



Polynomials. Secret Sharing.

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A number from 0 to 10.

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Share secret among *n* people.

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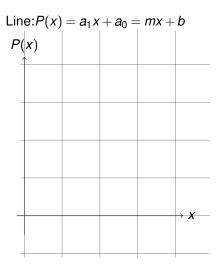
Polynomials P(x) with arithmetic modulo p: ¹ $a_i \in \{0, ..., p-1\}$ and

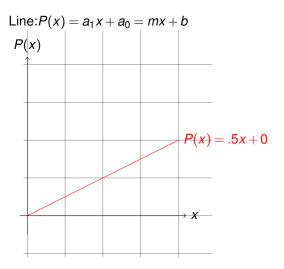
$$P(x) = a_d x^d + a_{d-1} x^{d-1} \cdots + a_0 \pmod{p},$$

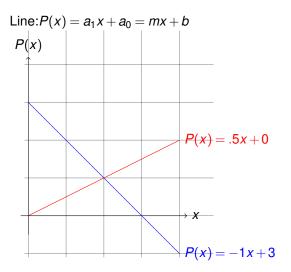
for $x \in \{0, \dots, p-1\}.$

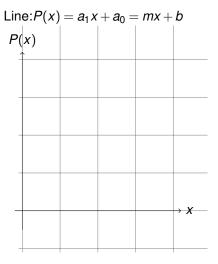
Line: $P(x) = a_1 x + a_0$

Line: $P(x) = a_1x + a_0 = mx + b$

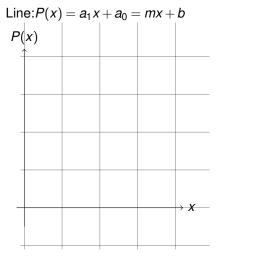




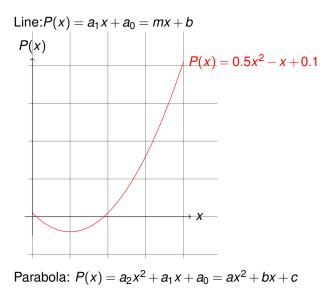


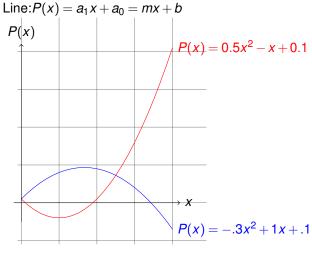


Parabola: $P(x) = a_2 x^2 + a_1 x + a_0$

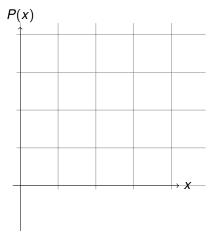


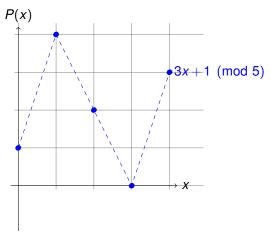
Parabola: $P(x) = a_2x^2 + a_1x + a_0 = ax^2 + bx + c$

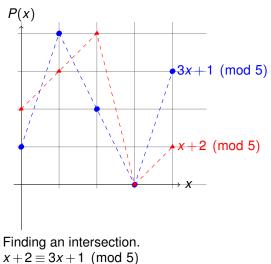




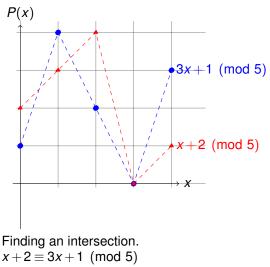
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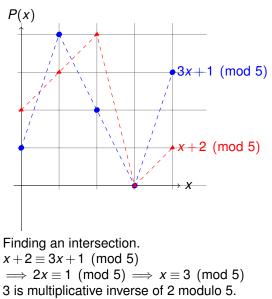


 \implies 2*x* \equiv 1 (mod 5)



 \implies 2x \equiv 1 (mod 5) \implies x \equiv 3 (mod 5) 2 is multiplicative inverse of 2 module 5

3 is multiplicative inverse of 2 modulo 5.



Good when modulus is prime!!

Two points make a line.

Fact: Exactly 1 degree $\leq d$ polynomial contains d + 1 points.²

²Points with different x values.

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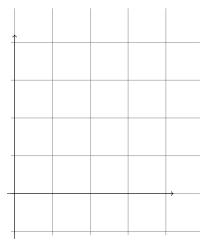
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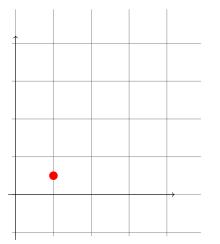
Fact: Exactly 1 degree $\leq d$ polynomial contains d + 1 points.²

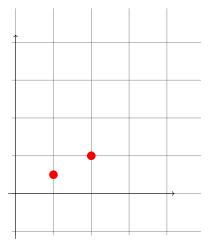
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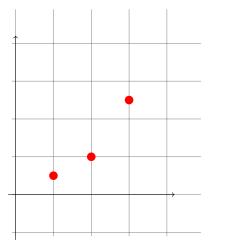
Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime *p* contains d + 1 pts.

