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CS61CL : Machine Structures

Lecture #4 – MIPS I: Registers, Memory, Decisions

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Review

- Data lives in 3 places in memory
  - Stack – local variables, function parameters
  - Heap – malloc (don't forget to free!)
  - Static – global variables
- Several techniques for managing heap w/ malloc/free: best-, first-, next-fit, slab,buddy
  - 2 types of memory fragmentation: internal & external; all suffer from some kind of frag.
  - Each technique has strengths and weaknesses, none is definitively best

Assembly Language

- Basic job of a CPU: execute lots of *instructions*.
- Instructions are the primitive operations that the CPU may execute.
- Different CPUs implement different sets of instructions. The set of instructions a particular CPU implements is an *Instruction Set Architecture* (*ISA*).
  - Examples: Intel 80x86 (Pentium 4), IBM/Motorola PowerPC (Macintosh), MIPS, Intel IA64, ...

MIPS Architecture

- MIPS – semiconductor company that built one of the first commercial RISC architectures
- We will study the MIPS architecture in some detail in this class (also used in upper division courses CS 152, 162, 164)
- Why MIPS instead of Intel 80x86?
  - MIPS is simple, elegant. Don't want to get bogged down in gritty details.
  - MIPS widely used in embedded apps, x86 little used in embedded, and more embedded computers than PCs

Most HP LaserJet workgroup printers are driven by MIPS-based™ 64-bit processors.

Assembly Variables: Registers (1/4)

- Unlike HLL like C or Java, assembly cannot use variables
  - Why not? Keep Hardware Simple
- Assembly Operands are registers
  - limited number of special locations built directly into the hardware
  - operations can only be performed on these!
- Benefit: Since registers are directly in hardware, they are very fast (faster than 1 billionth of a second)

Assembly Variables: Registers (2/4)

- Drawback: Since registers are in hardware, there are a predetermined number of them
  - Solution: MIPS code must be very carefully put together to efficiently use registers
- 32 registers in MIPS
  - Why 32? Smaller is faster
- Each MIPS register is 32 bits wide
  - Groups of 32 bits called a word in MIPS

Assembly Variables: Registers (3/4)

- Registers are numbered from 0 to 31
- Each register can be referred to by number or name
- Number references:

  `$0, $1, $2, … $30, $31`

Assembly Variables: Registers (4/4)

- By convention, each register also has a name to make it easier to code
- For now:

  `$16 - $23` ➔ `$s0 - $s7`

  (correspond to C variables)

  `$8 - $15` ➔ `$t0 - $t7`

  (correspond to temporary variables)

  Later will explain other 16 register names
- In general, use names to make your code more readable

C, Java variables vs. registers

- In C (and most High Level Languages) variables declared first and given a type
  - Example:

    ```
    int fahr, celsius;
    char a, b, c, d, e;
    ```
- Each variable can ONLY represent a value of the type it was declared as (cannot mix and match `int` and `char` variables).
- In Assembly Language, the registers have no type; operation determines how register contents are treated

Comments in Assembly

- Another way to make your code more readable: comments!
- Hash (#) is used for MIPS comments
  - anything from hash mark to end of line is a comment and will be ignored
  - This is just like the C99 //
- Note: Different from C.
  - C comments have format `/* comment */` so they can span many lines

Assembly Instructions

- In assembly language, each statement (called an Instruction), executes exactly one of a short list of simple commands
- Unlike in C (and most other High Level Languages), each line of assembly code contains at most 1 instruction
- Instructions are related to operations (=, +, -, *, /) in C or Java
- Ok, enough already…gimme my MIPS!

