inst.eecs.berkeley.edu/~cs61c

CS61C : Machine Structures

Lecture #18
Single Cycle CPU Control

CPS today!

2005-11-02

There is one handout today at the front and back of the room!

Lecturer PSOE, new dad Dan Garcia

www.cs.berkeley.edu/~ddgarcia

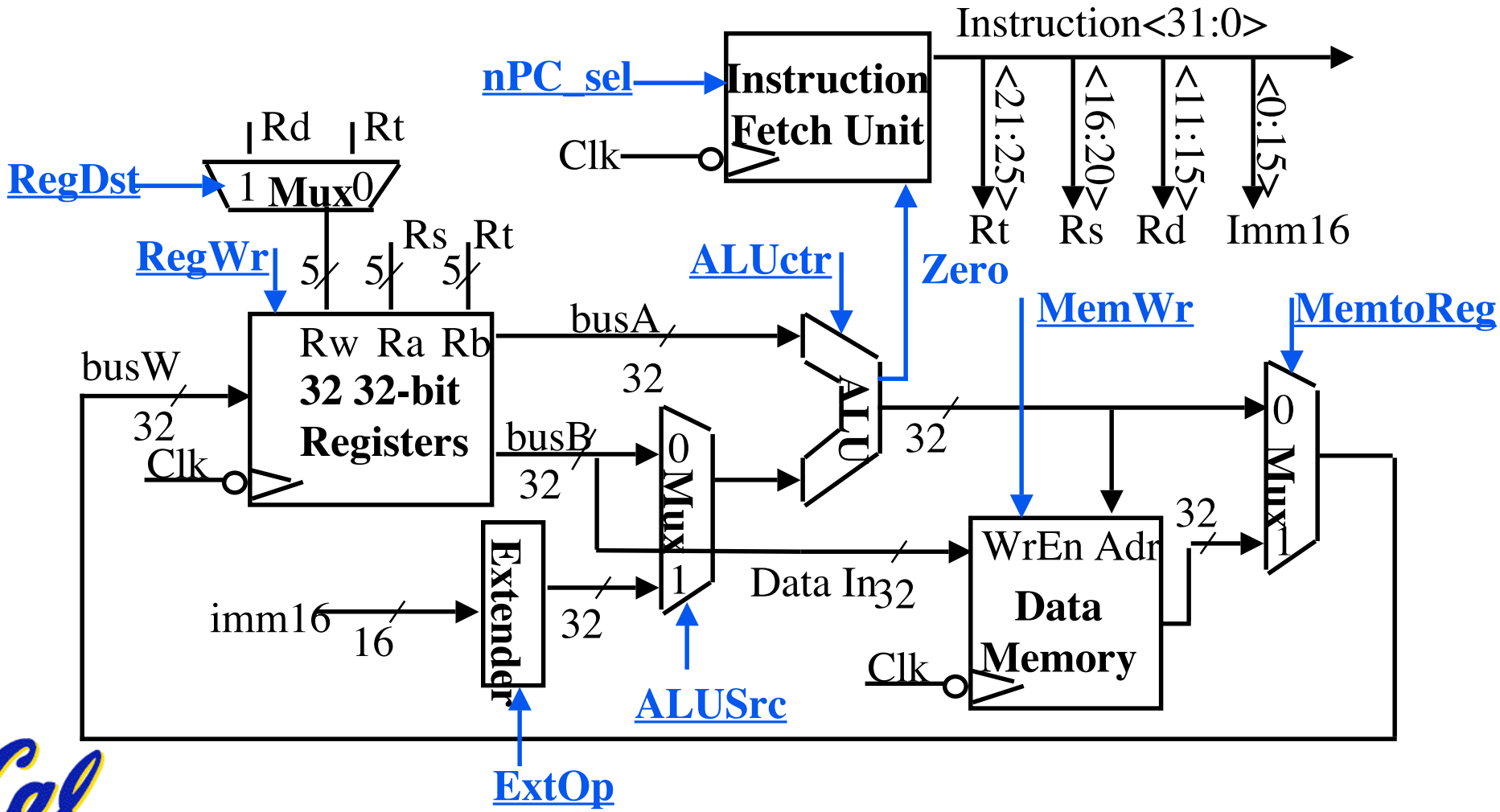
Data Transfer Record! ⇒

A Japanese company (Kansai Electric) claims fiber-optic cables on power-transmitting steel towers have achieved 1 Terabit/sec, 100 times faster than normal.

news.softpedia.com/news/Japanese-Claim-New-Fiber-Optic-Transmission-Record-11257.shtml

Summary: A Single Cycle Datapath

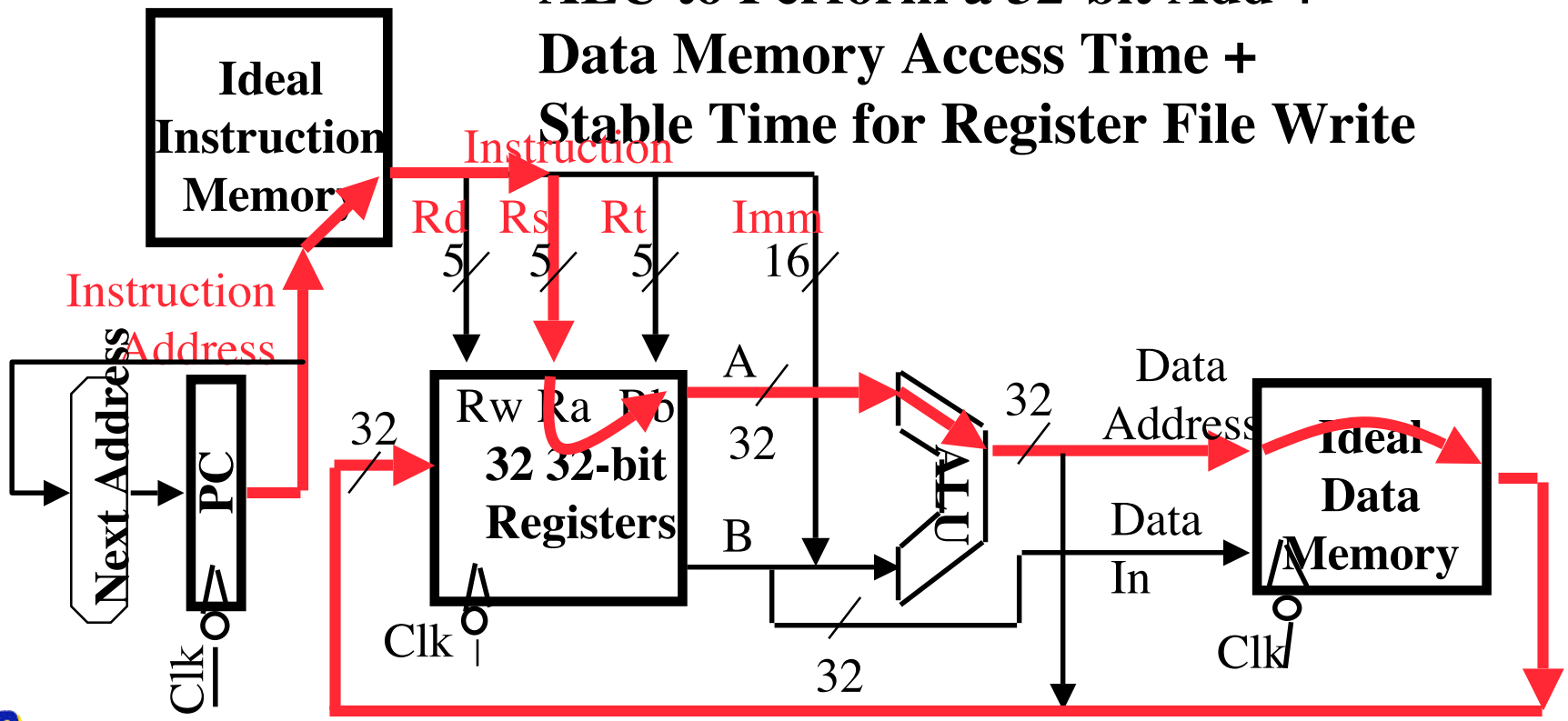
- Rs, Rt, Rd, Imed16 connected to datapath
- We have everything except control signals



An Abstract View of the Critical Path

- This affects how much you can overclock your PC!

Critical Path (Load Operation) =
Delay clock through PC (FFs) +
Instruction Memory's Access Time +
Register File's Access Time, +
ALU to Perform a 32-bit Add +
Data Memory Access Time +
Stable Time for Register File Write

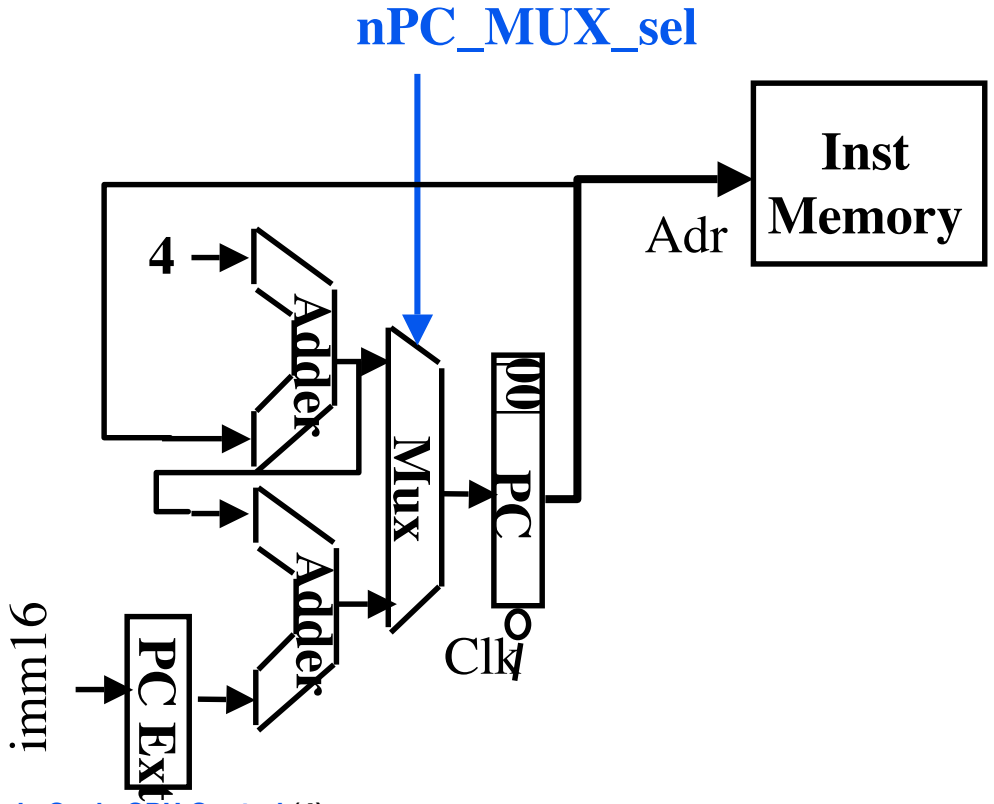


Recap: Meaning of the Control Signals

- **nPC_MUX_sel:** **0 ⇒ PC <– PC + 4**
  1 ⇒ PC <– PC + 4 + {SignExt(Im16) , 00 }

