

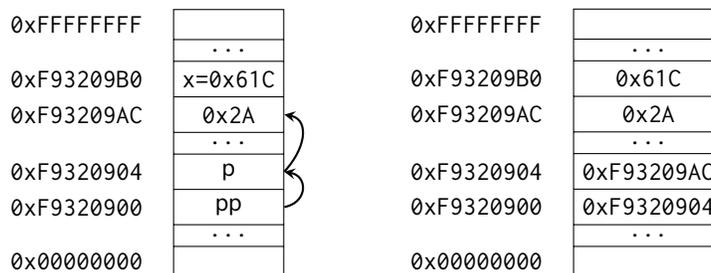
# 1 C

C is syntactically similar to Java, but there are a few key differences:

1. C is function-oriented, not object-oriented; there are no objects.
2. C does not automatically handle memory for you.
  - Stack memory, or *things allocated the way you're accustomed to*: data is garbage immediately after the *function in which it was defined* returns.
  - Heap memory, or *things allocated with malloc, calloc, or realloc commands*: data is freed only when the programmer explicitly frees it!
  - In any case, allocated memory always holds garbage until it is initialized!
3. C uses pointers explicitly. `*p` tells us to use the value that `p` points to, rather than the value of `p`, and `&x` gives the address of `x` rather than the value of `x`.

On the left is the memory represented as a box-and-pointer diagram.

On the right, we see how the memory is really represented in the computer.



Let's assume that `int* p` is located at `0xF9320904` and `int x` is located at `0xF93209B0`. As we can observe:

- `*p` should return `0x2A` ( $42_{10}$ ).
- `p` should return `0xF93209AC`.
- `x` should return `0x61C`.
- `&x` should return `0xF93209B0`.

Let's say we have an `int **pp` that is located at `0xF9320900`.

1.1 What does `pp` evaluate to? How about `*pp`? What about `**pp`?

`pp` evaluates to `0xF9320904`. `*pp` evaluates to `0xF93209AC`. `**pp` evaluates to `0x2A`.

1.2 The following functions are syntactically-correct C, but written in an incomprehensible style. Describe the behavior of each function in plain English.

- (a) Recall that the ternary operator evaluates the condition before the `?` and returns the value before the colon (`:`) if true, or the value after it if false.

```
1 int foo(int *arr, size_t n) {
2     return n ? arr[0] + foo(arr + 1, n - 1) : 0;
3 }
```

Returns the sum of the first  $N$  elements in `arr`.

- (b) Recall that the negation operator, `!`, returns 0 if the value is non-zero, and 1 if the value is 0. The `~` operator performs a *bitwise not* (NOT) operation.

```
1 int bar(int *arr, size_t n) {
2     int sum = 0, i;
3     for (i = n; i > 0; i--)
4         sum += !arr[i - 1];
5     return ~sum + 1;
6 }
```

Returns  $-1$  times the number of zeroes in the first  $N$  elements of `arr`.

- (c) Recall that `^` is the *bitwise exclusive-or* (XOR) operator.

```
1 void baz(int x, int y) {
2     x = x ^ y;
3     y = x ^ y;
4     x = x ^ y;
5 }
```

Ultimately does not change the value of either `x` or `y`.

## 2 Programming with Pointers

2.1 Implement the following functions so that they work as described.

- (a) Swap the value of two `ints`. *Remain swapped after returning from this function.*

```
1 void swap(int *x, int *y) {
2     int temp = *x;
3     *x = *y;
4     *y = temp;
5 }
```

- (b) Return the number of bytes in a string. *Do not use `strlen`.*

```
1 int mystrlen(char* str) {
2     int count = 0;
3     while (*str++) {
4         count++;
5 }
```

```

5     }
6     return count;
7 }

```

2.2 The following functions may contain logic or syntax errors. Find and correct them.

- (a) Returns the sum of all the elements in `summands`.

It is necessary to pass a size alongside the pointer.

```

1 int sum(int* summands, size_t n) {
2     int sum = 0;
3     for (int i = 0; i < n; i++)
4         sum += *(summands + i);
5     return sum;
6 }

```

- (b) Increments all of the letters in the `string` which is stored at the front of an array of arbitrary length, `n >= strlen(string)`. Does not modify any other parts of the array's memory.