MIPS Addition and Subtraction (1/4)

- Syntax of Instructions:
  - 1 2,3,4
  - where:
  - 1) operation by name
  - 2) operand getting result ("destination")
  - 3) 1st operand for operation ("source1")
  - 4) 2nd operand for operation ("source2")
- Syntax is rigid:
  - 1 operator, 3 operands
  - Why? Keep Hardware simple via regularity

Addition and Subtraction of Integers (2/4)

- Addition in Assembly
  - Example: `add $s0,$s1,$s2` (in MIPS)
  - Equivalent to: `a = b + c` (in C)
  - where MIPS registers `$s0,$s1,$s2` are associated with C variables `a, b, c`
- Subtraction in Assembly
  - Example: `sub $s3,$s4,$s5` (in MIPS)
  - Equivalent to: `d = e - f` (in C)
  - where MIPS registers `$s3,$s4,$s5` are associated with C variables `d, e, f`

Addition and Subtraction of Integers (3/4)

- How do the following C statement?

```
a = b + c + d - e;
```

- Break into multiple instructions

```
add $t0, $s1, $s2 # temp = b + c
add $t0, $t0, $s3 # temp = temp + d
sub $s0, $t0, $s4 # a = temp - e
```

- Notice: A single line of C may break up into several lines of MIPS.
- Notice: Everything after the hash mark on each line is ignored (comments)

Addition and Subtraction of Integers (4/4)

- How do we do this?

```
f = (g + h) - (i + j);
```

- Use intermediate temporary register

```
add $t0,$s1,$s2      # temp = g + h
add $t1,$s3,$s4      # temp = i + j
sub $s0,$t0,$t1      # f=(g+h)-(i+j)
```


Register Zero

- One particular immediate, the number zero (0), appears very often in code.
- So we define register zero (`$0` or `$zero`) to always have the value 0; eg
  - `add $s0,$s1,$zero` (in MIPS)
  - `f = g` (in C)
  - where MIPS registers `$s0,$s1` are associated with C variables `f, g`
- defined in hardware, so an instruction
  - `add $zero,$zero,$s0`
- will not do anything!

Immediates

- Immediates are numerical constants.
- They appear often in code, so there are special instructions for them.
- Add Immediate:
  - `addi $s0,$s1,10` (in MIPS)
  - `f = g + 10` (in C)
  - where MIPS registers `$s0,$s1` are associated with C variables `f, g`
- Syntax similar to `add` instruction, except that last argument is a number instead of a register.

Assembly Operands: Memory

- C variables map onto registers; what about large data structures like arrays?
- 1 of 5 components of a computer: memory contains such data structures
- But MIPS arithmetic instructions only operate on registers, never directly on memory.
- Data transfer instructions transfer data between registers and memory:
  - Memory to register
  - Register to memory

Data Transfer: Memory to Reg (1/4)

- To transfer a word of data, we need to specify two things:
  - Register: specify this by # (`$0` - `$31`) or symbolic name (`$s0,…,$t0,…`)
  - Memory address: more difficult
    - Think of memory as a single one-dimensional array, so we can address it simply by supplying a pointer to a memory address.
    - Other times, we want to be able to offset from this pointer.
- Remember: "Load FROM memory"

Data Transfer: Memory to Reg (2/4)

- To specify a memory address to copy from, specify two things:
  - A register containing a pointer to memory
  - A numerical offset (in bytes)
- The desired memory address is the sum of these two values.
- Example: `8($t0)`
  - specifies the memory address pointed to by the value in `$t0`, plus 8 bytes

Data Transfer: Memory to Reg (3/4)

- Load Instruction Syntax:

  `1    2,3(4)`

  - where
    1) operation name
    2) register that will receive value
    3) numerical offset in bytes
    4) register containing pointer to memory
- MIPS Instruction Name:
  - `lw` (meaning Load Word, so 32 bits or one word are loaded at a time)

Data Transfer: Memory to Reg (4/4)

Data flow

Example: `lw $t0,12($s0)`

This instruction will take the pointer in `$s0`, add `12` bytes to it, and then load the value from the memory pointed to by this calculated sum into register `$t0`

- Notes:
  - `$s0` is called the base register
  - `12` is called the offset
  - offset is generally used in accessing elements of array or structure: base reg points to beginning of array or structure (note offset must be a constant known at assembly time)

Data Transfer: Reg to Memory

- Also want to store from register into memory
  - Store instruction syntax is identical to Load's
- MIPS Instruction Name:

  `sw` (meaning Store Word, so 32 bits or one word is stored at a time)

Data flow

- Example: `sw $t0,12($s0)`

  This instruction will take the pointer in `$s0`, add `12` bytes to it, and then store the value from register `$t0` into that memory address
- Remember: "Store INTO memory"

