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CS61C : Machine Structures

Lecture #9 – MIPS Logical & Shift Ops, and Instruction Representation I

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There is one handout today at the front and back of the room!

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Hearing keys?! ⇒ Prof Doug Tygar revealed that you can just listen to the sound of keystrokes to find out what's typed…beware!

www.earthtimes.org/articles/show/4103.html

Review

- Functions called with `jal`, return with `jr $ra`.
- The stack is your friend: Use it to save anything you need. Just be sure to leave it the way you found it.
- Instructions we know so far
  - Arithmetic: `add, addi, sub, addu, addiu, subu`
  - Memory: `lw, sw, lb, sb, lbu`
  - Decision: `beq, bne, slt, slti, sltu, sltiu`
  - Unconditional Branches (Jumps): `j, jal, jr`
- Registers we know so far
  - All of them!
  - There are CONVENTIONS when calling procedures!

Bitwise Operations

- Up until now, we've done arithmetic (`add`, `sub`,`addi`), memory access (`lw` and `sw`), and branches and jumps.
- All of these instructions view contents of register as a single quantity (such as a signed or unsigned integer)
- New Perspective: View register as 32 raw bits rather than as a single 32-bit number
- Since registers are composed of 32 bits, we may want to access individual bits (or groups of bits) rather than the whole.
- Introduce two new classes of instructions:
  - Logical & Shift Ops

Logical Operators (1/3)

- Two basic logical operators:
  - AND: outputs 1 only if both inputs are 1
  - OR: outputs 1 if at least one input is 1
- Truth Table: standard table listing all possible combinations of inputs and resultant output for each. E.g.,

A	B	A AND B	A OR B
0	0	0	0
0	1	0	1
1	0	0	1
1	1	1	1

Logical Operators (2/3)

- Logical Instruction Syntax:

  1 2,3,4
  - where
    1) operation name
    2) register that will receive value
    3) first operand (register)
    4) second operand (register) or immediate (numerical constant)
- In general, can define them to accept > 2 inputs, but in the case of MIPS assembly, these accept exactly 2 inputs and produce 1 output
  - Again, rigid syntax, simpler hardware

Logical Operators (3/3)

- Instruction Names:
  - `and`, `or`: Both of these expect the third argument to be a register
  - `andi`, `ori`: Both of these expect the third argument to be an immediate
- MIPS Logical Operators are all bitwise, meaning that bit 0 of the output is produced by the respective bit 0's of the inputs, bit 1 by the bit 1's, etc.
  - C: Bitwise AND is `&` (e.g., `z = x & y;`)
  - C: Bitwise OR is `|` (e.g., `z = x | y;`)

Uses for Logical Operators (1/3)

- Note that anding a bit with 0 produces a 0 at the output while anding a bit with 1 produces the original bit.
- This can be used to create a mask.
  - Example:

```
     1011 0110 1010 0100 0011 1101 1001 1010
mask:0000 0000 0000 0000 0000 1111 1111 1111
```

  - The result of anding these:

```
     0000 0000 0000 0000 0000 1101 1001 1010
```

mask last 12 bits

Uses for Logical Operators (2/3)

- The second bitstring in the example is called a mask. It is used to isolate the rightmost 12 bits of the first bitstring by masking out the rest of the string (e.g. setting it to all 0s).
- Thus, the and operator can be used to set certain portions of a bitstring to 0s, while leaving the rest alone.
  - In particular, if the first bitstring in the above example were in `$t0`, then the following instruction would mask it:

```
andi   $t0,$t0,0xFFF
```

Uses for Logical Operators (3/3)

- Similarly, note that oring a bit with 1 produces a 1 at the output while oring a bit with 0 produces the original bit.
- This can be used to force certain bits of a string to 1s.
  - For example, if `$t0` contains `0x12345678`, then after this instruction:

```
ori  $t0, $t0, 0xFFFF
```

  - … `$t0` contains `0x1234FFFF` (e.g. the high-order 16 bits are untouched, while the low-order 16 bits are forced to 1s).

