1 More Running Time

Give the worst case and best case running time in \( \Theta(\cdot) \) notation in terms of \( M \) and \( N \).

(a) Assume that \( \text{slam()} \in \Theta(1) \) and returns a boolean.

```java
public void comeon() {
    int j = 0;
    for (int i = 0; i < N; i += 1) {
        for (; j < M; j += 1) {
            if (slam(i, j))
                break;
        }
    }
    for (int k = 0; k < 1000 * N; k += 1) {
        System.out.println("space jam");
    }
}
```

2 Recursive Running Time

For the following recursive functions, give the worst case and best case running time in \( \Theta(\cdot) \) notation.

(a) Give the running time in terms of \( N \).

```java
public void andslam(int N) {
    if (N > 0) {
        for (int i = 0; i < N; i += 1) {
            System.out.println("bigballer.jpg");
        }
        andslam(N / 2);
    }
}
```
(b) Give the running time for `andwelcome(arr, 0, N)` where `N` is the length of the input array `arr`.

```java
public static void andwelcome(int[] arr, int low, int high) {
    for (int i = low; i < high; i += 1) {
        System.out.print("loyal ");
    }
    System.out.println("]");
    if (high - low > 0) {
        double coin = Math.random();
        if (coin > 0.5) {
            andwelcome(arr, low, low + (high - low) / 2);
        } else {
            andwelcome(arr, low, low + (high - low) / 2);
            andwelcome(arr, low + (high - low) / 2, high);
        }
    }
}
```

(c) Give the running time in terms of `N`.

```java
public int tothe(int N) {
    if (N <= 1) {
        return N;
    }
    return tothe(N - 1) + tothe(N - 1) + tothe(N - 1);
}
```

(d) Give the running time in terms of `N`.

```java
public static void spacejam(int N) {
    if (N == 1) {
        return;
    }
    for (int i = 0; i < N; i += 1) {
        spacejam(N-1);
    }
}
```
3 Hey you watchu gon do?

For each example below, there are two algorithms solving the same problem. Given the asymptotic runtimes for each, is one of the algorithms guaranteed to be faster? If so, which? And if neither is always faster, explain why. Assume the algorithms have very large input (so \( N \) is very large).

(a) Algorithm 1: \( \Theta(N) \), Algorithm 2: \( \Theta(N^2) \)
(b) Algorithm 1: \( \Omega(N) \), Algorithm 2: \( \Omega(N^2) \)
(c) Algorithm 1: \( O(N) \), Algorithm 2: \( O(N^2) \)
(d) Algorithm 1: \( \Theta(N^2) \), Algorithm 2: \( O(\log N) \)
(e) Algorithm 1: \( O(N \log N) \), Algorithm 2: \( \Omega(N \log N) \)

Why did we need to assume that \( N \) was large?

4 Big Ballin’ Bounds

1. Prove the following bounds by finding some constant \( M > 0 \) and input \( N > 0 \) for \( M \in \mathbb{R}, N \in \mathbb{N} \) such that \( f \) and \( g \) satisfy the relationship.

   (a) \( f \in O(g) \) for \( f = 2n, g = n^2 \)
   (b) \( f \in \Omega(g) \) for \( f = 0.1n, g = 40 \)
   (c) \( f \in \Theta(g) \) for \( f = \log(n), g = \log(n^a) \), for \( a > 0 \).

2. Answer the following claims with true or false. If false, provide a counterexample.

   (a) If \( f(n) \in O(g(n)) \), then \( 500f(n) \in O(g(n)) \).
   (b) If \( f(n) \in \Theta(g(n)) \), then \( 2^f(n) \in \Theta(2^g(n)) \).