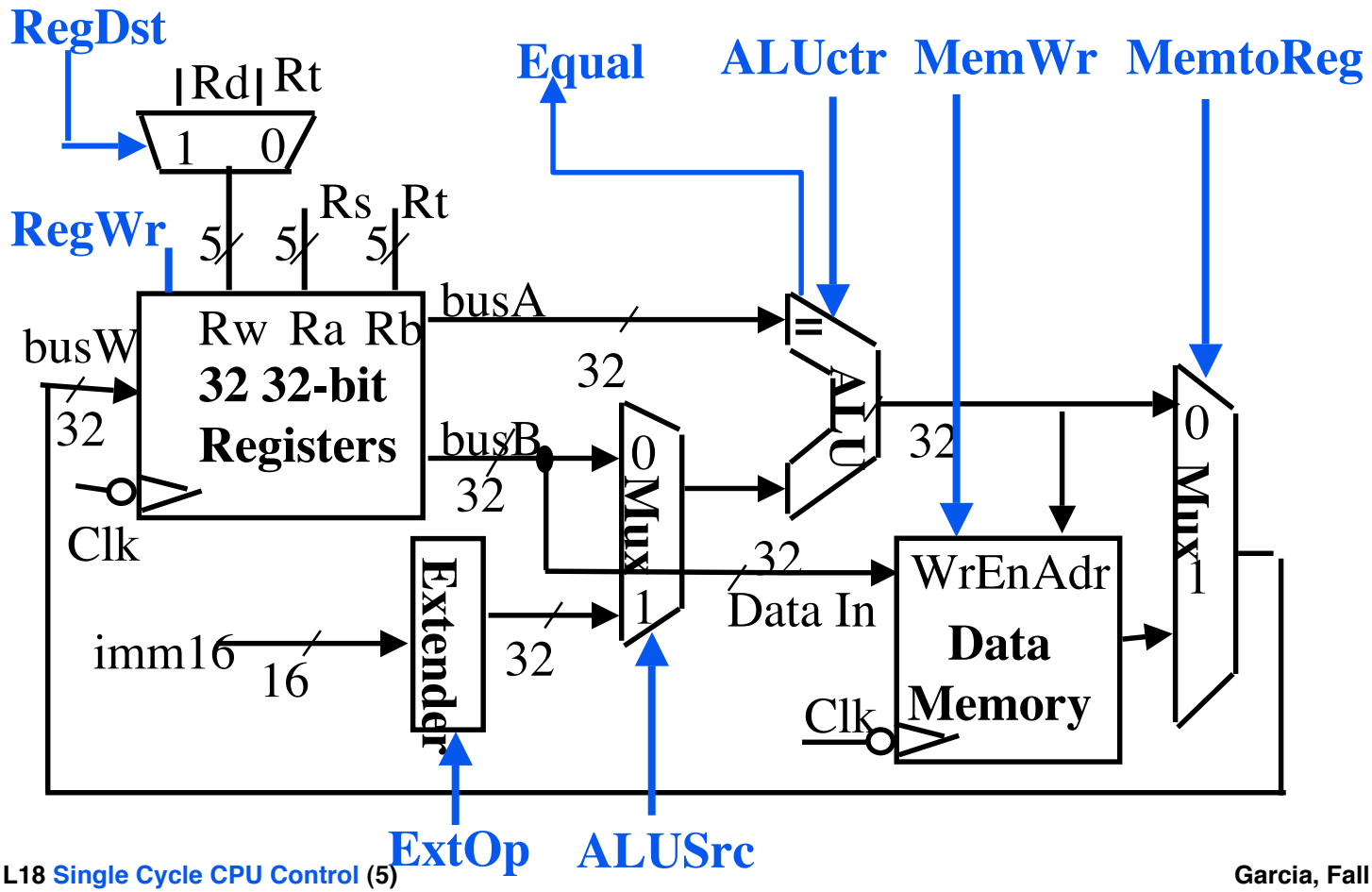
  "n"=next

- **Later in lecture: higher-level connection between mux and branch cond**

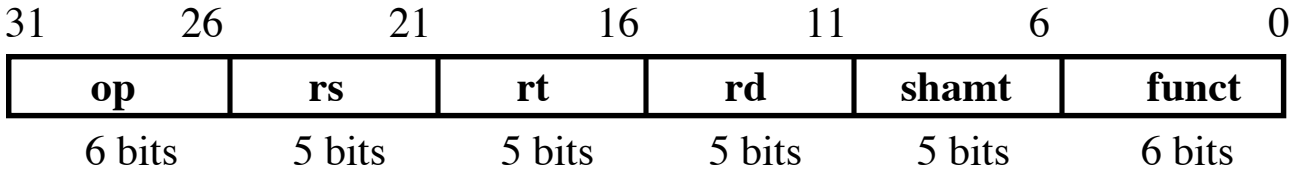


Recap: Meaning of the Control Signals

- **ExtOp:** "zero", "sign"
- **ALUsrc:** 0 ⇒ regB; 1 ⇒ immed
- **ALUctr:** "add", "sub", "or"
- **MemWr:** 1 ⇒ write memory
- **MemtoReg:** 0 ⇒ ALU; 1 ⇒ Mem
- **RegDst:** 0 ⇒ "rt"; 1 ⇒ "rd"
- **RegWr:** 1 ⇒ write register



RTL: The `Add` Instruction

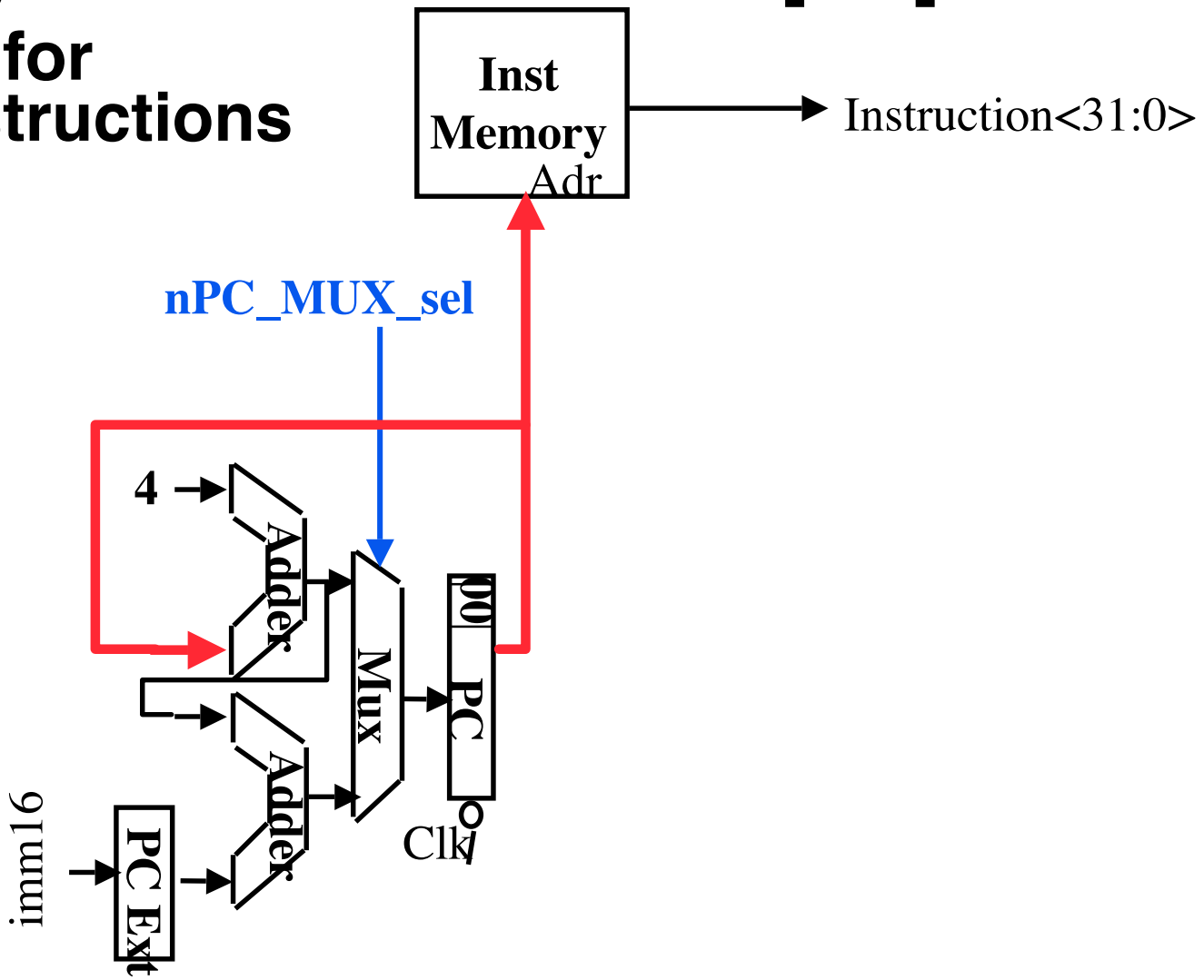



```
add rd, rs, rt
```

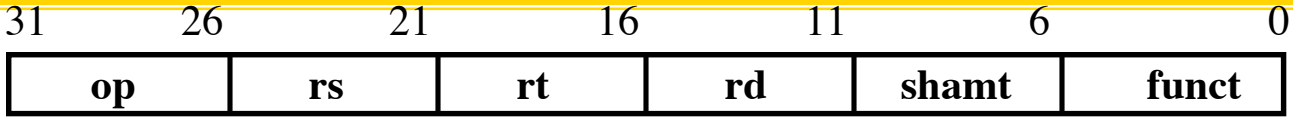
- MEM[PC] Fetch the instruction from memory
- R[rd] = R[rs] + R[rt] The actual operation
- PC = PC + 4 Calculate the next instruction's address

Instruction Fetch Unit at the Beginning of Add

- Fetch the instruction from Instruction memory: Instruction = MEM[PC]
  - same for all instructions



