# CS61B Lecture #14: Integers

#### Ann

- .

# Integer Types and Literals

Announcement:		Туре	Bits	Signed?	Literals			
<ul> <li>Project #0 due Tuesday night.</li> </ul>		byte		Yes	Cast from int: (byte) 3			
<ul> <li>Programming contest SATURDAY! You can still sig</li> </ul>	n un at	short	16	Yes	None. Cast from int: (short) 4096			
https://inst.eecs.berkeley.edu/~ctest/conte	•				'a' // (char) 97 '\n' // newline ((char) 10)			
• Test #1 will be Tuesday, 6 October 2015, 8-10PM	C C	char	16	No	$'\t'$ // tab ((char) 8)			
					'\\' // backslash			
<ul> <li>Test #2 will be Tuesday, 10 November 2015, 7-9P</li> </ul>	<i>M</i> .				'A', '\101', '\u0041' // == (char) 65			
<b>Today:</b> Integer Types; Readings: A Java Referen First Java, Chapter 10.	ce, §6.3-4. Head	int	32	Yes	123 0100 // Octal for 64 0x3f, 0xffffffff // Hexadecimal 63, -1 (!)			
Readings for Upcoming Topics: Data Structures (	Into Java), Chap-	long	64	Yes	123L, 01000L, 0x3fL 1234567891011L			
ter 1.		<ul> <li>Negative numerals are just negated (positive) literals.</li> </ul>						
		• "N bits" means that there are $2^N$ integers in the domain of the type:						
				- If signed, range of values is $-2^{N-1} \dots 2^{N-1} - 1$ .				
		- If ur	nsigne	d, only no	on-negative numbers, and range is $02^N - 1.$			
Last modified: Mon Sep 28 12:10:19 2015	CS61B: Lecture #14 1	Last modified: Mon	Sep 28 12	:10:19 2015	C561B: Lecture #14 2			
Modular Arithmetic		Modular Arithmetic: Examples						
• Problem: How do we handle overflow, such as occur	s in 10000*10000*10000?	• (byte)	(64*	8) <b>yields</b>	<b>0</b> , since $512 - 0 = 2 \times 2^8$ .			
<ul> <li>Some languages throw an exception (Ada), some g sults (C, C++)</li> </ul>	• (byte) (64*2) and (byte) (127+1) yield -128, since $128 - (-128) = 1 \times 2^8$ .							
<ul> <li>Java defines the result of any arithmetic operation or conversion on integer types to "wrap around"—modular arithmetic.</li> <li>That is, the "next number" after the largest in an integer type is the smallest (like "clock arithmetic").</li> </ul>		• (byte) (101*99) yields 15, since $9999 - 15 = 39 \times \cdot 2^8$ .						
		• (byte) (-30*13) yields 122, since $-390 - 122 = -2 \times 2^8$ .						
		• (char) (-1) yields $2^{16} - 1$ , since $-1 - (2^{16} - 1) = -1 \times 2^{16}$ .						
• E.g., (byte) 128 == (byte) (127+1) == (byte)	-128							
• In general,								
<ul> <li>If the result of some arithmetic subexpression have type T, an n-bit integer type,</li> <li>then we compute the real (mathematical) value,</li> <li>and yield a number, x', that is in the range of the subscription of t</li></ul>	<i>x</i> ,							
equivalent to $x$ modulo $2^n$ .								
– (That means that $x - x'$ is a multiple of $2^n$ .)								

#### Modular Arithmetic and Bits

- Why wrap around?
- Java's definition is the natural one for a machine that uses binary arithmetic.
- For example, consider bytes (8 bits):

Decimal	Binary		
101	1100101		
×99	1100011		
9999	100111 00001111		
- 9984	100111 00000000		
15	00001111		

- In general, bit n, counting from 0 at the right, corresponds to  $2^n$ .
- The bits to the left of the vertical bars therefore represent multiples of  $2^8=256.$
- So throwing them away is the same as arithmetic module 256.

```
Last modified: Mon Sep 28 12:10:19 2015
```

```
CS61B: Lecture #14 5
```

#### Conversion

- In general Java will silently convert from one type to another if this makes sense and no information is lost from value.
- Otherwise, cast explicitly, as in (byte) x.
- Hence, given

byte aByte; char aChar; short aShort; int anInt; long aLong;

```
// OK:
```

aShort = aByte; anInt = aByte; anInt = aShort; anInt = aChar; aLong = anInt;

// Not OK, might lose information: anInt = aLong; aByte = anInt; aChar = anInt; aShort = anInt; aShort = aChar; aChar = aShort; aChar = aByte;

// OK by special dispensation: aByte = 13; // 13 is compile-time constant aByte = 12+100 // 112 is compile-time constant

- Why this representation for -1?
  - $\begin{array}{c|c|c} 1 & 00000001_2 \\ + & -1 & 1111111_2 \\ = & 0 & 1|00000000_2 \end{array}$

Only 8 bits in a byte, so bit 8 falls off, leaving 0.

- The truncated bit is in the  $2^8$  place, so throwing it away gives an equal number modulo  $2^8$ . All bits to the left of it are also divisible by  $2^8$ .
- $\bullet$  On unsigned types (char), arithmetic is the same, but we choose to represent only non-negative numbers modulo  $2^{16}$ :



```
Last modified: Mon Sep 28 12:10:19 2015
```

CS61B: Lecture #14 6

### Promotion

- $\bullet$  Arithmetic operations (+, \*, ...) promote operands as needed.
- Promotion is just implicit conversion.
- For integer operations,
  - if any operand is long, promote both to long.
  - otherwise promote both to int.

```
• So,
```

aByte + 3 == (int) aByte + 3 // Type int aLong + 3 == aLong + (long) 3 // Type long 'A' + 2 == (int) 'A' + 2 // Type int aByte = aByte + 1 // ILLEGAL (why?)

• But fortunately,

```
aByte += 1; // Defined as aByte = (byte) (aByte+1)
```

• Common example:

```
// Assume aChar is an upper-case letter
char lowerCaseChar = (char) ('a' + aChar - 'A'); // why cast?
```

# Bit twiddling

- Java (and C, C++) allow for handling integer types as sequences of bits. No "conversion to bits" needed: they already are.
- Operations and their uses:

	Mask		Set		Flip		Flip all
	00101100		00101100		00101100		
&	10100111	Ι	10100111	^	10100111	~	10100111
	00100100		10101111		10001011		01011000

• Shifting:

Left	Arithmetic Right	Logical Right			
10101101 << 3	10101101 >> 3	10101100 >>> 3			
01101000	11110101	00010101			
• What is: (-1) >>> x << n? x >> n? (x >>> 3)	29 <b>?</b> & ((1<<5)-1)?				

Last modified: Mon Sep 28 12:10:19 2015

CS61B: Lecture #14 9

•