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UCB CS61C : Machine Structures

Lecture 36 – Performance
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How fast is your computer?

CRAY XT5-HE IS FASTEST SUPERCOMPUTER!

Every 6 months (Nov/June), the fastest supercomputers in the world face off. The fastest computer is now the Jaguar, a Linux Cray XT5-HE Opteron Six Core 2.6 GHz, with 224,256 cores, achieves 2.3 PFlops. They use LINPACK floating point benchmark (A x = B).

www.top500.org/lists/2009/11
www.nccs.gov/computing-resources/jaguar/

Why Performance? Faster is better!

- **Purchasing Perspective**: given a collection of machines (or upgrade options), which has the
  - best performance ?
  - least cost ?
  - best performance / cost ?
- **Computer Designer Perspective**: faced with design options, which has the
  - best performance improvement ?
  - least cost ?
  - best performance / cost ?
- All require basis for comparison and metric for evaluation!
  - Solid metrics lead to solid progress!

Two Notions of "Performance"

Plane	DC to Paris	Top Speed	Passen-gers	Throughput (pmph)
Boeing 747	6.5 hours	610 mph	470	286,700
BAD/Sud Concorde	3 hours	1350 mph	132	178,200

- Which has higher performance?
  - Interested in time to deliver 100 passengers?
  - Interested in delivering as many passengers per day as possible?
- In a computer, time for one task called
  Response Time or Execution Time
- In a computer, tasks per unit time called
  Throughput or Bandwidth

Definitions

- Performance is in units of things per sec
  - bigger is better
- If mostly concerned with response time
  - $\text{performance}(x) = \frac{1}{\text{execution\_time}(x)}$
- " F(ast) is *n* times faster than S(low) " means:

$$n = \frac{\text{performance(F)}}{\text{performance(S)}} = \frac{\text{execution\_time(S)}}{\text{execution\_time(F)}}$$

Example of Response Time v. Throughput

- **Time of Concorde vs. Boeing 747?**
  - Concord is 6.5 hours / 3 hours
    = 2.2 times faster
  - Concord is 2.2 times ("120%") faster in terms of flying time (response time)
- **Throughput of Boeing vs. Concorde?**
  - Boeing 747: 286,700 pmph / 178,200 pmph
    = 1.6 times faster
  - Boeing is 1.6 times ("60%") faster in terms of throughput
- **We will focus primarily on response time.**

Words, Words, Words…

- Will (try to) stick to "n times faster"; its less confusing than "m % faster"
- As faster means both decreased execution time and increased performance, to reduce confusion we will (and you should) use "improve execution time" or "improve performance"

What is Time?

- **Straightforward definition of time:**
  - Total time to complete a task, including disk accesses, memory accesses, I/O activities, operating system overhead, ...
  - "real time", "response time" or "elapsed time"
- **Alternative: just time processor (CPU) is working only on your program (since multiple processes running at same time)**
  - "CPU execution time" or "CPU time"
  - Often divided into system CPU time (in OS) and user CPU time (in user program)

How to Measure Time?

- **Real Time ⇒ Actual time elapsed**

- **CPU Time: Computers constructed using a clock that runs at a constant rate and determines when events take place in the hardware**
  - These discrete time intervals called clock cycles (or informally clocks or cycles)
  - Length of clock period: clock cycle time (e.g., ½ nanoseconds or ½ ns) and clock rate (e.g., 2 gigahertz, or 2 GHz), which is the inverse of the clock period; use these!

Measuring Time using Clock Cycles (1/2)

- CPU execution time for a program
  - Units of [seconds / program] or [s/p]

= Clock Cycles for a program x Clock Period

  - Units of [s/p] = [cycles / p] x [s / cycle] = [c/p] x [s/c]

- Or

$$= \frac{\text{Clock Cycles for a program [c / p]}}{\text{Clock Rate [c / s]}}$$

Measuring Time using Clock Cycles (2/2)

- One way to define clock cycles:

Clock Cycles for program [c/p]

= Instructions for a program [i/p]
(called "Instruction Count")

x Average Clock cycles Per Instruction [c/i]
(abbreviated "CPI")

- CPI one way to compare two machines with same instruction set, since Instruction Count would be the same

Performance Calculation (1/2)

- CPU execution time for program [s/p]
  = Clock Cycles for program [c/p]
  x Clock Cycle Time [s/c]

- Substituting for clock cycles:

  CPU execution time for program [s/p]
  = (Instruction Count [i/p] x CPI [c/i])
  x Clock Cycle Time [s/c]

= Instruction Count x CPI x Clock Cycle Time

Performance Calculation (2/2)

$$\text{CPU time} = \frac{\text{Instructions}}{\text{Program}} \times \frac{\text{Cycles}}{\text{Instruction}} \times \frac{\text{Seconds}}{\text{Cycle}}$$

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$$\text{CPU time} = \frac{\text{Seconds}}{\text{Program}}$$

Product of all 3 terms: if missing a term, can't predict time, the real measure of performance

How Calculate the 3 Components?

