

Continuous Probability 3

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Independent Continuous Random Variables Definition: The continuous RVs X and Y are independent if

 $Pr[X \in A, Y \in B] = Pr[X \in A]Pr[Y \in B], \forall A, B.$

Theorem: The continuous RVs X and Y are independent if and only if

$$f_{X,Y}(x,y) = f_X(x)f_Y(y)$$

Proof: As in the discrete case. **Definition:** The continuous RVs X_1, \ldots, X_n are mutually independent if

 $Pr[X_1 \in A_1, \ldots, X_n \in A_n] = Pr[X_1 \in A_1] \cdots Pr[X_n \in A_n], \forall A_1, \ldots, A_n.$

Theorem: The continuous RVs $X_1, ..., X_n$ are mutually independent if and only if

$$f_{\mathbf{X}}(x_1,\ldots,x_n)=f_{X_1}(x_1)\cdots f_{X_n}(x_n).$$

Proof: As in the discrete case.

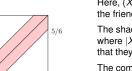
Review: CDF and PDF.

Key idea: For a continuous RV, Pr[X = x] = 0 for all $x \in \mathfrak{N}$. Examples: Uniform in [0, 1]; throw a dart in a target. Thus, one cannot define Pr[outcome], then Pr[event]. Instead, one starts by defining Pr[event]. Thus, one defines $Pr[X \in (-\infty, X]] = Pr[X \le x] =: F_X(x), x \in \mathfrak{N}$. Then, one defines $f_X(x) := \frac{d}{dx}F_X(x)$. Hence, $f_X(x)\varepsilon = Pr[X \in (x, x + \varepsilon)]$. $F_X(\cdot)$ is the cumulative distribution function (CDF) of X. $f_X(\cdot)$ is the probability density function (PDF) of X.

Meeting at a Restaurant

Two friends go to a restaurant independently uniformly at random between noon and 1pm.

They agree they will wait for 10 minutes. What is the probability they meet?



Here, (X, Y) are the times when the friends reach the restaurant. The shaded area are the pairs

where |X - Y| < 1/6, i.e., such that they meet.

The complement is the sum of two rectangles. When you put them together, they form a square with sides 5/6.

Thus, $Pr[\text{meet}] = 1 - (\frac{5}{6})^2 = \frac{11}{36}$.

 $0^{-1/6}$

Expectation

Definitions: (a) The **expectation** of a random variable *X* with pdf f(x) is defined as $E[X] = \int_{-\infty}^{\infty} x f_X(x) dx.$

(b) The expectation of a function of a random variable is defined as

$$E[h(X)] = \int_{-\infty}^{\infty} h(x) f_X(x) dx.$$

(c) The expectation of a function of multiple random variables is defined as

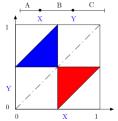
$$E[h(\mathbf{X})] = \int \cdots \int h(\mathbf{x}) f_{\mathbf{X}}(\mathbf{x}) dx_1 \cdots dx_n.$$

Justifications: Think of the discrete approximations of the continuous RVs.

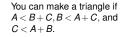
Breaking a Stick

You break a stick at two points chosen independently uniformly at random.

What is the probability you can make a triangle with the three pieces?

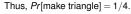


Let X, Y be the two break points along the [0, 1] stick.



If X < Y, this means X < 0.5, Y < X + 0.5, Y > 0.5. This is the blue triangle.

If X > Y, we get the red triangle, by symmetry.



Pr[make triangle] – 1 /4

Maximum of Two Exponentials

Let $X = Expo(\lambda)$ and $Y = Expo(\mu)$ be independent. Define $Z = \max\{X, Y\}$. Calculate E[Z]. We compute f_Z , then integrate. One has Pr[Z < z] = Pr[X < z, Y < z] = Pr[X < z]Pr[Y < z] $= (1 - e^{-\lambda z})(1 - e^{-\mu z}) = 1 - e^{-\lambda z} - e^{-\mu z} + e^{-(\lambda + \mu)z}$. Thus, $f_Z(z) = \lambda e^{-\lambda z} + \mu e^{-\mu z} - (\lambda + \mu)e^{-(\lambda + \mu)z}, \forall z > 0.$

 $E[Z] = \int_0^\infty z f_Z(z) dz = \frac{1}{\lambda} + \frac{1}{\mu} - \frac{1}{\lambda + \mu}.$

Quantization Noise

We saw that $E[Z^2] = \frac{1}{3}2^{-2(n+1)}$ and $E[X^2] = \frac{1}{3}$. The **signal to noise ratio** (SNR) is the power of the signal divided by the power of the noise.

Thus,

Hence,

 $SNR = 2^{2(n+1)}$.