Shift Instructions (1/4)

- Move (shift) all the bits in a word to the left or right by a number of bits.
  - Example: shift right by 8 bits

```
0001 0010 0011 0100 0101 0110 0111 1000

0000 0000 0001 0010 0011 0100 0101 0110
```

  - Example: shift left by 8 bits

```
0001 0010 0011 0100 0101 0110 0111 1000

0011 0100 0101 0110 0111 1000 0000 0000
```

Shift Instructions (2/4)

- Shift Instruction Syntax:
  - 1 2,3,4
  - where
    - 1) operation name
    - 2) register that will receive value
    - 3) first operand (register)
    - 4) shift amount (constant < 32)
- MIPS shift instructions:
  1. `sll` (shift left logical): shifts left and fills emptied bits with 0s
  2. `srl` (shift right logical): shifts right and fills emptied bits with 0s
  3. `sra` (shift right arithmetic): shifts right and fills emptied bits by sign extending

Shift Instructions (3/4)

- Example: shift right arith by 8 bits

```
0001 0010 0011 0100 0101 0110 0111 1000

0000 0000 0001 0010 0011 0100 0101 0110
```

- Example: shift right arith by 8 bits

```
1001 0010 0011 0100 0101 0110 0111 1000

1111 1111 1001 0010 0011 0100 0101 0110
```

Shift Instructions (4/4)

- Since shifting may be faster than multiplication, a good compiler usually notices when C code multiplies by a power of 2 and compiles it to a shift instruction:
  - `a *= 8;` (in C)
  - would compile to:
  - `sll    $s0,$s0,3` (in MIPS)
- Likewise, shift right to divide by powers of 2
  - remember to use `sra`

Peer Instruction

```
r: ...     # R/W $s0,$v0,$t0,$a0,$sp,$ra,mem
   ...     ### PUSH REGISTER(S) TO STACK?
   jal e   # Call e
   ...     # R/W $s0,$v0,$t0,$a0,$sp,$ra,mem
   jr $ra  # Return to caller of r

e: ...     # R/W $s0,$v0,$t0,$a0,$sp,$ra,mem
   jr $ra  # Return to r
```

What does `r` have to push on the stack before "`jal e`"?

```
1:  1 of ($s0,$sp,$v0,$t0,$a0,$ra)
2:  2 of ($s0,$sp,$v0,$t0,$a0,$ra)
3:  3 of ($s0,$sp,$v0,$t0,$a0,$ra)
4:  4 of ($s0,$sp,$v0,$t0,$a0,$ra)
5:  5 of ($s0,$sp,$v0,$t0,$a0,$ra)
6:  6 of ($s0,$sp,$v0,$t0,$a0,$ra)
7:  0 of ($s0,$sp,$v0,$t0,$a0,$ra)
```

Overview – Instruction Representation

- Big idea: stored program
  - consequences of stored program
- Instructions as numbers
- Instruction encoding
- MIPS instruction format for Add instructions
- MIPS instruction format for Immediate, Data transfer instructions

Big Idea: Stored-Program Concept

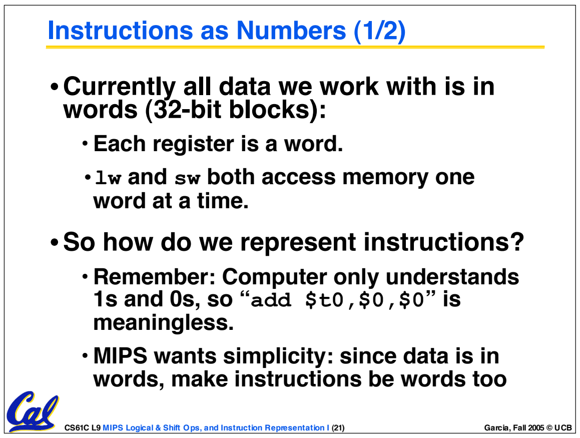
- Computers built on 2 key principles:
  1) Instructions are represented as numbers.
  2) Therefore, entire programs can be stored in memory to be read or written just like numbers (data).
- Simplifies SW/HW of computer systems:
  - Memory technology for data also used for programs