- **Clock Cycle Time: in specification of computer (Clock Rate in advertisements)**
- **Instruction Count:**
  - Count instructions in loop of small program
  - Use simulator to count instructions
  - Hardware counter in spec. register
    - (Pentium II,III,4)
- **CPI:**
  - Calculate: $\frac{\text{Execution Time / Clock cycle time}}{\text{Instruction Count}}$
  - Hardware counter in special register (PII,III,4)

Calculating CPI Another Way

- First calculate CPI for each individual instruction (`add`, `sub`, `and`, etc.)
- Next calculate frequency of each individual instruction
- Finally multiply these two for each instruction and add them up to get final CPI (the weighted sum)

Example (RISC processor)

Op	$Freq_i$	CPI_i	Prod	(% Time)
ALU	50%	1	.5	(23%)
Load	20%	5	1.0	(45%)
Store	10%	3	.3	(14%)
Branch	20%	2	.4	(18%)
			2.2	

Instruction Mix

(Where time spent)

- What if Branch instructions twice as fast?

What Programs Measure for Comparison?

- **Ideally run typical programs with typical input before purchase,**
  or before even build machine
  - Called a "workload"; For example:
  - Engineer uses compiler, spreadsheet
  - Author uses word processor, drawing program, compression software
- **In some situations its hard to do**
  - Don't have access to machine to "benchmark" before purchase
  - Don't know workload in future

Benchmarks

- Obviously, apparent speed of processor depends on code used to test it
- Need industry standards so that different processors can be fairly compared
- Companies exist that create these benchmarks: "typical" code used to evaluate systems
- Need to be changed every ~5 years since designers could (and do!) target for these standard benchmarks

Example Standardized Benchmarks (1/2)

- **Standard Performance Evaluation Corporation (SPEC) SPEC CPU2006**
  - CINT2006 12 integer (perl, bzip, gcc, go, ...)
  - CFP2006 17 floating-point (povray, bwaves, ...)
  - All relative to base machine (which gets 100)
    Sun Ultra Enterprise 2 w/296 MHz UltraSPARC II
  - They measure
    - System speed (SPECint2006)
    - System throughput (SPECint_rate2006)
  - **`www.spec.org/osg/cpu2006/`**

Example Standardized Benchmarks (2/2)

- SPEC
  - Benchmarks distributed in source code
  - Members of consortium select workload
    - 30+ companies, 40+ universities, research labs
  - Compiler, machine designers target benchmarks, so try to change every 5 years
  - SPEC CPU2006:

CINT2006

perlbench	C	Perl Programming language
bzip2	C	Compression
gcc	C	C Programming Language Compiler
mcf	C	Combinatorial Optimization
gobmk	C	Artificial Intelligence : Go
hmmer	C	Search Gene Sequence
sjeng	C	Artificial Intelligence : Chess
libquantum	C	Simulates quantum computer
h264ref	C	H.264 Video compression
omnetpp	C++	Discrete Event Simulation
astar	C++	Path-finding Algorithms
xalancbmk	C++	XML Processing