Pointers v. Values

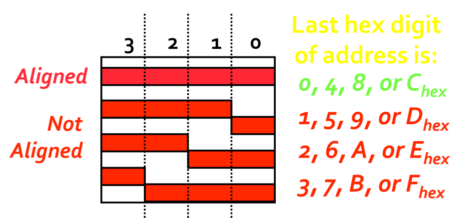
- Key Concept: A register can hold any 32-bit value. That value can be a `char`, an `int`, a pointer (memory addr), and so on
  - E.g., If you write: `add $t2,$t1,$t0` then `$t0` and `$t1` better contain values that can be added
  - E.g., If you write: `lw $t2,0($t0)` then `$t0` better contain a pointer
- Don't mix these up!

Notes about Memory

- Pitfall: Forgetting that sequential word addresses in machines with byte addressing do not differ by 1.
  - Many an assembly language programmer has toiled over errors made by assuming that the address of the next word can be found by incrementing the address in a register by 1 instead of by the word size in bytes.
  - Also, remember that for both `lw` and `sw`, the sum of the base address and the offset must be a multiple of 4 (to be word aligned)

More Notes about Memory: Alignment

- MIPS requires that all words start at byte addresses that are multiples of 4 bytes



- Called Alignment: objects fall on address that is multiple of their size

Role of Registers vs. Memory

- What if more variables than registers?
  - Compiler tries to keep most frequently used variable in registers
  - Less common variables in memory: spilling
- Why not keep all variables in memory?
  - Smaller is faster: registers are faster than memory
  - Registers more versatile:
    - MIPS arithmetic instructions can read 2, operate on them, and write 1 per instruction
    - MIPS data transfer only read or write 1 operand per instruction, and no operation

Administrivia

- HW2 due tomorrow.
- HW3 is up.
- Proj1 will be up soon… start early
- Future "Wednesday" assignments will be moved to Thursday due dates.
- Check the newsgroup often and ask there for help.

So Far...

- All instructions so far only manipulate data…we've built a calculator of sorts.
- In order to build a computer, we need ability to make decisions…
- C (and MIPS) provide labels to support "`goto`" jumps to places in code.
  - C: Horrible style; MIPS: Necessary!

MIPS Decision Instructions

- Decision instruction in MIPS:

```
beq   register1, register2, L1
```

`beq` is "Branch if (registers are) equal"
Same meaning as (using C):

```
if  (register1==register2) goto L1
```

- Complementary MIPS decision instruction

```
bne   register1, register2, L1
```

`bne` is "Branch if (registers are) not equal"
Same meaning as (using C):

```
if  (register1!=register2) goto L1
```

- Called conditional branches

MIPS Goto Instruction

- In addition to conditional branches, MIPS has an unconditional branch:

```
j label
```

- Called a Jump Instruction: jump (or branch) directly to the given label without needing to satisfy any condition
- Same meaning as (using C): `goto label`
- Technically, it's the same effect as:

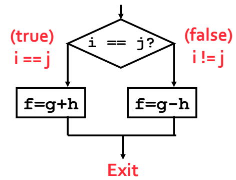
```
beq $0,$0,label
```

since it always satisfies the condition.

Compiling C `if` into MIPS (1/2)

- Compile by hand

```
if (i == j) f=g+h;
   else f=g-h;
```



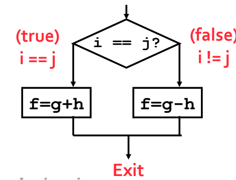
- Use this mapping:

f: $s0
g: $s1
h: $s2
i: $s3
j: $s4

Compiling C `if` into MIPS (2/2)

- Compile by hand

```
if (i == j) f=g+h;
   else f=g-h;
```



- Use this mapping:
  f: $s0 g: $s1 h: $s2 i: $s3 j: $s4
- Final compiled MIPS code:

```
      beq $s3,$s4,True  # branch i==j
      sub $s0,$s1,$s2   # f=g-h(false)
      j   Fin           # goto Fin
True: add $s0,$s1,$s2   # f=g+h (true)
Fin:
```

Note: Compiler automatically creates labels to handle decisions (branches). Generally not found in HLL code.