Consequence #1: Everything Addressed

- Since all instructions and data are stored in memory as numbers, everything has a memory address: instructions, data words
  - both branches and jumps use these
- C pointers are just memory addresses: they can point to anything in memory
  - Unconstrained use of addresses can lead to nasty bugs; up to you in C; limits in Java
- One register keeps address of instruction being executed: "Program Counter" (PC)
  - Basically a pointer to memory: Intel calls it Instruction Address Pointer, a better name

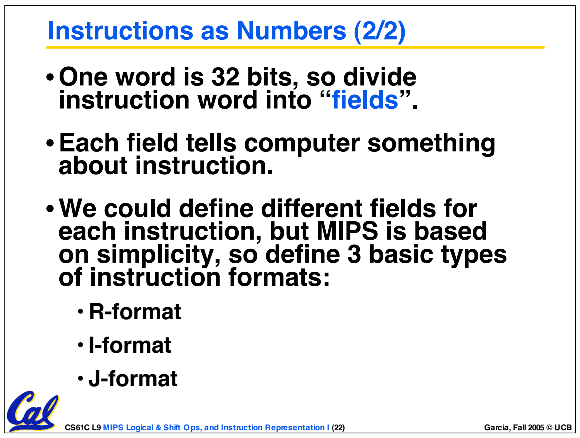
Consequence #2: Binary Compatibility

- Programs are distributed in binary form
  - Programs bound to specific instruction set
  - Different version for Macintoshes and PCs
- New machines want to run old programs ("binaries") as well as programs compiled to new instructions
- Leads to instruction set evolving over time
- Selection of Intel 8086 in 1981 for 1st IBM PC is major reason latest PCs still use 80x86 instruction set (Pentium 4); could still run program from 1981 PC today



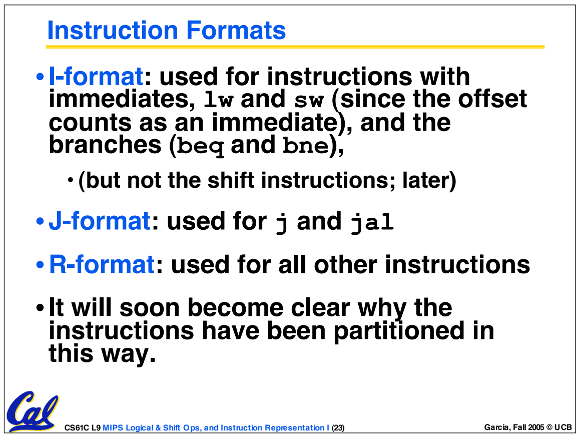
Instructions as Numbers (1/2)

- Currently all data we work with is in words (32-bit blocks):
  - Each register is a word.
  - `lw` and `sw` both access memory one word at a time.
- So how do we represent instructions?
  - Remember: Computer only understands 1s and 0s, so "`add $t0,$0,$0`" is meaningless.
  - MIPS wants simplicity: since data is in words, make instructions be words too



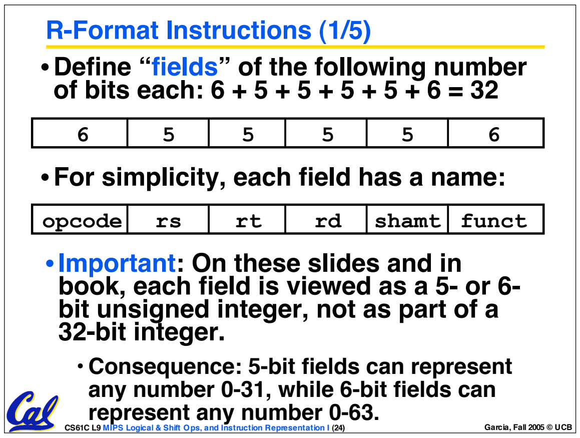
Instructions as Numbers (2/2)

- One word is 32 bits, so divide instruction word into "fields".
- Each field tells computer something about instruction.
- We could define different fields for each instruction, but MIPS is based on simplicity, so define 3 basic types of instruction formats:
  - R-format
  - I-format
  - J-format



Instruction Formats

- I-format: used for instructions with immediates, `lw` and `sw` (since the offset counts as an immediate), and the branches (`beq` and `bne`),
  - (but not the shift instructions; later)
- J-format: used for `j` and `jal`
- R-format: used for all other instructions
- It will soon become clear why the instructions have been partitioned in this way.