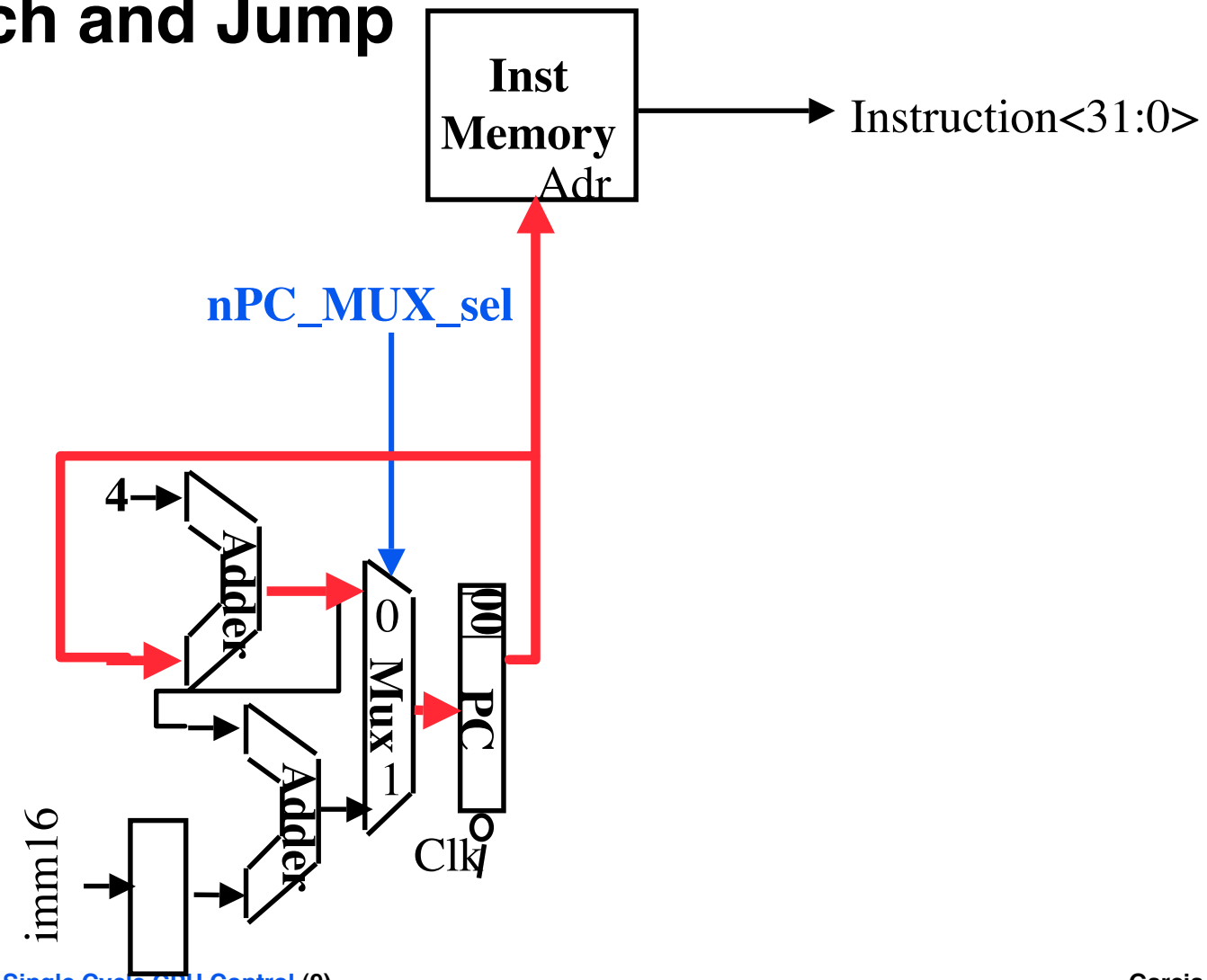
The Single Cycle Datapath during Add



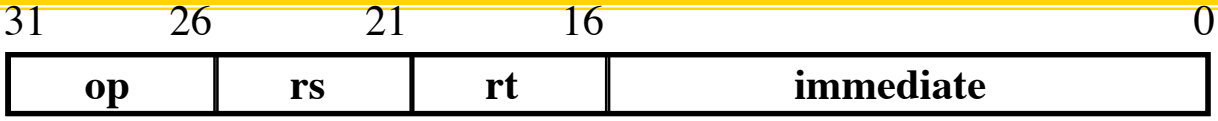
- R[rd] = R[rs] + R[rt]

Instruction Fetch Unit at the End of Add

- PC = PC + 4
  - This is the same for all instructions except: Branch and Jump

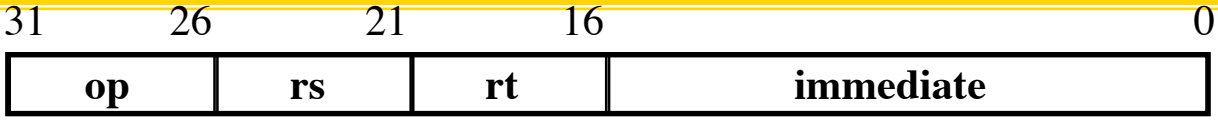


Single Cycle Datapath during Or Immediate?



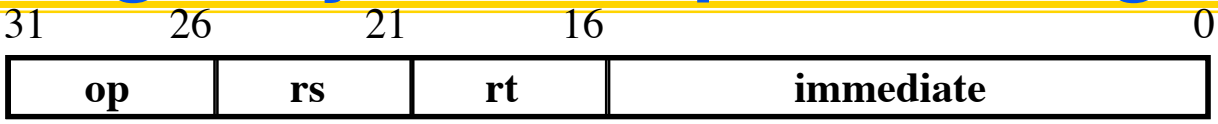
- R[rt] = R[rs] OR ZeroExt[Imm16]

Single Cycle Datapath during Or Immediate?

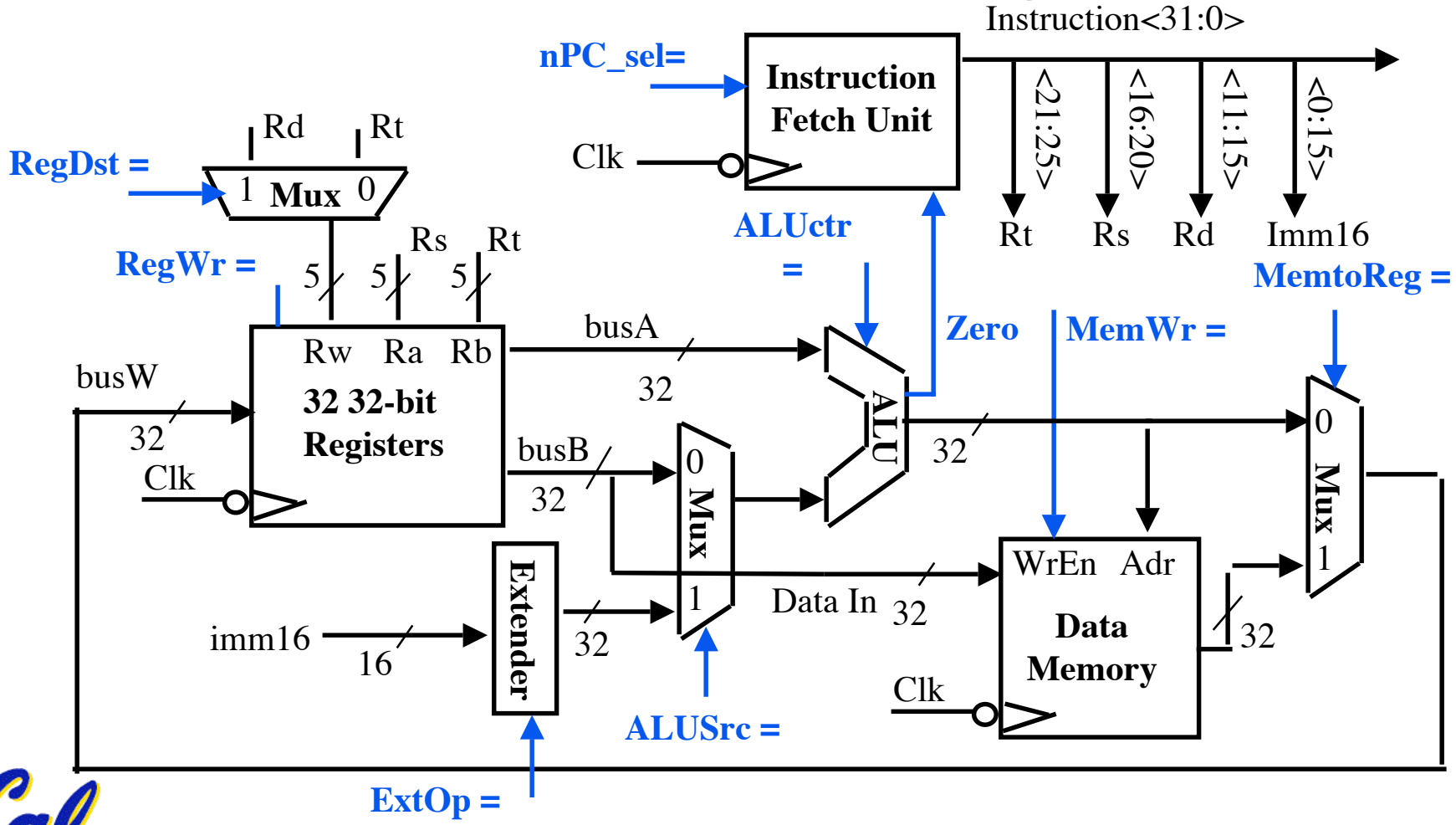


- R[rt] = R[rs] OR ZeroExt[Imm16]

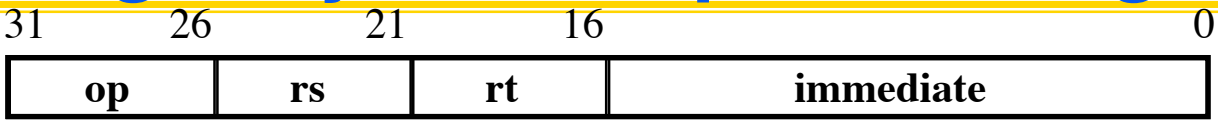
The Single Cycle Datapath during Load?