Loops in C/Assembly (1/3)

- Simple loop in C; `A[]` is an array of `int`s

```
do {  g = g + A[i];
      i = i + j;
} while (i !=  );
```

- Rewrite this as:

```
Loop:    g = g + A[i];
         i = i + j;
         if (i !=  ) goto Loop;
```

- Use this mapping:

```
 g,    ,  i,   j, base of A
$s1,   , $s3, $s4, $s5
```


Loops in C/Assembly (2/3)

- Final compiled MIPS code:

```
Loop:sll  $t1,$s3,2     # $t1= 4*I
     addu $t1,$t1,$s5   # $t1=addr A+4i
     lw   $t1,0($t1)    # $t1=A[i]
     addu $s1,$s1,$t1   # g=g+A[i]
     addu $s3,$s3,$s4   # i=i+j
     bne  $s3,    ,Loop # goto Loop
                        # if i!=h
```

- Original code:

```
Loop:    g = g + A[i];
         i = i + j;
         if (i !=  ) goto Loop;
```


Loops in C/Assembly (3/3)

- There are three types of loops in C:
  - `while`
  - `do… while`
  - `for`
- Each can be rewritten as either of the other two, so the method used in the previous example can be applied to these loops as well.
- Key Concept: Though there are multiple ways of writing a loop in MIPS, the key to decision-making is conditional branch

Inequalities in MIPS (1/4)

- Until now, we've only tested equalities (== and != in C). General programs need to test < and > as well.
- Introduce MIPS Inequality Instruction:
  - "Set on Less Than"
  - Syntax: `slt reg1,reg2,reg3`
  - Meaning: `reg1 = (reg2 < reg3);`

```
if (reg2 < reg3)
    reg1 = 1;
else reg1 = 0;
```

Same thing…

"set" means "change to 1", "reset" means "change to 0".

Inequalities in MIPS (2/4)

- How do we use this? Compile by hand:
  `if (g < h) goto Less; #g:$s0, h:$s1`
- Answer: compiled MIPS code…

```
    slt $t0,$s0,$s1 # $t0 = 1 if g<h
bne $t0,$0,Less  # goto Less
                 # if $t0!=0
                 # (if (g<h)) Less:
```

- Register `$0` always contains the value `0`, so `bne` and `beq` often use it for comparison after an `slt` instruction.
- A `slt` ➔ `bne` pair means `if(… < …)goto…`

Inequalities in MIPS (3/4)

- Now we can implement <, but how do we implement >, ≤ and ≥ ?
- We could add 3 more instructions, but:
  - MIPS goal: Simpler is Better
- Can we implement ≤ in one or more instructions using just `slt` and branches?
  - What about >?
  - What about ≥?

Inequalities in MIPS (4/4)

```
# a:$s0, b:$s1
slt $t0,$s0,$s1 # $t0 = 1 if a<b
beq $t0,$0,skip # skip if a >= b
    <stuff>      # do if a<b
skip:
```

Two independent variations possible:

Use `slt $t0,$s1,$s0` instead of `slt $t0,$s0,$s1`

Use `bne` instead of `beq`

Immediates in Inequalities

- There is also an immediate version of `slt` to test against constants: `slti`
  - Helpful in `for` loops

C
```
if (g >= 1) goto Loop
```

MIPS
```
Loop: . . .
slti $t0,$s0,1     # $t0 = 1 if
                   # $s0<1 (g<1)
beq  $t0,$0,Loop   # goto Loop
                   # if $t0==0
                   # (if (g>=1))
```

An `slt` ➔ `beq` pair means `if(… ≥ …)goto…`

"And in Conclusion…"

- In MIPS Assembly Language:
  - Registers replace C variables
  - One Instruction (simple operation) per line
  - Simpler is Better
  - Smaller is Faster
- New Instructions:
  `add, addi, sub`
- New Registers:
  C Variables: `$s0` - `$s7`
  Temporary Variables: `$t0` - `$t9`
  Zero: `$zero`

"And in Conclusion…"