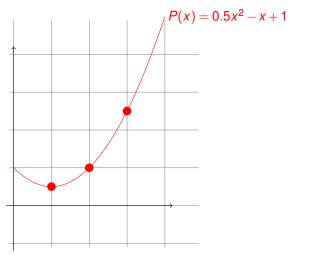
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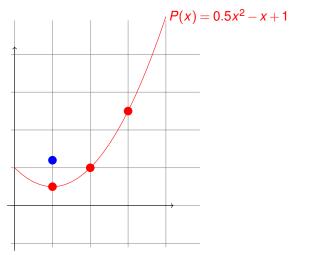


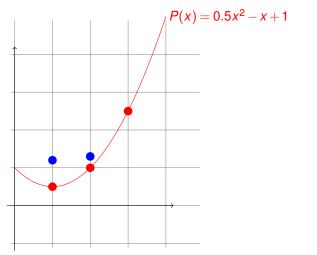


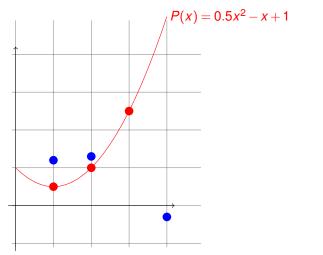


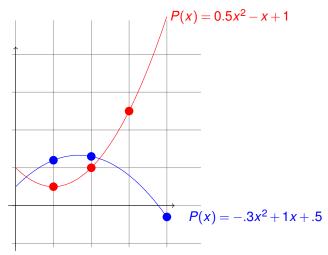




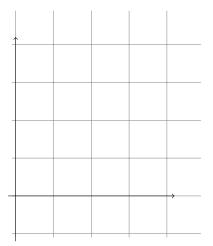


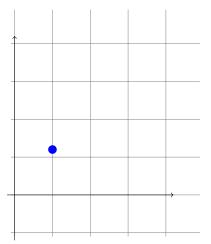


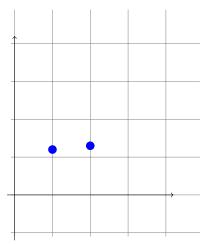


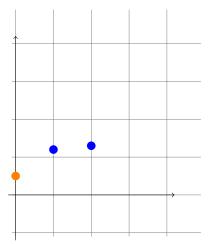


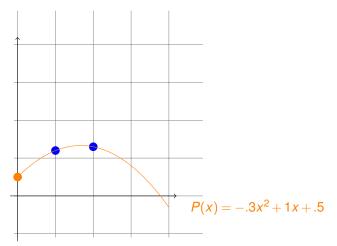
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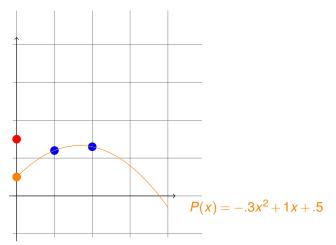


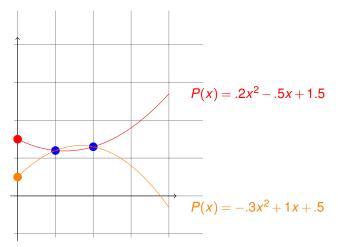


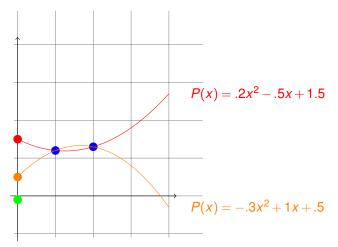


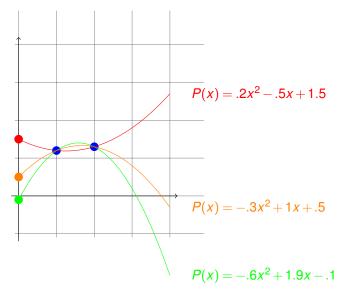


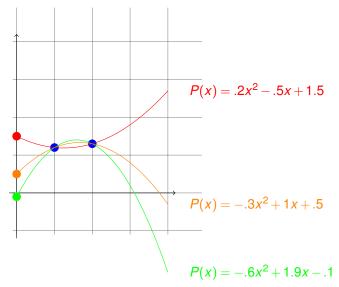












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Shamir's *k* out of *n* Scheme:

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Shares: points on a line.

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Shares: points on a line. Secret: *y*-intercept.

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Shares: points on a line. Secret: *y*-intercept. Arithmetic Modulo 11.

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Shares: points on a line. Secret: *y*-intercept. Arithmetic Modulo 11.

What's my secret?

For a line, $a_1x + a_0 = mx + b$ contains points (1,3) and (2,4).

P(1) =

$$P(1) = m(1) + b \equiv m + b$$

$$P(1) = m(1) + b \equiv m + b \equiv 3 \pmod{5}$$

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 $P(2) = m(2) + b \equiv 2m + b \equiv 4 \pmod{5}$

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Subtract first from second ..