R-Format Instructions (1/5)

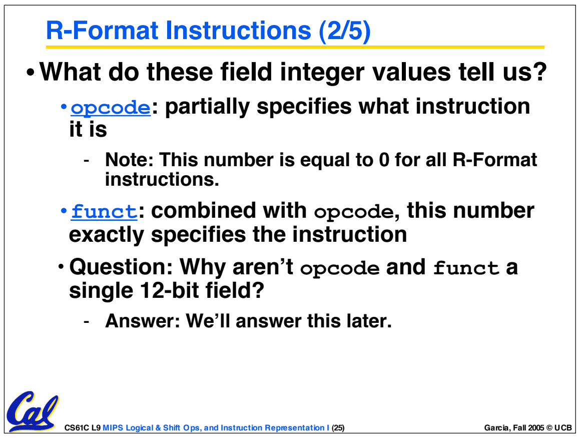
- Define "fields" of the following number of bits each: 6 + 5 + 5 + 5 + 5 + 6 = 32

6	5	5	5	5	6

- For simplicity, each field has a name:

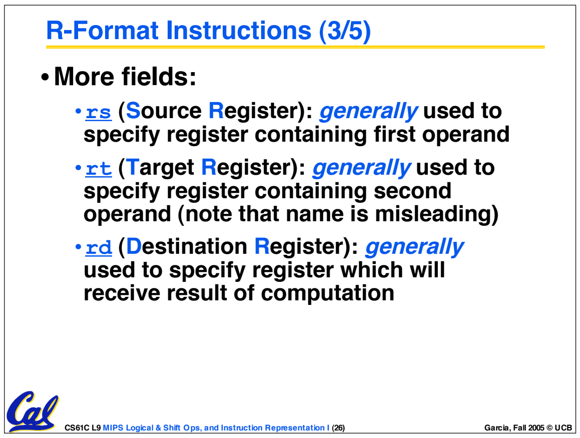
`opcode`	`rs`	`rt`	`rd`	`shamt`	`funct`

- Important: On these slides and in book, each field is viewed as a 5- or 6-bit unsigned integer, not as part of a 32-bit integer.
  - Consequence: 5-bit fields can represent any number 0-31, while 6-bit fields can represent any number 0-63.



R-Format Instructions (2/5)

- What do these field integer values tell us?
  - `opcode`: partially specifies what instruction it is
    - Note: This number is equal to 0 for all R-Format instructions.
  - `funct`: combined with `opcode`, this number exactly specifies the instruction
  - Question: Why aren't `opcode` and `funct` a single 12-bit field?
    - Answer: We'll answer this later.



R-Format Instructions (3/5)

- More fields:
  - `rs` (Source Register): *generally* used to specify register containing first operand
  - `rt` (Target Register): *generally* used to specify register containing second operand (note that name is misleading)
  - `rd` (Destination Register): *generally* used to specify register which will receive result of computation

R-Format Instructions (4/5)

- Notes about register fields:
  - Each register field is exactly 5 bits, which means that it can specify any unsigned integer in the range 0-31. Each of these fields specifies one of the 32 registers by number.
  - The word "generally" was used because there are exceptions that we'll see later. E.g.,
    - `mult` and `div` have nothing important in the `rd` field since the dest registers are `hi` and `lo`
    - `mfhi` and `mflo` have nothing important in the `rs` and `rt` fields since the source is determined by the instruction (p. 264 P&H)

R-Format Instructions (5/5)

- Final field:
  - `shamt`: This field contains the amount a shift instruction will shift by. Shifting a 32-bit word by more than 31 is useless, so this field is only 5 bits (so it can represent the numbers 0-31).
  - This field is set to 0 in all but the shift instructions.
- For a detailed description of field usage for each instruction, see green insert in COD 3/e
  - (You can bring with you to all exams)

R-Format Example (1/2)

- MIPS Instruction:

```
add   $8,$9,$10
```

`opcode` = 0 (look up in table in book)

`funct` = 32 (look up in table in book)

`rd` = 8 (destination)

`rs` = 9 (first *operand*)

`rt` = 10 (second *operand*)

`shamt` = 0 (not a shift)