- Memory is byte-addressable, but `lw` and `sw` access one word at a time.
- A pointer (used by `lw` and `sw`) is just a memory address, we can add to it or subtract from it (using offset).
- A Decision allows us to decide what to execute at run-time rather than compile-time.
- C Decisions are made using conditional statements within `if`, `while`, `do while`, `for`.
- MIPS Decision making instructions are the conditional branches: `beq` and `bne`.
- New Instructions:
  `lw, sw, beq, bne, j`

"And in conclusion…"

- To help the conditional branches make decisions concerning inequalities, we introduce: "Set on Less Than" called `slt`, `slti`, `sltu`, `sltiu`
- One can store and load (signed and unsigned) bytes as well as words with `lb`, `lbu`
- Unsigned add/sub don't cause overflow
- New MIPS Instructions:

Bonus slides

- These are extra slides that used to be included in lecture notes, but have been moved to this, the "bonus" area to serve as a supplement.
- The slides will appear in the order they would have in the normal presentation

Bonus

Example: The C Switch Statement (1/3)

- Choose among four alternatives depending on whether `k` has the value 0, 1, 2 or 3. Compile this C code:

```
switch (k) {
 case 0: f=i+j; break; /* k=0 */
 case 1: f=g+h; break; /* k=1 */
 case 2: f=g–h; break; /* k=2 */
 case 3: f=i–j; break; /* k=3 */
}
```


Example: The C Switch Statement (2/3)

- This is complicated, so **simplify**.
- Rewrite it as a chain of if-else statements, which we already know how to compile:

```
if(k==0) f=i+j;
  else if(k==1) f=g+h;
    else if(k==2) f=g–h;
      else if(k==3) f=i–j;
```

- Use this mapping:

```
f:$s0, g:$s1, h:$s2,
i:$s3, j:$s4, k:$s5
```


Example: The C Switch Statement (3/3)

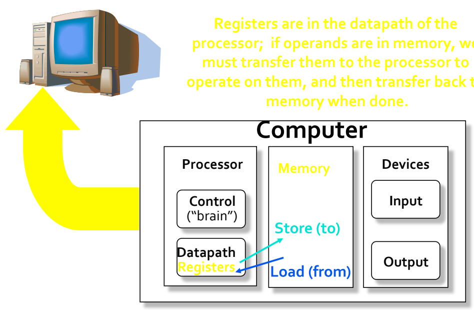
- Final compiled MIPS code:

```
      bne $s5,$0,L1     # branch k!=0
     add  $s0,$s3,$s4 #k==0 so f=i+j
     j    Exit        # end of case so Exit
L1:  addi $t0,$s5,-1  # $t0=k-1
     bne  $t0,$0,L2   # branch k!=1
     add  $s0,$s1,$s2 #k==1 so f=g+h
     j    Exit        # end of case so Exit
L2:  addi $t0,$s5,-2  # $t0=k-2
     bne  $t0,$0,L3   # branch k!=2
     sub  $s0,$s1,$s2 #k==2 so f=g-h
     j    Exit        # end of case so Exit
L3:  addi $t0,$s5,-3  # $t0=k-3
     bne  $t0,$0,Exit # branch k!=3
     sub  $s0,$s3,$s4 #k==3 so f=i-j
Exit:
```


Immediates

- There is no Subtract Immediate in MIPS: Why?
- Limit types of operations that can be done to absolute minimum
  - if an operation can be decomposed into a simpler operation, don't include it
  - `addi …, -X` = `subi …, X` => so no `subi`
- `addi $s0,$s1,-10` (in MIPS)
  `f = g - 10` (in C)
  where MIPS registers `$s0,$s1` are associated with C variables `f, g`

Anatomy: 5 components of any Computer



These are "data transfer" instructions…

Addressing: Byte vs. Word

- Every word in memory has an **address**, similar to an index in an array
- Early computers numbered words like C numbers elements of an array:
  - `Memory[0]`, `Memory[1]`, `Memory[2]`, …

  Called the "address" of a word
- Computers needed to access 8-bit bytes as well as words (4 bytes/word)
- Today machines address memory as bytes, (i.e., "Byte Addressed") hence 32-bit (4 byte) word addresses differ by 4
  - `Memory[0]`, `Memory[4]`, `Memory[8]`