- R[rt] = Data Memory {R[rs] + SignExt[imm16]}



The Single Cycle Datapath during Load



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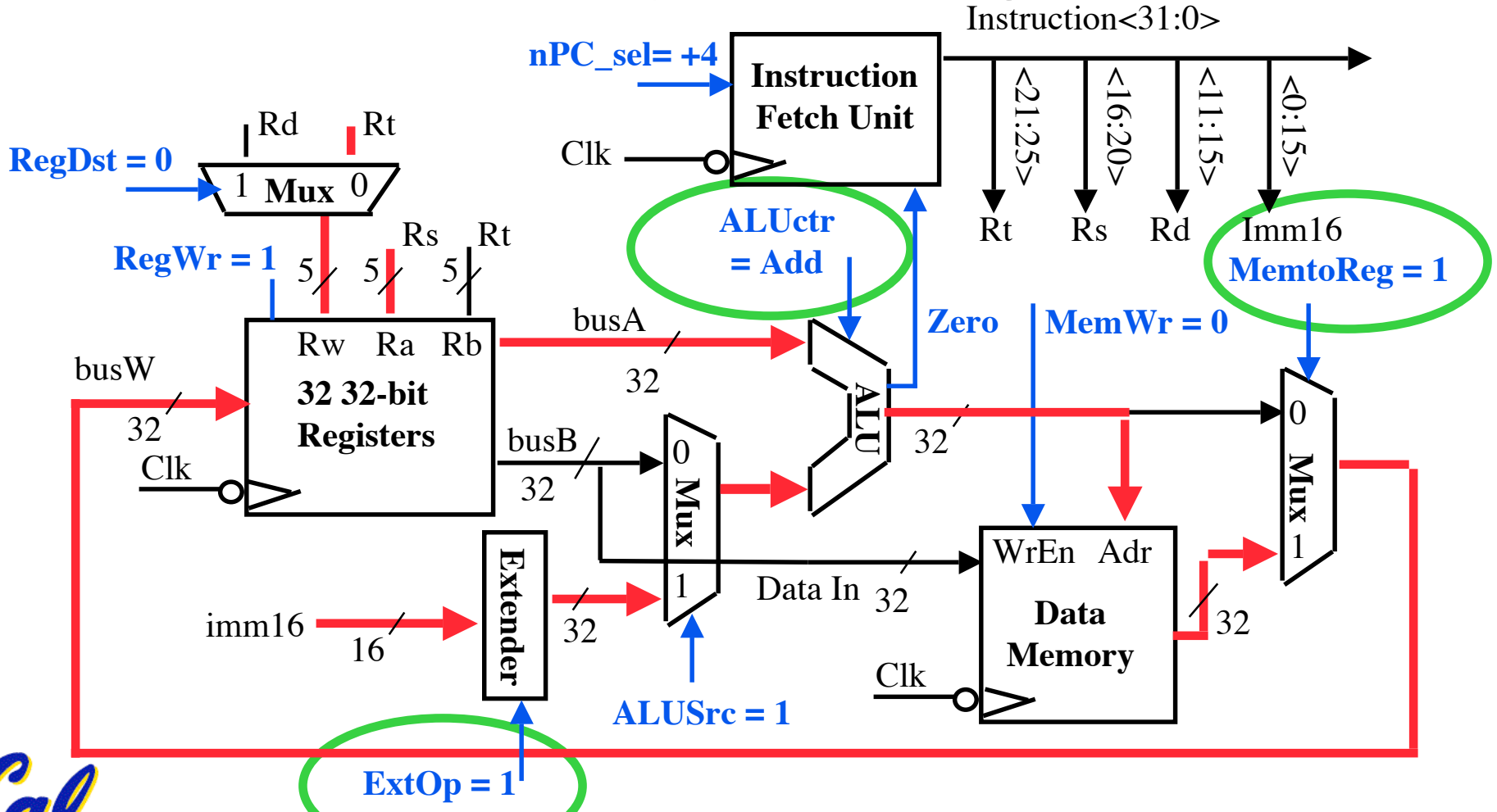
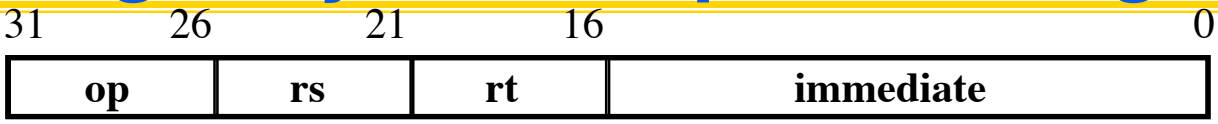
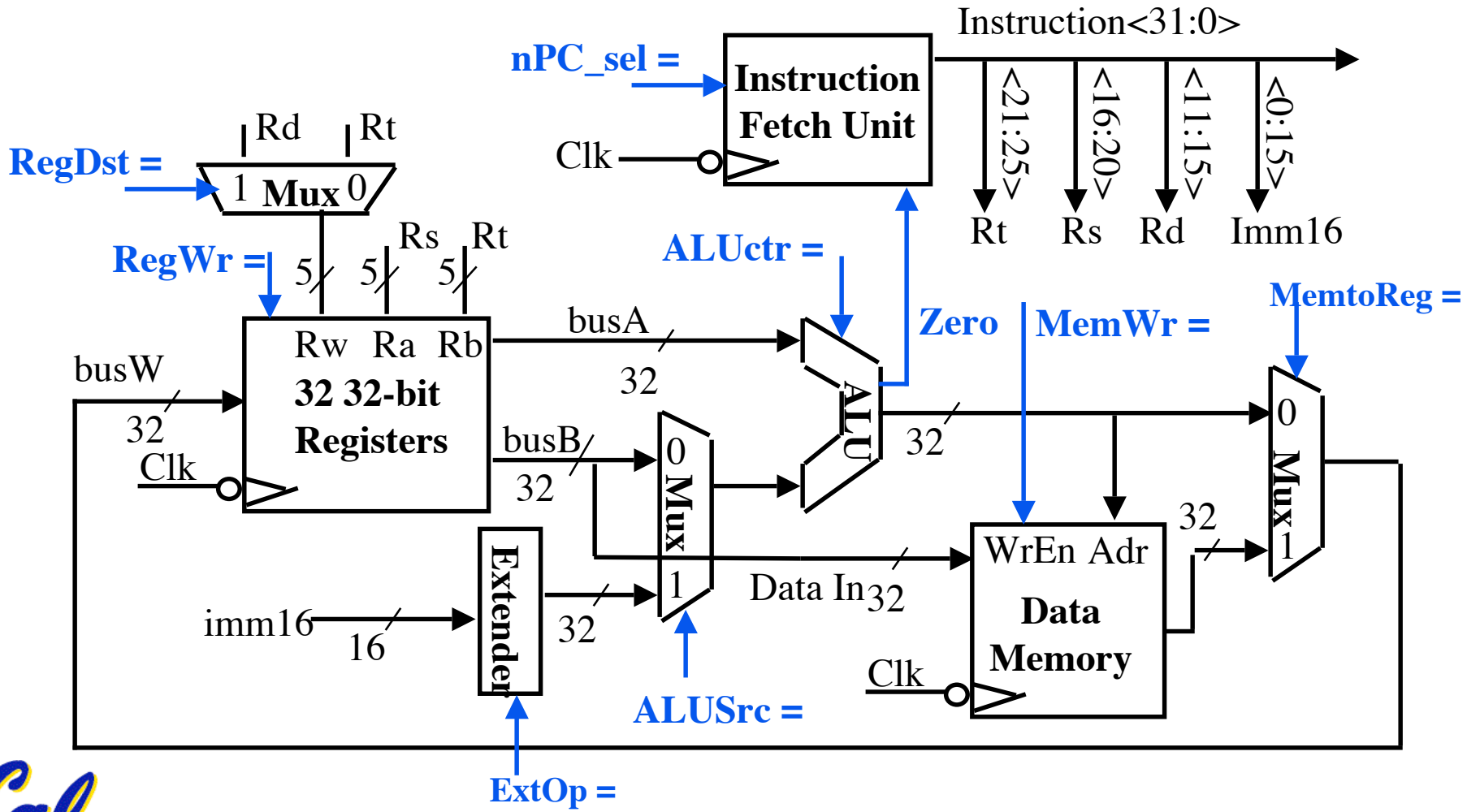


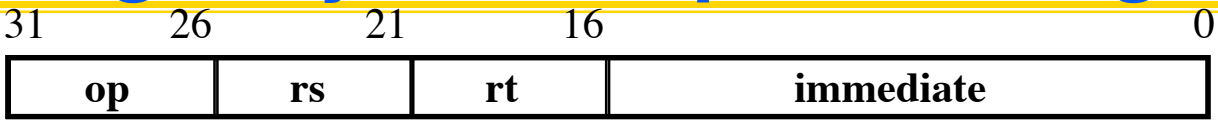
The Single Cycle Datapath during Store?



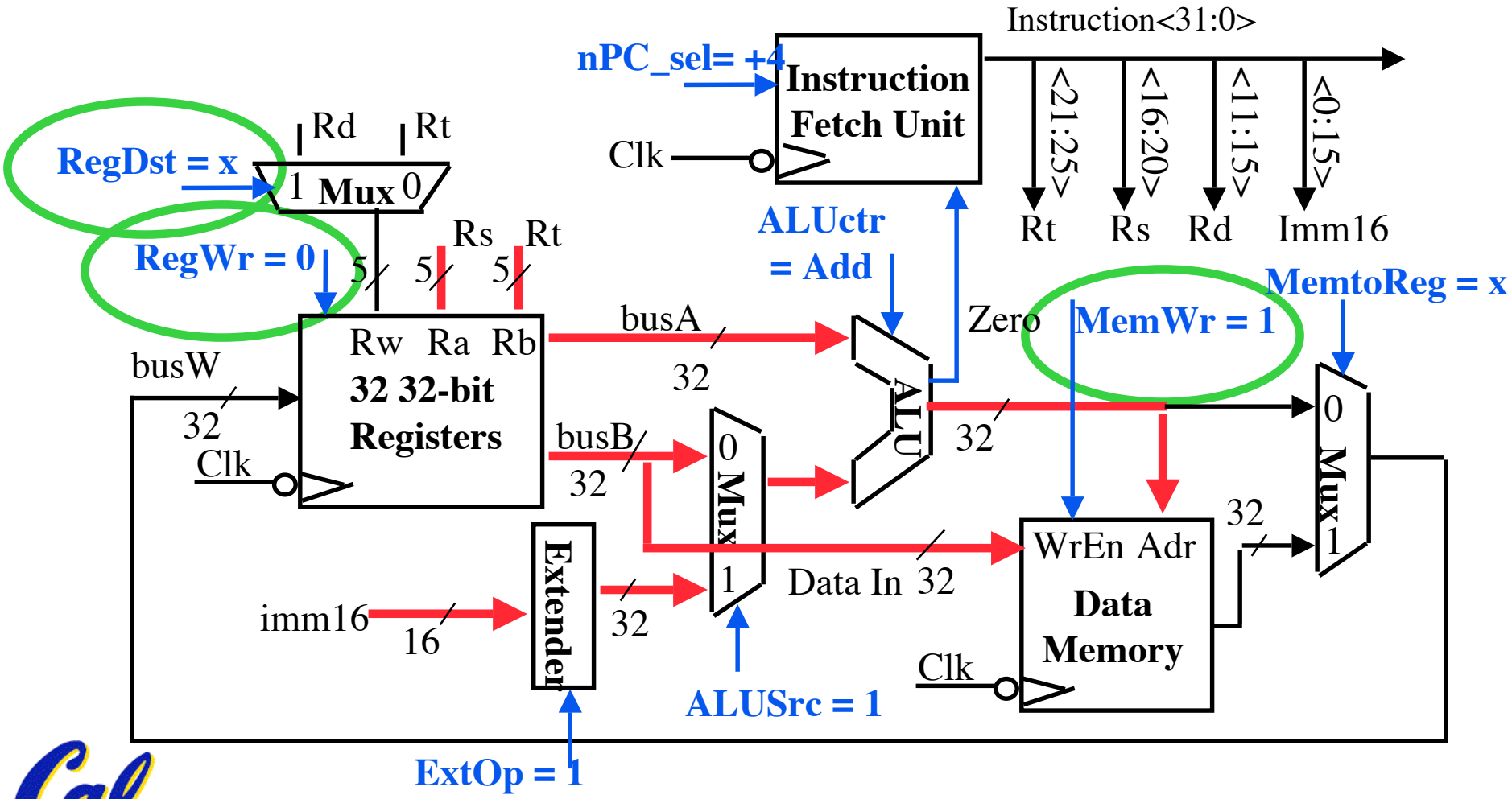
- Data Memory {R[rs] + SignExt[imm16]} = R[rt]



