: 251A only - Optional Challenge Question for 151 [5 pts]

Dennard's scaling says that if all transistor dimensions and the operating voltage is scaled linearly, the the resulting transitors will speed up linearly and that power density (power per unit area) will remain constant. However, as Moore's Law progressed and circuits were scaled down, processor chip performance and power both increased at a rate faster than linear. How can you explain why frequencies scaled at a rate higher than linear with process scaling. What was the effect of this frequency scaling on power consumption?

## Solution:

1. Voltages didn't scale down
2. From a power perspective, the die actually got bigger, so the total power consumption went up
3. Architectural changes, e.g. critical path reduction, deep pipelines, etc, let the frequency go up (which then explains why the power went up)

## Problem 4: Full Custom versus Standard Cell [6 pts]

Consider custom chip design with the two methods described in class: Standard Cells and Full Custom.

Breifly describe for each the main advantage and main disadvantage relative to the other.

## Solution:

Standard Cell Approach:
Main advantage: Layout Fully automated and cells pre-verified and leading to higher designer productivity and reduced NRE costs Main disadvantage: Circuits Less optimized for performance, cost, power as may be needed by the application. No analog processing possible.

Full Custom Approach:
Main advantage: Circuits and layout can be optimized for performance, cost, power as needed by particular application.

Main disadvantage: Long design process, for sixing transistors, generating layout, verifying layout, etc.

## Problem 5: Implementation Alternatives [12 pts]

Consider the following implementation approaches for building some digital system: Full Custom, Standard Cell, Gate Array, FPGA, micro-processor.

For the best implementation possible, rank order (a partial ordering might be okay) from highest to lowest a) for NRE costs [ 3 pts ], b) per part costs [ 3 pts ], c) performance [3 pts], and d) flexibility [3 pts].

## Solution:

a) NRE costs Full Custom, Standard Cell, Gate Array, FPGA, micro-processor.
b) Per part costs Micro-processor, FPGA, Gate Array, Standard Cell, full-custom. Also accept (in the worst case): Full-custom, micro-processor, FPGA, Gate Array, Standard Cell.
c) Performance Full Custom, Standard Cell, Gate Array, FPGA, micro-processor. Also accept: Full Custom, Standard Cell, Gate Array, micro-processor, FPGA.
d) Flexibilty for reuse in other applications Micro-processor, FPGA, Gate Array/Standard-Cell/Full-Custom. Also accept: FPGA, Micro-processor, Gate Array/Standard-Cell/Full-Custom.

Alternative orderings with sensible explanations may also be accepted.

## Problem 6: Combinational Logic [4 pts]

Consider a combinational logic block that multiplies two N-bit numbers yeilding a 2 N -bit result. How many rows are in the truth table that represents the function of this block? How many rows are there when $\mathrm{N}=8$ ?

## Solution:

A truth table has a 1 row for every combination of input values. Since there are two $N$-bit inputs, the total input width is $2 N$ bits, and each bit has 2 possible values ( 0 or 1 ). This means that there will be $2^{2 N}$ possible input states and $2^{2 N}$ truth table rows. When there are two $N=8$ bit inputs, the table has $2^{(2 \times 8)}=65536$ entries.

## Problem 7: Combinational Logic Functions [4 pts]

Consider the set of CL blocks with N inputs and 1 output. How many unique such blocks exist?

## Solution:

$2^{2^{N}}$. Why? There are $2^{N}$ input combinations to a CL block with $N$ inputs an 1 output. For each of these input values, the CL produces a 1-bit output, which has 2 possible states. So there are $2^{2^{N}}$ different mappings from inputs to output state that the CL logic can encode. Blocks implementing different mappings are unique.

## Problem 8: Combinational Logic Gates [6 pts]

In class we said that any combinational logic function can be implemented with either AND-gates along with inverters or with OR-gates along with inverters. There is a hypothesis that NAND-gates alone or NOR-gates alone are sufficient to implement any combinational logic function. Prove or disprove this hypothesis.

## Solution:

You can turn a NAND gate into an inverter (connect inputs together, or fix one input high). With this inverter, you can turn a NAND gate into an AND gate. Then you have AND-gates and inverters, which we know to be sufficient to implement any combinational logic function. Similarly for NOR gates.