Compilation with Memory

- What offset in `lw` to select `A[5]` in C?
- 4x5=20 to select `A[5]`: byte v. word
- Compile by hand using registers:
  `g = h + A[5];`
  - `g`: `$s1`, `h`: `$s2`, `$s3`: base address of `A`
- 1st transfer from memory to register:

```
lw    $t0,20($s3)  # $t0 gets A[5]
```

  - Add 20 to `$s3` to select `A[5]`, put into `$t0`
- Next add it to `h` and place in `g`

```
add $s1,$s2,$t0   # $s1 = h+A[5]
```


C Decisions: `if` Statements

- 2 kinds of `if` statements in C
  - `if` (*condition*) *clause*
  - `if` (*condition*) *clause1* `else` *clause2*
- Rearrange 2nd `if` into following:

```
  if   (condition) goto L1;
     clause2;
     goto L2;
  L1: clause1;
  L2:
```

- Not as elegant as `if-else`, but same meaning

Last time: Loading, Storing bytes 1/2

- In addition to word data transfers (`lw`, `sw`), MIPS has **byte** data transfers:
  - load byte: `lb`
  - store byte: `sb`
- same format as `lw`, `sw`
- E.g., `lb $s0, 3($s1)`
  - *contents of memory location with address = sum of "3" + contents of register `s1` is copied to the low byte position of register `s0`.*

Loading, Storing bytes 2/2

- What do with other 24 bits in the 32 bit register?
  - `lb`: sign extends to fill upper 24 bits

`xxxx xxxx xxxx xxxx xxxx xxxx xzzz zzzz`

...is copied to "sign-extend" ← This bit | byte loaded

- Normally don't want to sign extend chars
  - MIPS instruction that doesn't sign extend when loading bytes:
    - load byte unsigned: `lbu`

Overflow in Arithmetic (1/2)

- Reminder: Overflow occurs when there is a mistake in arithmetic due to the limited precision in computers.
- Example (4-bit unsigned numbers):

+15	1111
+3	0011
+18	10010

  - But we don't have room for 5-bit solution, so the solution would be 0010, which is +2, and wrong.

Overflow in Arithmetic (2/2)

- Some languages detect overflow (Ada), some don't (C)
- MIPS solution is 2 kinds of arithmetic instructs:
  - These cause overflow to be detected
    - add (`add`)
    - add immediate (`addi`)
    - subtract (`sub`)
  - These do not cause overflow detection
    - add unsigned (`addu`)
    - add immediate unsigned (`addiu`)
    - subtract unsigned (`subu`)
- Compiler selects appropriate arithmetic

What about unsigned numbers?

- Also unsigned inequality instructions:

    `sltu`, `sltiu`

...which sets result to `1` or `0` depending on unsigned comparisons

- What is value of `$t0`, `$t1`?

(`$s0 = FFFF FFFA`$_{hex}$, `$s1 = 0000 FFFA`$_{hex}$)

```
slt $t0, $s0, $s1
sltu $t1, $s0, $s1
```


MIPS Signed vs. Unsigned – diff meanings!

- MIPS terms Signed/Unsigned "overloaded":
  - Do/Don't sign extend
    - (`lb, lbu`)
  - Do/Don't overflow
    - (`add, addi, sub, mult, div`)
    - (`addu, addiu, subu, multu, divu`)
  - Do signed/unsigned compare
    - (`slt, slti/sltu, sltiu`)

Two "Logic" Instructions

- Here are 2 more new instructions
- Shift Left: `sll $s1,$s2,2 #s1=s2<<2`
  - Store in `$s1` the value from `$s2` shifted 2 bits to the left, inserting 0's on right; << in C
  - Before: `0000 0002`$_{hex}$
    `0000 0000 0000 0000 0000 0000 0000 0010`$_{two}$
  - After: `0000 0008`$_{hex}$
    `0000 0000 0000 0000 0000 0000 0000 1000`$_{two}$
  - What arithmetic effect does shift left have?
- Shift Right: `srl` is opposite shift; >>