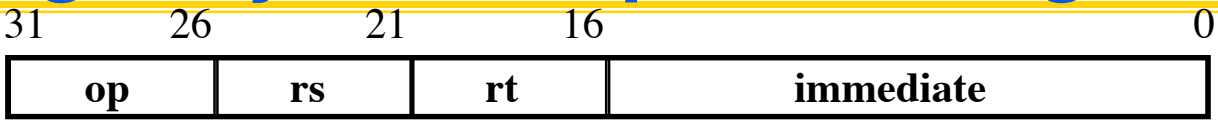
The Single Cycle Datapath during Store



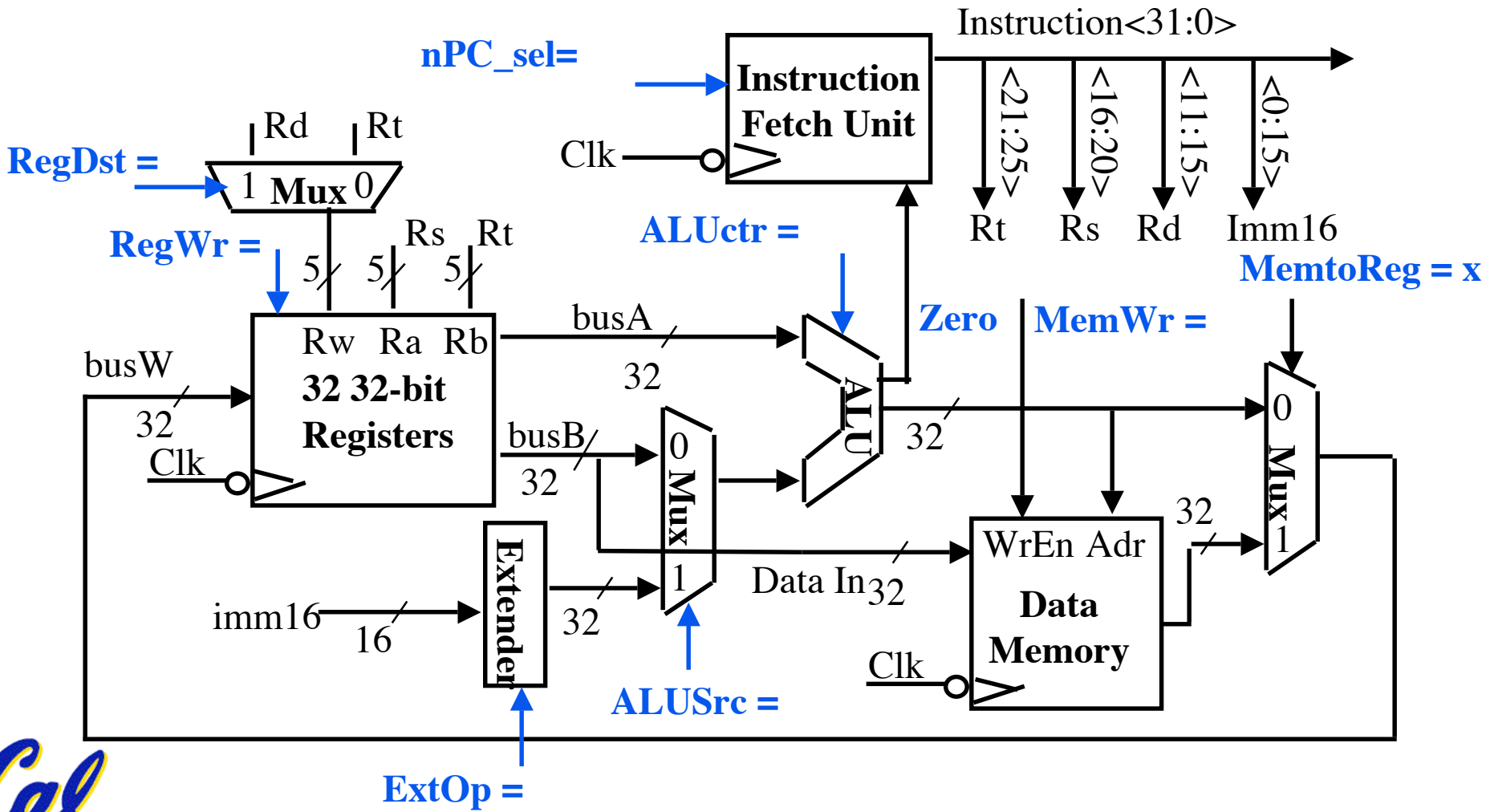
- Data Memory {R[rs] + SignExt[imm16]} = R[rt]



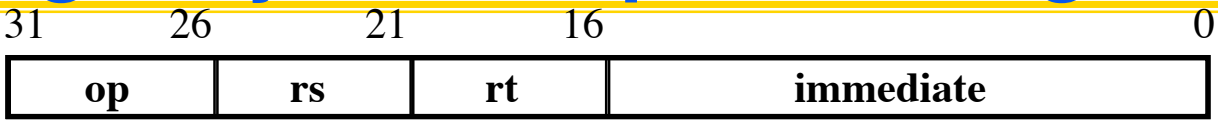
The Single Cycle Datapath during Branch?



- if (R[rs] - R[rt] == 0) then Zero = 1 ; else Zero = 0



The Single Cycle Datapath during Branch



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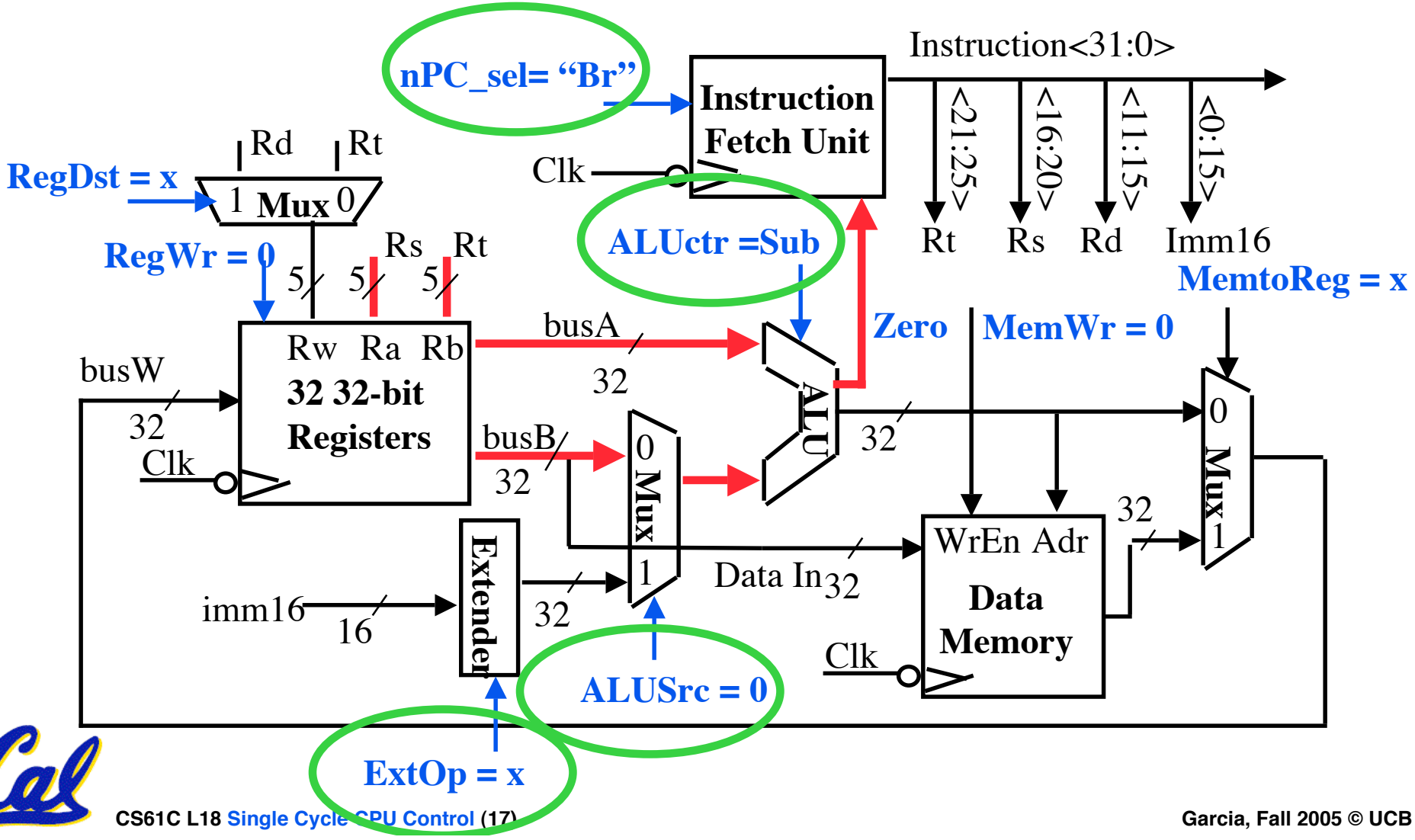
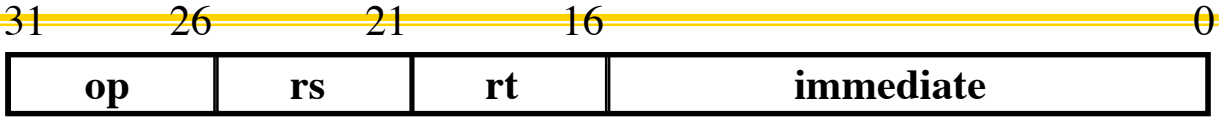
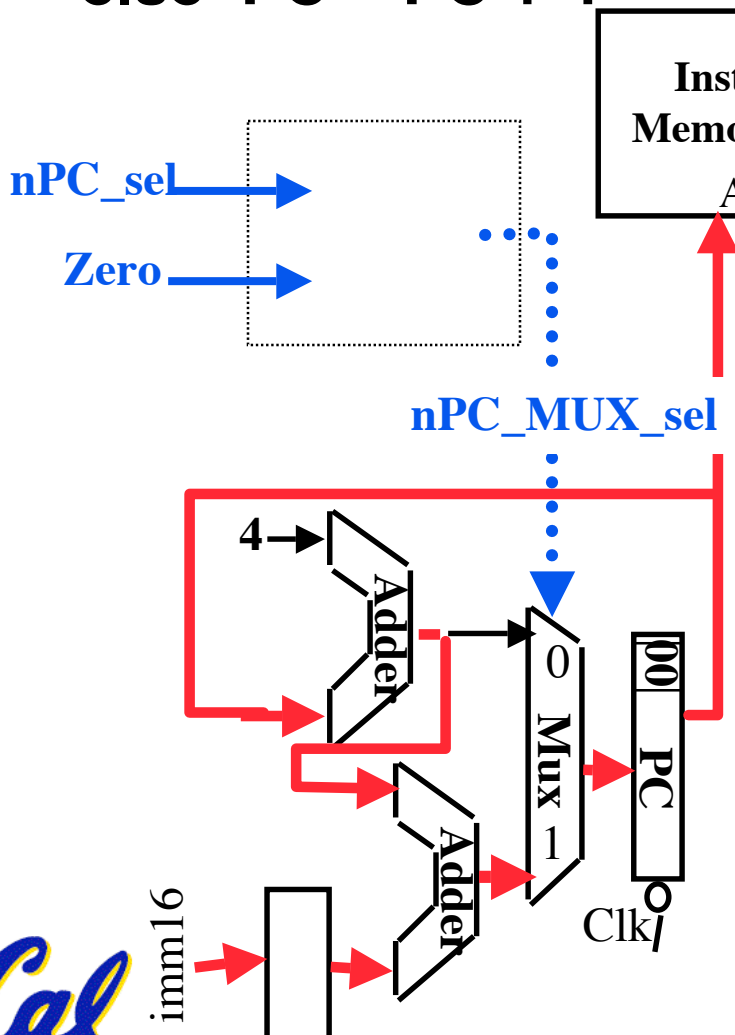