The ends of strings are denoted by the null terminator rather than `n`. Simply having space for `n` characters in the array does not mean the string stored inside is also of length `n`.

```

1 void increment(char* string) {
2     for (i = 0; string[i] != 0; i++)
3         string[i]++; // or (*(string + i))++;
4 }

```

Another common bug to watch out for is the corner case that occurs when incrementing the character with the value `0xFF`. Adding 1 to `0xFF` will overflow back to 0, producing a null terminator and unintentionally shortening the string.

- (c) Copies the string `src` to `dst`.

```

1 void copy(char* src, char* dst) {
2     while (*dst++ = *src++);
3 }

```

No errors.

- (d) Overwrites an input string `src` with "61C is awesome!" if there's room. Does nothing if there is not. Assume that `length` correctly represents the length of `src`.

```

1 void cs61c(char* src, size_t length) {
2     char *srcptr, replaceptr;
3     char replacement[16] = "61C is awesome!";
4     srcptr = src;
5     replaceptr = replacement;
6     if (length >= 16) {

```

```

7         for (int i = 0; i < 16; i++)
8             *srcptr++ = *replaceptr++;
9     }
10 }
```

`char *srcptr, replaceptr` initializes a `char` pointer, and a `char`—not two `char` pointers.

The correct initialization should be, `char *srcptr, *replaceptr`.

### 3 Memory Management

3.1 For each part, choose one or more of the following memory segments where the data could be located: **code**, **static**, **heap**, **stack**.

(a) Static variables

Static

(b) Local variables

Stack

(c) Global variables

Static

(d) Constants

Code, static, or stack

Constants can be compiled directly into the code. `x = x + 1` can compile with the number 1 stored directly in the machine instruction in the code. That instruction will always increment the value of the variable `x` by 1, so it can be stored directly in the machine instruction without reference to other memory. This can also occur with pre-processor macros.

```

1 #define y 5
2
3 int plus_y(int x) {
4     x = x + y;
5     return x;
6 }
```

Constants can also be found in the stack or static storage depending on if it's declared in a function or not.

```

1 const int x = 1;
2
3 int sum(int* arr) {
4     int total = 0;
5     ...
6 }
```

In this example, `x` is a variable whose value will be stored in the static storage, while `total` is a local variable whose value will be stored on the stack. Variables declared `const` are not allowed to change, but the usage of `const` can get more tricky when combined with pointers.

- (e) Machine Instructions

Code

- (f) Result of `malloc`

Heap

- (g) String Literals

Static or stack.

When declared in a function, string literals can be stored in different places. `char* s = "string"` is stored in the static memory segment while `char[7] s = "string"` will be stored in the stack.

**3.2** Write the code necessary to allocate memory on the heap in the following scenarios

- (a) An array `arr` of  $k$  integers

```
arr = (int *) malloc(sizeof(int) * k);
```

- (b) A string `str` containing  $p$  characters

```
str = (char *) malloc(sizeof(char) * (p + 1));
```

Don't forget the null terminator!

- (c) An  $n \times m$  matrix `mat` of integers initialized to zero.

```
mat = (int *) calloc(n * m, sizeof(int));
```

Alternative solution. This might be needed if you wanted to efficiently permute the rows of the matrix.

```
1 mat = (int **) calloc(n, sizeof(int *));
2 for (int i = 0; i < n; i++)
3     mat[i] = (int *) calloc(m, sizeof(int));
```

Suppose we've defined a linked list `struct` as follows. Assume `*lst` points to the first element of the list, or is `NULL` if the list is empty.

```
struct ll_node {
    int first;
    struct ll_node* rest;
}
```

**3.3** Implement `prepend`, which adds one new value to the front of the linked list.

```
1 void prepend(struct ll_node** lst, int value) {
2     struct ll_node* item = (struct ll_node*) malloc(sizeof(struct ll_node));
```

```
3     item->first = value;
4     item->rest = *lst;
5     *lst = item;
6 }
```

3.4 Implement `free_ll`, which frees all the memory consumed by the linked list.

```
1 void free_ll(struct ll_node** lst) {
2     if (*lst) {
3         free_ll(&((*lst)->rest));
4         free(*lst);
5     }
6     *lst = NULL; // Make writes to **lst fail instead of writing to unusable memory.
7 }
```