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 $m \equiv 1 \pmod{5}$

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Backsolve: $b \equiv 2 \pmod{5}$.

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$$P(1) = m(1) + b \equiv m + b \equiv 3 \pmod{5}$$

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Subtract first from second..

$$m+b \equiv 3 \pmod{5}$$

 $m \equiv 1 \pmod{5}$

Backsolve: $b \equiv 2 \pmod{5}$. Secret is 2. And the line is...

 $x+2 \mod 5$.

$$P(1) = m(1) + b \equiv 5 \pmod{11}$$

 $P(3) = m(3) + b \equiv 9 \pmod{11}$

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Subtract first from second.

$$P(1) = m(1) + b \equiv 5 \pmod{11}$$

 $P(3) = m(3) + b \equiv 9 \pmod{11}$

Subtract first from second.

 $2m \equiv 4 \pmod{11}$

$$P(1) = m(1) + b \equiv 5 \pmod{11}$$

 $P(3) = m(3) + b \equiv 9 \pmod{11}$

Subtract first from second.

$$2m \equiv 4 \pmod{11}$$

Multiplicative inverse of 2 (mod 11) is 6:

$$P(1) = m(1) + b \equiv 5 \pmod{11}$$

 $P(3) = m(3) + b \equiv 9 \pmod{11}$

Subtract first from second.

 $2m \equiv 4 \pmod{11}$

Multiplicative inverse of 2 (mod 11) is 6: $6 \times 2 \equiv 12 \equiv 1 \pmod{11}$

$$P(1) = m(1) + b \equiv 5 \pmod{11}$$

 $P(3) = m(3) + b \equiv 9 \pmod{11}$

Subtract first from second.

$$2m \equiv 4 \pmod{11}$$

Multiplicative inverse of 2 (mod 11) is 6: $6 \times 2 \equiv 12 \equiv 1 \pmod{11}$ Multiply both sides by 6.

$$P(1) = m(1) + b \equiv 5 \pmod{11}$$

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Subtract first from second.

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Multiplicative inverse of 2 (mod 11) is 6: $6 \times 2 \equiv 12 \equiv 1 \pmod{11}$ Multiply both sides by 6.

 $12m = 24 \pmod{11}$

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Multiplicative inverse of 2 (mod 11) is 6: $6 \times 2 \equiv 12 \equiv 1 \pmod{11}$ Multiply both sides by 6.

> $12m = 24 \pmod{11}$ $m = 2 \pmod{11}$

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Backsolve: $2 + b \equiv 5 \pmod{11}$.

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What's my secret?

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So polynomial is $2x^2 + 1x + 4 \pmod{5}$

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Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime *p* contains *d* + 1 pts.

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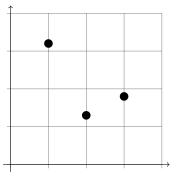
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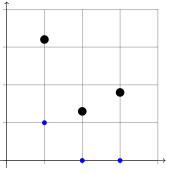
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Construction proves the existence of a degree *d* polynomial!

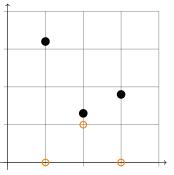


Points: (1,3.2), (2,1.3), (3,1.8).



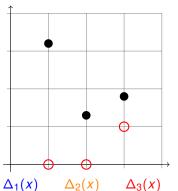
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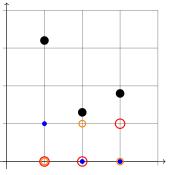


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 $\Delta_1(x) \qquad \Delta_2(x)$



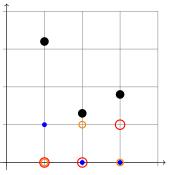
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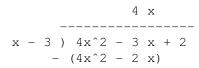
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Must prove Roots fact.

4 xx - 3) $4x^2 - 3x + 2$



$$\begin{array}{c} 4 \ x + 4 \\ x - 3 \) \ 4x^2 - 3 \ x + 2 \\ - \ (4x^2 - 2 \ x) \\ - - - - - \\ 4 \ x + 2 \end{array}$$

$$4 x + 4 r 4$$

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$$- (4x^{2} - 2 x)$$

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 $4x^2 - 3x + 2 \equiv (x - 3)(4x + 4) + 4 \pmod{5}$ In general, divide P(x) by (x - a) gives Q(x) and remainder r.

Polynomial Division. Divide $4x^2 - 3x + 2$ by (x - 3) modulo 5.

$$\begin{array}{c} 4 \ x + 4 \ r \ 4 \\ x - 3 \) \ 4x^2 - 3 \ x + 2 \\ - \ (4x^2 - 2 \ x) \\ - \\ - \ (4 \ x - 2) \\ - \\ 4 \end{array}$$

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In general, divide $P(x)$ by $(x - a)$ gives $Q(x)$ and remainder r .
That is, $P(x) = (x - a)Q(x) + r$

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- Intuitively, a field is a set with operations corresponding to addition, multiplication, and division.