Instruction Fetch Unit at the End of Branch



- if (Zero == 1) then PC = PC + 4 + SignExt[imm16]*4 ; else PC = PC + 4



- What is encoding of nPC_sel?
  - Direct MUX select?
  - Branch / not branch
- Let's pick 2nd option

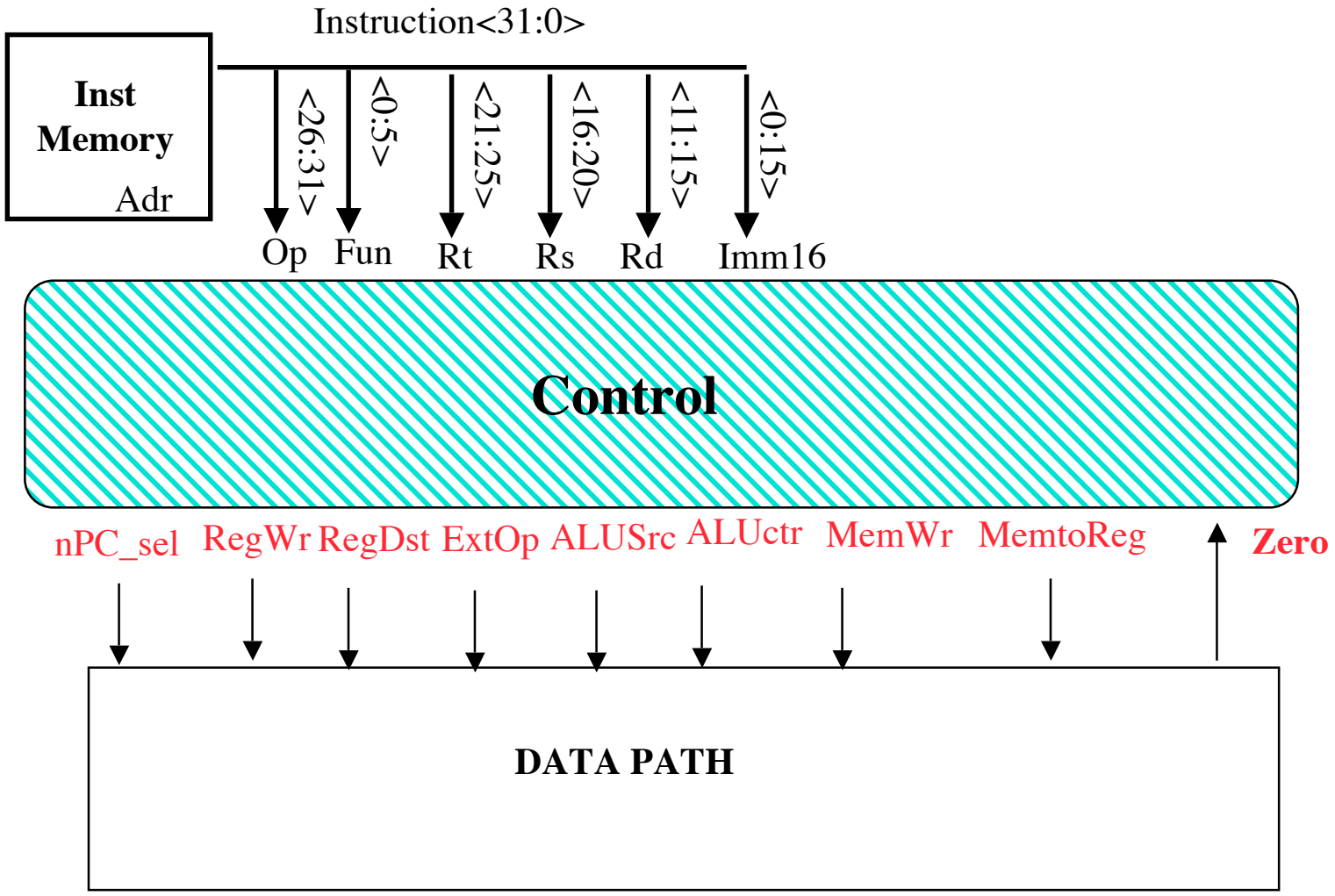
nPC_sel	zero?	MUX
0	x	0
1	0	0
1	1	1

Q: What logic gate?

Administrivia

- Any Administrivia?

Step 4: Given Datapath: RTL ⇒ Control



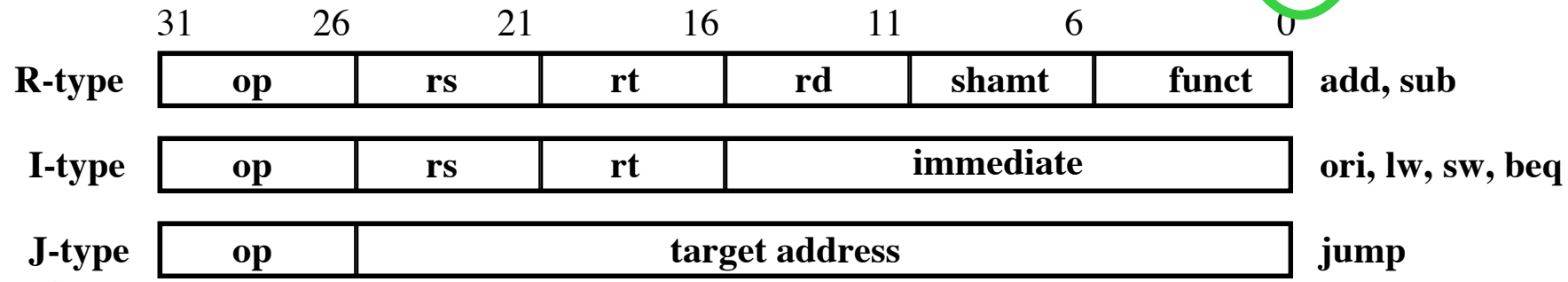
A Summary of the Control Signals (1/2)

inst	Register Transfer	
ADD	R[rd] <– R[rs] + R[rt];	PC <– PC + 4
	ALUsrc = RegB, ALUctr = "add", RegDst = rd, RegWr, nPC_sel = "+4"	
SUB	R[rd] <– R[rs] – R[rt];	PC <– PC + 4
	ALUsrc = RegB, ALUctr = "sub", RegDst = rd, RegWr, nPC_sel = "+4"	
ORi	R[rt] <– R[rs] + zero_ext(Imm16);	PC <– PC + 4
	ALUsrc = Im, Extop = "Z", ALUctr = "or", RegDst = rt, RegWr, nPC_sel ="+4"	
LOAD	R[rt] <– MEM[R[rs] + sign_ext(Imm16)];	PC <– PC + 4
	ALUsrc = Im, Extop = "Sn", ALUctr = "add", MemtoReg, RegDst = rt, RegWr,	nPC_sel = "+4"
STORE	MEM[R[rs] + sign_ext(Imm16)] <– R[rs];	PC <– PC + 4
	ALUsrc = Im, Extop = "Sn", ALUctr = "add", MemWr, nPC_sel = "+4"	
BEQ	if (R[rs] == R[rt]) then PC <– PC + sign_ext(Imm16)] || 00 else PC <– PC + 4	
	nPC_sel = "Br", ALUctr = "sub"	

A Summary of the Control Signals (2/2)

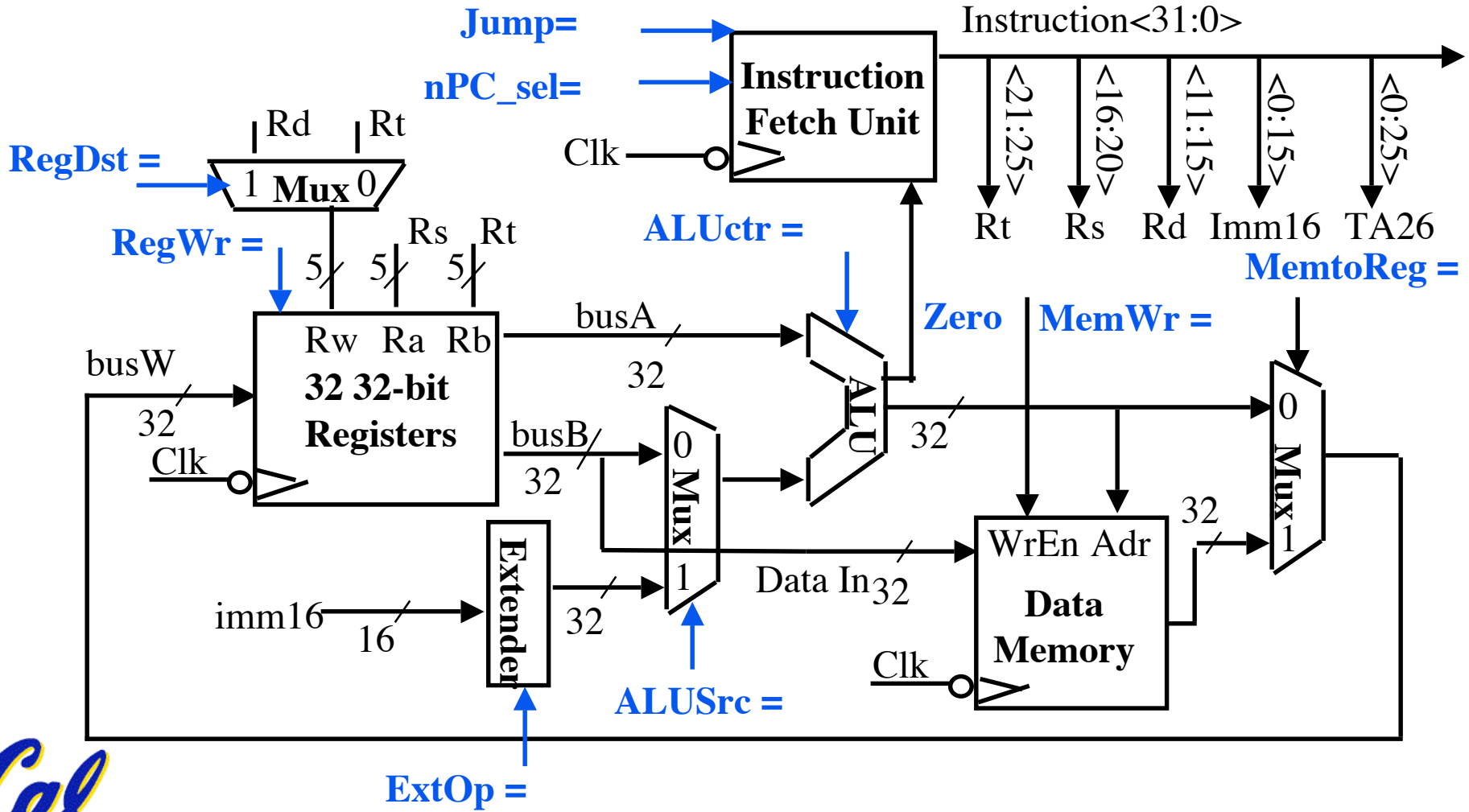
See Appendix A

	add	sub	ori	lw	sw	beq	jump
func	10 0000	10 0010	**We Don't Care :-)**				
op	00 0000	00 0000	00 1101	10 0011	10 1011	00 0100	00 0010
RegDst	1	1	0	0	x	x	x
ALUSrc	0	0	1	1	1	0	x
MemtoReg	0	0	0	1	x	x	x
RegWrite	1	1	1	1	0	0	0
MemWrite	0	0	0	0	1	0	0
nPCsel	0	0	0	0	0	1	0
Jump	0	0	0	0	0	0	1
ExtOp	x	x	0	1	1	x	x
ALUctr<2:0>	Add	Subtract	Or	Add	Add	Subtract	x