Expressed in decibels, one has

 $SNR(dB) = 10\log_{10}(SNR) = 20(n+1)\log_{10}(2) \approx 6(n+1).$

For instance, if n = 16, then $SNR(dB) \approx 112dB$.

Maximum of *n* i.i.d. Exponentials

Let X_1, \ldots, X_n be i.i.d. *Expo*(1). Define $Z = \max\{X_1, X_2, \ldots, X_n\}$. Calculate E[Z].

We use a recursion. The key idea is as follows:

 $Z = \min\{X_1, \ldots, X_n\} + V$

where V is the maximum of n-1 i.i.d. Expo(1). This follows from the memoryless property of the exponential.

Let then $A_n = E[Z]$. We see that

$$A_n = E[\min\{X_1, ..., X_n\}] + A_{n-1}$$

= $\frac{1}{n} + A_{n-1}$

because the minimum of *Expo* is *Expo* with the sum of the rates. Hence,

 $E[Z] = A_n = 1 + \frac{1}{2} + \dots + \frac{1}{n} = H(n).$

Replacing Light Bulbs

Say that light bulbs have i.i.d. Expo(1) lifetimes. We turn a light on, and replace it as soon as it burns out. How many light bulbs do we need to replace in *t* units of time? **Theorem:** The number X_t of replaced light bulbs is P(t). That is, $Pr[X_t = n] = \frac{t^n}{n}e^{-t}$. **Proof:** We see how X_t increases over the next $\varepsilon \ll 1$ time units. Let *A* be the event that a burns out during $[t, t + \varepsilon]$. Then, $Pr[X_{t+\varepsilon} = n] \approx Pr[X_t = n, A^c] + Pr[X_t = n - 1, A]$ $= Pr[X_t = n]Pr[A^c] + Pr[X_t = n - 1]Pr[A]$

 $= Pr[X_t = n]Pr[A^c] + Pr[X_t = n-1]Pr[A]$ $\approx Pr[X_t = n](1-\varepsilon) + Pr[X_t = n-1]\varepsilon.$

Hence, $g(n,t) := Pr[X_t = n]$ is such that

 $g(n,t+\varepsilon) \approx g(n,t) - g(n,t)\varepsilon + g(n-1,t)\varepsilon.$

Quantization Noise

In digital video and audio, one represents a continuous value by a finite number of bits.

This introduces an error perceived as noise: the quantization noise. What is the power of that noise?

Model: X = U[0, 1] is the continuous value. *Y* is the closest multiple of 2^{-n} to *X*. Thus, we can represent *Y* with *n* bits. The error is Z := X - Y.

The power of the noise is $E[Z^2]$.

Analysis: We see that *Z* is uniform in $[0, a = 2^{-(n+1)}]$. Thus,

$$E[Z^2] = \frac{a^2}{3} = \frac{1}{3}2^{-2(n+1)}.$$

The power of the signal X is $E[X^2] = \frac{1}{3}$.

Replacing Light Bulbs

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Theorem: The number X_t of replaced light bulbs is P(t).

That is, $Pr[X_t = n] = \frac{t^n}{n!}e^{-t}$.

Proof: (continued) We saw that

 $g(n,t+\varepsilon) \approx g(n,t) - g(n,t)\varepsilon + g(n-1,t)\varepsilon.$

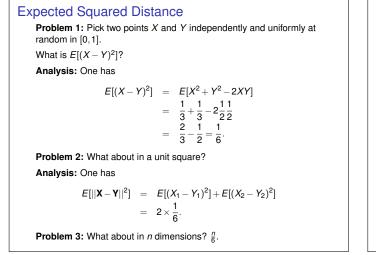
Subtracting g(n, t), dividing by ε , and letting $\varepsilon \rightarrow 0$, one gets

$$g'(n,t) = -g(n,t) + g(n-1,t).$$

You can check that these equations are solved by $g(n, t) = \frac{t^n}{n!} e^{-t}$. Indeed, then

$$g'(n,t) = \frac{t^{n-1}}{(n-1)!}e^{-t} - g(n,t)$$

= $g(n-1,t) - g(n,t).$



Geometric and Exponential

The geometric and exponential distributions are similar. They are both memoryless. Consider flipping a coin every 1/N second with Pr[H] = p/N, where $N \gg 1$. Let *X* be the time until the first *H*. Fact: $X \approx Expo(p)$. Analysis: Note that $Pr[X > t] \approx Pr[first Nt flips are tails]$ $= (1 - \frac{p}{N})^{Nt} \approx \exp\{-pt\}.$

Indeed, $(1-\frac{a}{N})^N \approx \exp\{-a\}.$

Summary

 Continuous Probability 3

 • Continuous RVs are essentially the same as discrete RVs

 • Think that
$$X \approx x$$
 with probability $f_X(x)\varepsilon$

 • Sums become integrals,

 • The exponential distribution is magical: memoryless.