CFP2006

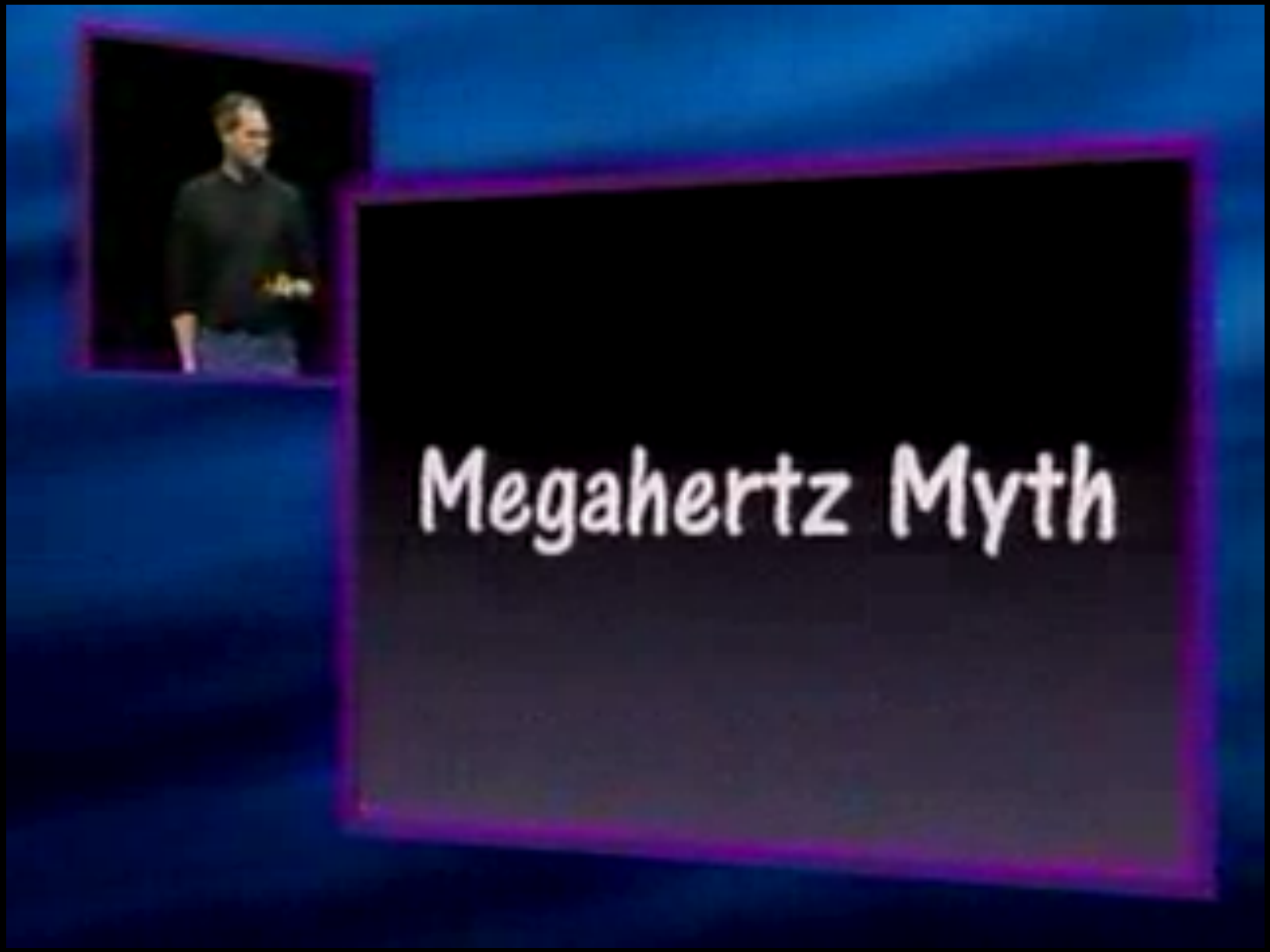
bwaves	Fortran	Fluid Dynamics
gamess	Fortran	Quantum Chemistry
milc	C	Physics / Quantum Chromodynamics
zeusmp	Fortran	Physics / CFD
gromacs	C,Fortran	Biochemistry / Molecular Dynamics
cactusADM	C,Fortran	Physics / General Relativity
leslie3d	Fortran	Fluid Dynamics
namd	C++	Biology / Molecular Dynamics
dealII	C++	Finite Element Analysis
soplex	C++	Linear Programming, Optimization
povray	C++	Image Ray-tracing
calculix	C,Fortran	Structural Mechanics
GemsFDTD	Fortran	Computational Electromegnetics
tonto	Fortran	Quantum Chemistry
lbm	C	Fluid Dynamics
wrf	C,Fortran	Weather
sphinx3	C	Speech recognition

Performance Evaluation: The Demo

If we're talking about performance, let's discuss the ways shady salespeople have fooled consumers (so you don't get taken!)

5. Never let the user touch it
4. Only run the demo through a script
3. Run it on a stock machine in which "no expense was spared"
2. Preprocess all available data
1. Play a movie

Megahertz Myth Marketing Movie



Peer Instruction

1) The Sieve of Eratosthenes and Quicksort were early effective benchmarks.
2) A program runs in 100 sec. on a machine, `mult` accounts for 80 sec. of that. If we want to make the program run 6 times faster, we need to up the speed of `mult`s by AT LEAST 6.

	12
a)	FF
b)	FT
c)	TF
d)	TT

Peer Instruction Answers

1) The Sieve of Eratosthenes and Quicksort were early effective benchmarks. FALSE

2) A program runs in 100 sec. on a machine, `mult` accounts for 80 sec of that. If we want to make the program run 6 times faster, we need to up the speed of `mult`s by AT LEAST 6. FALSE

1. Early benchmarks? Yes.
   Effective? No. Too simple!
2. 6 times faster = 16 sec.
   mults must take -4 sec!
   I.e., impossible!

12
a) FF
b) FT
c) TF
d) TT

"And in conclusion…"

$$\text{CPU time} = \frac{\text{Instructions}}{\text{Program}} \times \frac{\text{Cycles}}{\text{Instruction}} \times \frac{\text{Seconds}}{\text{Cycle}}$$

- Latency v. Throughput
- Performance doesn't depend on any single factor: need Instruction Count, Clocks Per Instruction (CPI) and Clock Rate to get valid estimations
- User Time: time user waits for program to execute: depends heavily on how OS switches between tasks
- CPU Time: time spent executing a single program: depends solely on design of processor (datapath, pipelining effectiveness, caches, etc.)
- Benchmarks
  - Attempt to predict perf, Updated every few years
  - Measure everything from simulation of desktop graphics programs to battery life
- Megahertz Myth
  - MHz ≠ performance, it's just one factor