The Single Cycle Datapath during Jump

- New PC = { PC[31..28], target address, 00 }



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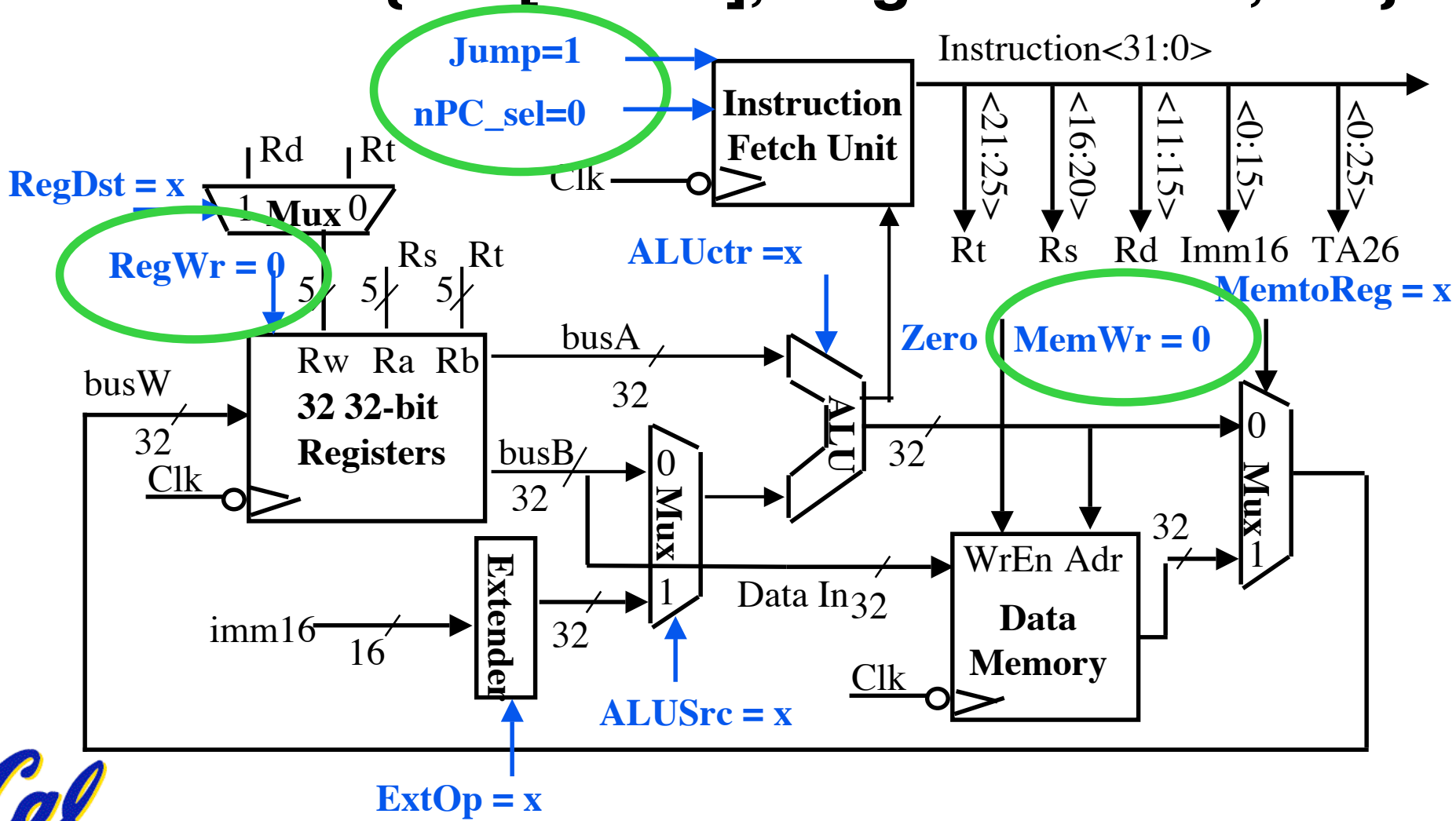

Instruction Fetch Unit at the End of Jump

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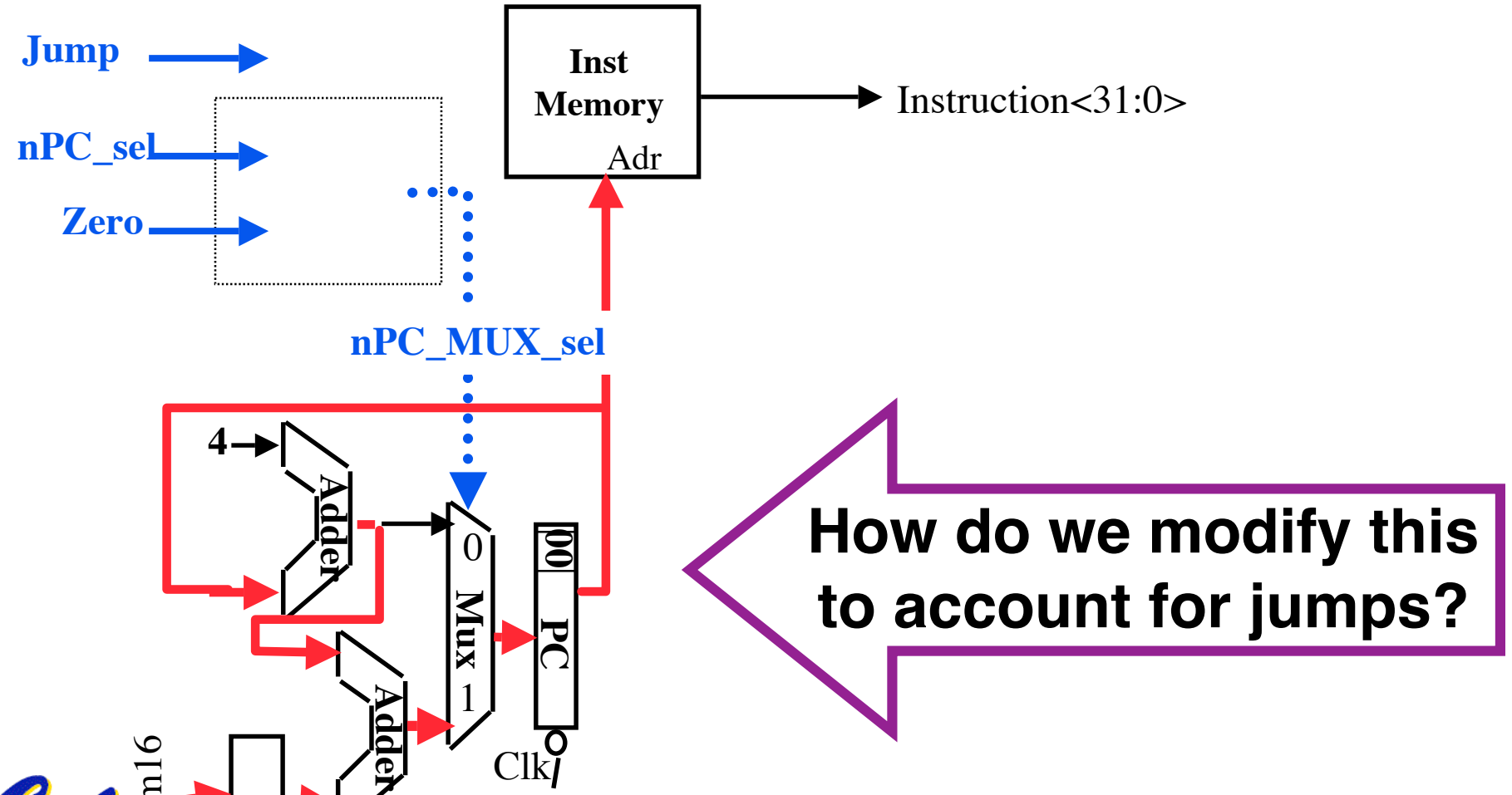
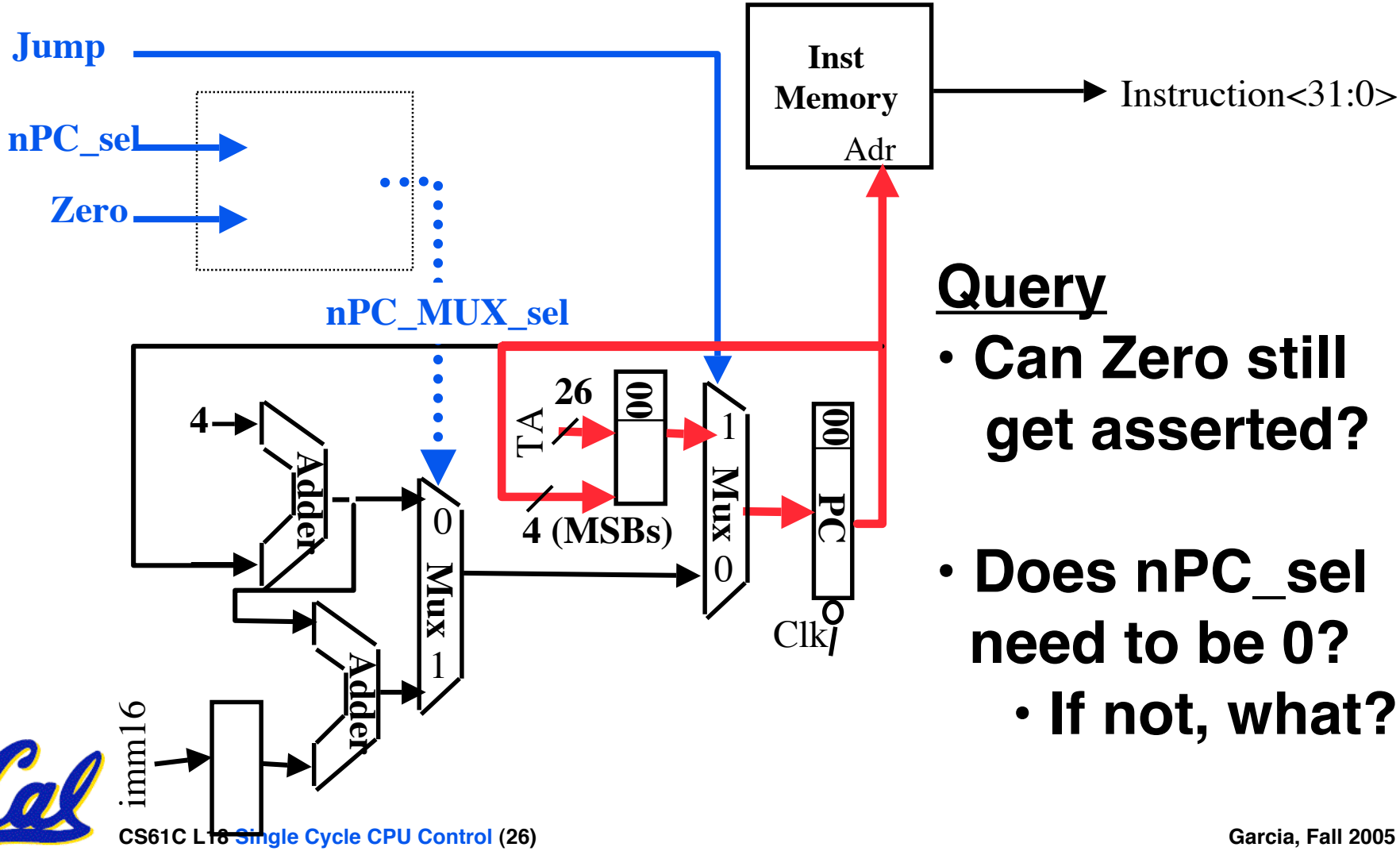