R-Format Example (2/2)

- MIPS Instruction:

```
add   $8,$9,$10
```

Decimal number per field representation:

0	9	10	8	0	32

Binary number per field representation:

000000	01001	01010	01000	00000	100000

hex

hex representation: 012A 4020$_{hex}$

decimal representation: 19,546,144$_{ten}$

- Called a Machine Language Instruction

I-Format Instructions (1/4)

- What about instructions with immediates?
  - 5-bit field only represents numbers up to the value 31: immediates may be much larger than this
  - Ideally, MIPS would have only one instruction format (for simplicity): unfortunately, we need to compromise
- Define new instruction format that is partially consistent with R-format:
  - First notice that, if instruction has immediate, then it uses at most 2 registers.

I-Format Instructions (2/4)

- Define "fields" of the following number of bits each: 6 + 5 + 5 + 16 = 32 bits

6	5	5	16

- Again, each field has a name:

`opcode`	`rs`	`rt`	`immediate`

- Key Concept: Only one field is inconsistent with R-format. Most importantly, `opcode` is still in same location.

I-Format Instructions (3/4)

- What do these fields mean?
  - `opcode`: same as before except that, since there's no `funct` field, `opcode` uniquely specifies an instruction in I-format
  - This also answers question of why R-format has two 6-bit fields to identify instruction instead of a single 12-bit field: in order to be consistent with other formats.
  - `rs`: specifies the *only* register operand (if there is one)
  - `rt`: specifies register which will receive result of computation (this is why it's called the *target* register "rt")

I-Format Instructions (4/4)

- The Immediate Field:
  - `addi`, `slti`, `sltiu`, the immediate is sign-extended to 32 bits. Thus, it's treated as a signed integer.
  - 16 bits ➔ can be used to represent immediate up to 2^{16} different values
  - This is large enough to handle the offset in a typical `lw` or `sw`, plus a vast majority of values that will be used in the `slti` instruction.
  - We'll see what to do when the number is too big in our next lecture…

I-Format Example (1/2)

- MIPS Instruction:

```
addi   $21,$22,-50
```

`opcode` = 8 (look up in table in book)

`rs` = 22 (register containing operand)

`rt` = 21 (target register)

`immediate` = -50 (by default, this is decimal)

I-Format Example (2/2)

- MIPS Instruction:

```
addi   $21,$22,-50
```

Decimal/field representation:

8	22	21	-50

Binary/field representation:

001000	10110	10101	1111111111001110

hexadecimal representation: 22D5 FFCE_{hex}

decimal representation: $584{,}449{,}998_{ten}$

Peer Instruction

Which instruction has same representation as 35_{ten}?

1. add $0, $0, $0	opcode	rs	rt	rd	shamt	funct
2. subu $s0,$s0,$s0	opcode	rs	rt	rd	shamt	funct
3. lw $0, 0($0)	opcode	rs	rt	offset		
4. addi $0, $0, 35	opcode	rs	rt	immediate		
5. subu $0, $0, $0	opcode	rs	rt	rd	shamt	funct

6. Trick question!
 Instructions are not numbers

Registers numbers and names:
 0: $0, .. 8: $t0, 9:$t1, ..15: $t7, 16: $s0, 17: $s1, .. 23: $s7

Opcodes and function fields (if necessary)

`add`: opcode = 0, funct = 32

`subu`: opcode = 0, funct = 35

`addi`: opcode = 8

`lw`: opcode = 35

In conclusion…

- Logical and Shift Instructions
  - Operate on individual bits (arithmetic operate on entire word)
  - Use to isolate fields, either by masking or by shifting back & forth
  - Use <u>shift left logical</u>, `sll`, for multiplication by powers of 2
  - Use <u>shift right arithmetic</u>, `sra`, for division by powers of 2
- Simplifying MIPS: Define instructions to be same size as data word (one word) so that they can use the same memory (compiler can use `lw` and `sw`).
- Computer actually stores programs as a series of these 32-bit numbers.
- MIPS Machine Language Instruction: 32 bits representing a single instruction

R	opcode	rs	rt	rd	shamt	funct
I	opcode	rs	rt	immediate		
J	opcode	target address				