

- Memory Mapped I/O**

Certain memory addresses correspond to registers in I/O devices and not normal memory.

**Control Register:** Indicates if it is okay to read/write data register

**Data Register:** Contains I/O data

Register	Location	Contains
Receiver Control	0xffff0000	Lowest two bits: Interrupt Enable Bit, Ready Bit
Receiver Data	0xffff0004	Received data stored at lowest byte
Transmitter Control	0xffff0008	Lowest two bits: Interrupt Enable Bit, Ready Bit
Transmitter Data	0xffff000c	Transmitted data stored at lowest byte

**Write MIPS code to read a byte from the receiver and immediately send it to the transmitter.**

- Polling and Interrupts**

Operation	Definition	Pro/Con	Good For
<b>Polling</b>			
<b>Interrupts</b>			

Interrupt Enable: Indicates whether or not to cause an interrupt when the ready bit is set

- Network Overhead vs. Bandwidth**

Assume we have two networks A and B.

Network A has a 200us overhead and a peak bandwidth of 10MB/s

Network B has a 500us overhead and a peak bandwidth of 100MB/s

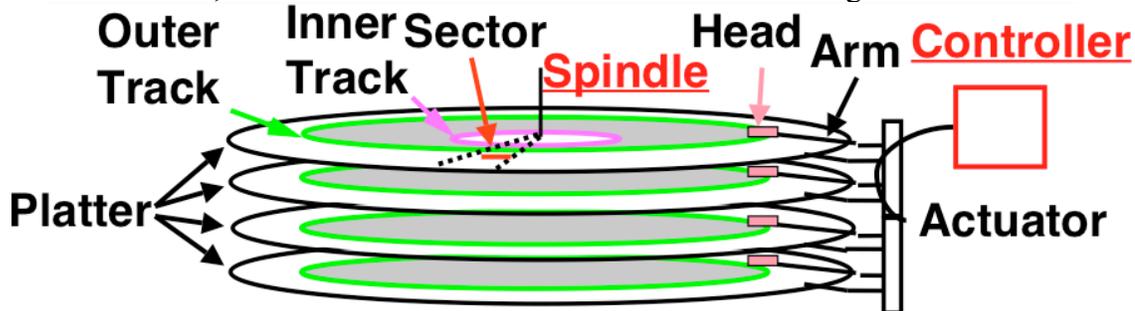
(MB = MebiByte)

**How long would it take to send 1000 Bytes over each network?**

**Which network is better for sending large amounts of data?**

## Disk Organization

Magnetic disks are one of the most common types of I/O devices. Bits are encoded by the controlling the polarity of magnetic fields on some sort of substrate. Since the magnetic fields do not require power to be maintained, disks are considered a form of non-volatile storage.



An additional term not shown here is that the collection of corresponding tracks across all the platters is called a cylinder.

There are two ways to address disks. Logical addressing treats the disk drive as one big array of blocks. Physical addressing uses a (cylinder, sector, platter) tuple to specify a blocks physical position in the disk drive.

## Disk Performance

Disk Latency = Seek Time + Rotation Time + Transfer Time + Controller Overhead

## RAID

Big disks are expensive (and dangerous). We can use an array of smaller disks to simulate the behavior of one larger disk with a more reasonable cost.

RAID 0	No redundancy, just multiple disks
RAID 1	Mirroring for redundancy, doubles read bandwidth
RAID 2	Bit-level striping, increases bandwidth further
RAID 3	Parity Disks, allows recovery from a single disk failure
RAID 4	Block-level striping with Parity Disk, increases bandwidth
RAID 5+	Striped Parity, reduces wear and tear

## Performance Metrics

In order to get any meaningful definition of performance (used for comparison), we need to develop a quantitative metric that we all can agree on. This is harder than it sounds. We briefly talked about these when discussing pipelining.

**Response Time/Execution Time/Latency** - the time it takes to complete one task

**Throughput/Bandwidth** - tasks completed per time unit

We say A has  $n$  times better performance than B if  $n = \text{performance}(A) / \text{performance}(B)$

Avoid using "percent" comparisons and "increased/decreased" to compare performances. This can be confusing and misleading. Instead we talk about performance improvements.

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Based on notes by Aaron Staley, which were in turn based on notes by David Jacobs.

**Megahertz Myth**

A processors performance is determined by more than just the clock speed.  
 CPU time = Instruction Count \* CPI \* clock period

**Exercises**

You are the lead developer of a new video game at AE, Inc. The graphics are quite sexy, but the frame rates (performance) are horrible. Doubly unfortunately, you have to show it off at a shareholder meeting tomorrow. What do you do?

You need to render your latest and greatest über-133t animation. If your rendering software contains the following mix of instructions, which processor is the best choice?

Operation	Frequency
ALU	30%
Load	30%
Store	20%
Branch	20%

A's CPI	B's CPI	C's CPI
1	1	1
3	5	3
2	3	4
3	2	2

What if the processors had different clock speeds? Assume A is a 1 Ghz processor, B is a 1.5 Ghz processor, and C is a 750 Mhz processor.

But wait, these processors are made by different manufacturers, and use different instruction sets. So the renderer (for the different architectures) takes a different number of instructions on each. Which is best if your main loop on A averages 1000 instructions; on B it averages 800 instructions; and on C it averages 1200 instructions?

**Parallel Computing**

Parallel computing refers more to multicore and multiprocessor machines. This is sometimes also called "supercomputing." Since the processors are physically closer together, there is a potential for much faster communications between them. However, synchronizing the processors can prove a difficult problem.

**Amdahl's Law**

The potential speedup from parallelization is limited by the amount a program can be parallelized. Let  $s$  be the fraction of the work that must be done sequentially and  $P$  be the number of processors. Then,

$$\text{Speedup}(P) \leq 1/s$$

**Exercise**

What are the contributing factors to Amdahl's law? Why isn't it an equality?