## CS 61C:

# Great Ideas in Computer Architecture Hardware intro, Digital Logic 

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## Levels of Representation/Interpretation



## Machine Interpretation

Hardware Architecture Description (e.g., block diagrams)

Architecture Implementation

## Logic Circuit Description

 (Circuit Schematic Diagrams)```
temp = v[k];
```

$v[k]=v[k+1] ;$
$v[k+1]=$ temp;

|  | \$2) | nything can be represente |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| lw | \$t1, 4(\$2) | as a number, i.e., data or instructions |  |  |  |  |
| sw | \$t1, 0(\$2) |  |  |  |  |  |
| sW | \$t0, 4(\$2) |  |  |  |  |  |
| 00 | 10011100 | 01 | 10 | 11 | 0101 | 1000 |
| 1010 | 11110101 | 1000 | 000 | 100 | 1100 | 011 |
| 1100 | 01101010 | 1111 | 010 | 1000 | 0000 | 100 |
| 0101 | 100000 | 100 |  |  | 1010 |  |



## Hardware Design

- Next several weeks: how a modern processor is built, starting with basic elements as building blocks
- Why study hardware design?
- Understand capabilities and limitations of HW in general and processors in particular
- What processors can do fast and what they can't do fast (avoid slow things if you want your code to run fast!)
- Background for more in-depth HW courses (EECS 151, CS 152)
- Hard to know what you'll need for next 30 years
- There is only so much you can do with standard processors: you may need to design own custom HW for extra performance
- Even some commercial processors today have customizable hardware!


## Synchronous Digital Systems

Hardware of a processor, such as the MIPS, is an example of a Synchronous Digital System

## Synchronous:

- All operations coordinated by a central clock
- "Heartbeat" of the system!

Digital:

- Represent all values by discrete values
- Two binary digits: 1 and 0
- Electrical signals are treated as 1's and 0's
- 1 and 0 are complements of each other
- High /low voltage for true / false, 1 / 0


## Switches: Basic Element of Physical Implementations

- Implementing a simple circuit (arrow shows action if wire changes to " 1 " or is asserted):


$$
Z \equiv A
$$

## Switches (cont'd)

- Compose switches into more complex ones (Boolean functions):


$$
Z \equiv A \text { or } B
$$

## Historical Note

- Early computer designers built ad hoc circuits from switches
- Began to notice common patterns in their work: ANDs, ORs, ...
- Master's thesis (by Claude Shannon, 1940) made link between work and 19 ${ }^{\text {th }}$ Century Mathematician George Boole
- Called it "Boolean" in his honor
- Could apply math to give theory to hardware design, minimization, ...


## Transistors

- High voltage $\left(\mathrm{V}_{\text {dd }}\right)$ represents 1 , or true
- In modern microprocessors, Vdd ~ 1.0 Volt
- Low voltage (0 Volt or Ground) represents 0, or false
- Pick a midpoint voltage to decide if a 0 or a 1
- Voltage greater than midpoint=1
- Voltage less than midpoint $=0$
- This removes noise as signals propagate - a big advantage of digital systems over analog systems
- If one switch can control another switch, we can build a computer!
- Our switches: CMOS transistors


## CMOS Transistor Networks

- Modern digital systems designed in CMOS
- MOS: Metal-Oxide on Semiconductor
- C for complementary: use pairs of normally-on and normally-off switches
- CMOS transistors act as voltage-controlled switches
- Similar, though easier to work with, than electromechanical relay switches from earlier era
- Use energy primarily when switching


## CMOS Transistors



- Three terminals: source, gate, and drain
- Switch action: if voltage on gate terminal is (some amount) higher/lower than source terminal then conducting path established between drain and source terminals (switch is closed)

$n$-channel transitor
off when voltage at Gate is low
on when:
voltage(Gate) > voltage (Threshold)

p-channel transistor
on when voltage at Gate is low off when:
voltage(Gate) > voltage (Threshold)



## Intel 14nm Technology



Side view of wiring layers

## Sense of Scale



Source: Mark Bohr, IDF14

## CMOS Circuit Rules

- Don't pass weak values => Use Complementary Pairs
- N-type transistors pass weak 1's ( $\left.\mathrm{V}_{\mathrm{dd}}-\mathrm{V}_{\mathrm{th}}\right)$
- N-type transistors pass strong 0's (ground)
- Use N-type transistors only to pass O's (N for negative)
- Converse for P-type transistors: Pass weak 0s, strong 1s
- Pass weak 0's $\left(\mathrm{V}_{\mathrm{th}}\right)$, strong 1's ( $\mathrm{V}_{\mathrm{dd}}$ )
- Use P-type transistors only to pass 1's (P for positive)
- Use pairs of N-type and P-type to get strong values
- Never leave a wire undriven
- Make sure there's always a path to $\mathrm{V}_{\mathrm{dd}}$ or GND
- Never create a path from $\mathrm{V}_{\mathrm{dd}}$ to GND (ground)
- This would short-circuit the power supply!


## CMOS Networks

p-channel transistor
on when voltage at Gate is low off when:
voltage(Gate) > voltage (Threshold)

$n$-channel transitor
off when voltage at Gate is low on when:
voltage(Gate) > voltage (Threshold)
Called an inverter or not gate

## Two-Input Networks



Called a NAND gate
(NOT AND)

## Clickers/Peer Instruction




## Administrivia

- Proj2-1 due Sunday 2/21!
- MT1 next Thursday night
- If you are in DSP or have an exam conflict, you should have received an email from Fred. If not, email Fred and William.
- EE16B students will take it at the normal time
- We will post room assignments on Piazza over the weekend
- Proj2-2 will be released late next week.


## Combinational Logic Symbols

- Common combinational logic systems have standard symbols called logic gates
- Buffer, NOT



## Boolean Algebra

- Use plus "+" for OR
- "logical sum"
- Use product for AND ( $a \bullet b$ or implied via $a b$ )
- "logical product"
- "Hat" to mean complement (NOT)
- Thus
$a b+a+\bar{c}$
$=a \bullet b+a+\bar{c}$
$=(\mathrm{a}$ AND b) OR a OR (NOT c )


## Representations of Combinational

 Logic (groups of logic gates)


## Truth Table Example \#1: $y=F(a, b): 1$ iff $a \neq b$

| $a$ | $b$ | $y$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |

Truth Table Example \#2: 2-bit Adder


How
Many
Rows?

## Truth Table Example \#3: 32-bit Unsigned Adder

| A | B | C |  |
| :---: | :---: | :--- | :--- |
| $000 \ldots 0$ | $000 \ldots 0$ | $000 \ldots 00$ |  |
| $000 \ldots 0$ | $000 \ldots 1$ | $000 \ldots 01$ |  |
| . | . | $\cdot$ | How |
| . | . | $\cdot$ | Many |
| . | . | - | Rows? |
| $111 \ldots 1$ | $111 \ldots 1$ | $111 \ldots 10$ |  |

## Truth Table Example \#4: 3-input Majority Circuit

$Y=$
This is called Sum of Products form; Just another way to represent the TT as a logical expression

More simplified forms (fewer gates and wires)

| a | b | c | y |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 |

## And in Conclusion, ...

- Multiple Hardware Representations
- Analog voltages quantized to represent logic 0 and logic 1
- Transistor switches form gates: AND, OR, NOT, NAND, NOR
- Truth table mapped to gates for combinational logic design
- Boolean algebra for gate minimization