How do we modify this to account for jumps?

Instruction Fetch Unit at the End of Jump

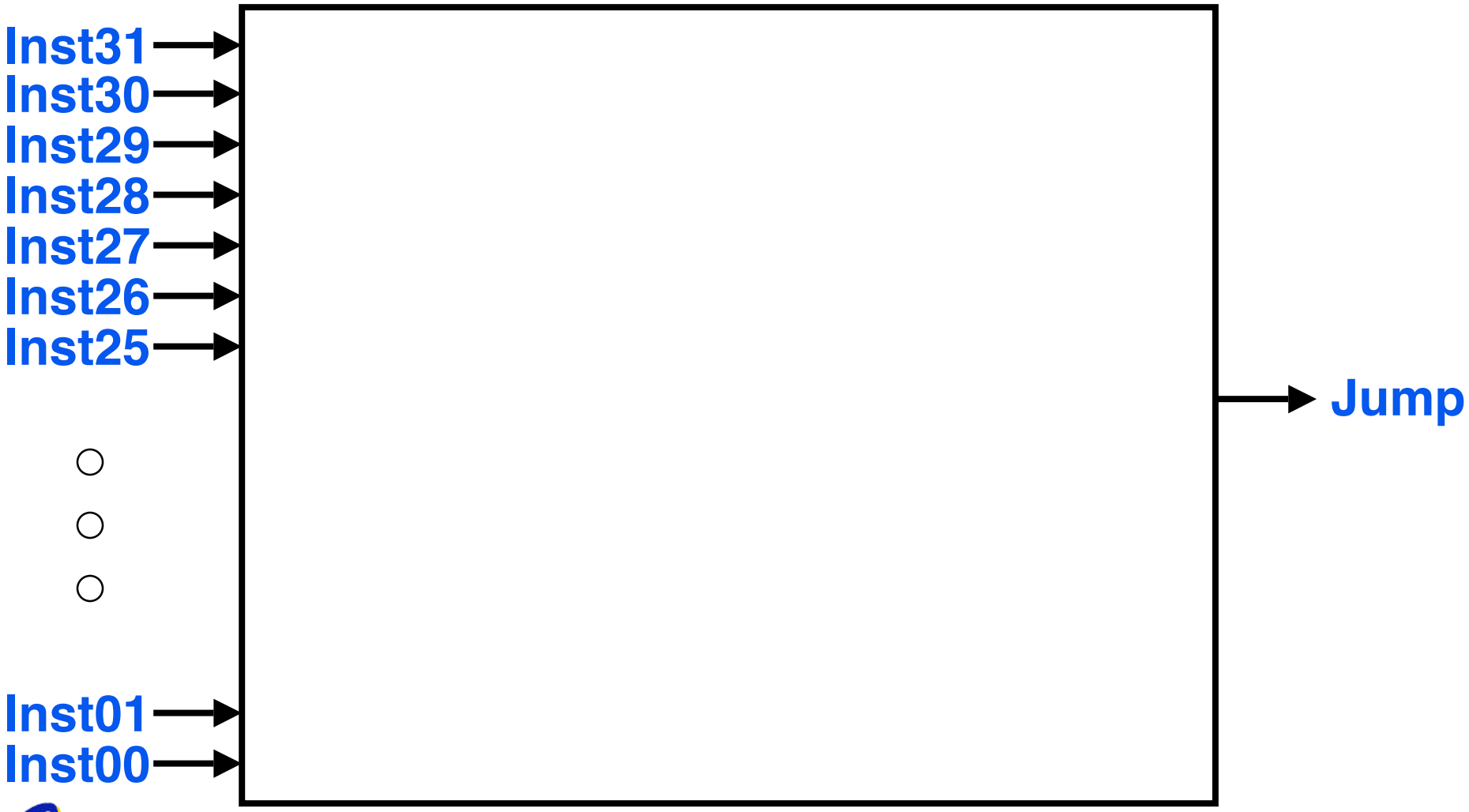
- New PC = { PC[31..28], target address, 00 }



Query

- Can Zero still get asserted?
- Does nPC_sel need to be 0?
  - If not, what?

Build CL to implement Jump on paper now



Peer Instruction

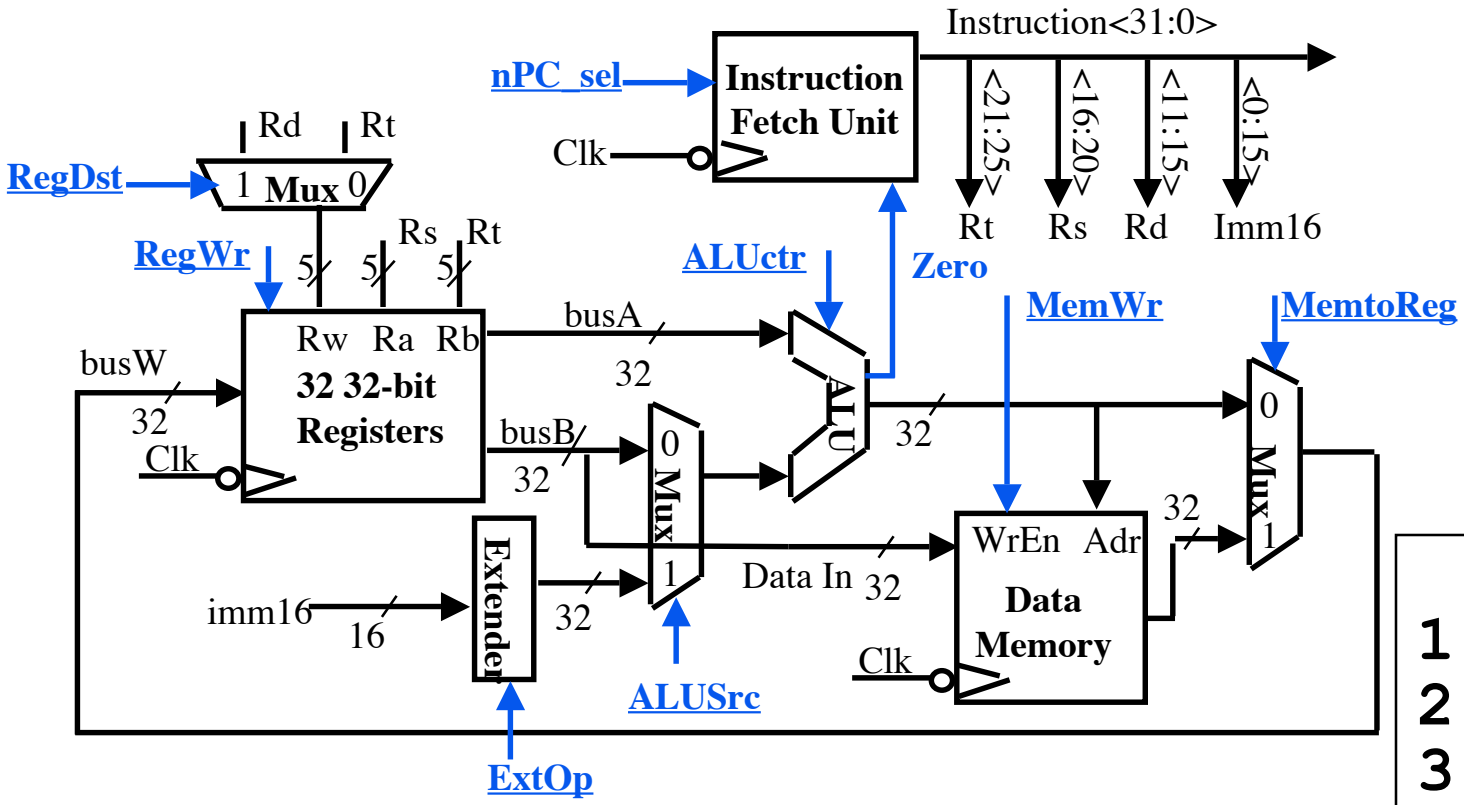
A. Our ALU is a synchronous device

B. We should use the main ALU to compute PC=PC+4

C. The ALU is inactive for memory reads or writes.

	ABC
1:	FFF
2:	FFT
3:	FTF
4:	FTT
5:	TFF
6:	TFT
7:	TTF
8:	TTT

Peer Instruction



	ABC
1:	SRF
2:	SRT
3:	SEF
4:	SET
5:	BRF
6:	BRT
7:	BEF
8:	BET

A. MemToReg='x' & ALUctr='sub'. SUB or BEQ?

B. ALUctr='add'. Which 1 signal is different for all 3 of: ADD, LW, & SW? RegDst or ExtOp?

C. "Don't Care" signals are useful because we can simplify our PLA personality matrix. F / T?

And in Conclusion… Single cycle control

- 5 steps to design a processor
  - 1. Analyze instruction set => datapath requirements
  - 2. Select set of datapath components & establish clock methodology
  - 3. Assemble datapath meeting the requirements
  - 4. Analyze implementation of each instruction to determine setting of control points that effects the register transfer.
  - 5. Assemble the control logic
- Control is the hard part
- MIPS makes that easier
  - Instructions same size
  - Source registers always in same place
  - Immediates same size, location
  - Operations always on registers/immediates