

CS61B Lecture #14: Integers

Announcement:

- Programming contest SATURDAY! You can still sign up.
- Hackers@Berkeley "HackJam"—a 12 hour hackathon hosted by Hackers@Berkeley and sponsored by Box.
 - There will be food served throughout the event and prizes awarded at the end.
 - Who should come: Anyone interested in hacking, regardless of experience. There will be helpful students and engineers from Box there to help anyone who wants to learn.
 - Time: 11am-11pm Saturday, September 29th.
 - Place: Wozniak Lounge, Soda Hall.

Today: Integer Types; Readings: *A Java Reference*, §6.3-4. *Head First Java*, Chapter 10.

Readings for Upcoming Topics: *Data Structures (Into Java)*, Chapter 1.

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Integer Types and Literals

Type	Bits	Signed?	Literals
byte	8	Yes	
short	16	Yes	
char	16	No	'a' // (char) 97 '\n' // newline ((char) 10) '\t' // tab ((char) 8) '\' ' // backslash 'A', '\101', '\u0041' // == (char) 65
int	32	Yes	123 0100 // Octal for 64 0x3f, 0xffffffff // Hexadecimal 63, -1 (!)
long	64	Yes	123L, 01000L, 0x3fL 1234567891011L

- " N bits" means that there are 2^N integers in the domain of the type.
- If signed, range of values is $-2^{N-1} .. 2^{N-1} - 1$.
- If unsigned, only non-negative numbers, and range is $0..2^N - 1$.
- Negative numerals are just negated (positive) literals.
- Use casting for **byte** and **short**: (byte) 12, (short) 2000.

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Modular Arithmetic

- **Problem:** How do we handle overflow, such as occurs in $10000 * 10000 * 10000$?
- Some languages throw an exception (Ada), some give undefined results (C, C++)
- Java defines the result of any arithmetic operation or conversion on integer types to "wrap around"—*modular arithmetic*.
- That is, the "next number" after the largest in an integer type is the smallest (like "clock arithmetic").
- E.g., (byte) 128 == (byte) (127+1) == (byte) -128
- In general,
 - If the result of some arithmetic subexpression is supposed to have type T , an n -bit integer type,
 - then we compute the real (mathematical) value, x ,
 - and yield a number, x' , that is in the range of T , and that is equivalent to x modulo 2^n .
 - (That means that $x - x'$ is a multiple of 2^n .)

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Modular Arithmetic II

- (byte) (64*8) yields 0, since $512 - 0 = 2 \cdot 2^8$.
- (byte) (64*2) and (byte) (127+1) yield -128, since $128 - (-128) = 1 \cdot 2^8$.
- (byte) (345*6) yields 22, since $2070 - 22 = 8 \cdot 2^8$.
- (byte) (-30*13) yields 122, since $-390 - 122 = -2 \cdot 2^8$.
- (char) (-1) yields $2^{16} - 1$, since $-1 - (2^{16} - 1) = -1 \cdot 2^{16}$.
- Natural definition for a machine that uses binary arithmetic:

Type char	Type byte
	0 = 00000000 ₂
0 = 0000000000000000 ₂	1 = 00000001 ₂
$2^{16} - 1 = 1111111111111111_2$	127 = 01111111 ₂
	-128 = 10000000 ₂
	-1 = 11111111 ₂

- Terminology: rightmost (units) bit is *bit 0*, 2s bit is *bit 1*.
- Hence, changing bit n modifies value by 2^n ; truncating on left to n bits computes modulo 2^n .

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Negative numbers

- Why this representation for -1?

$$\begin{array}{r|l} 1 & 00000001_2 \\ + -1 & 11111111_2 \\ \hline = 0 & 1|00000000_2 \end{array}$$

Only 8 bits in a byte, so bit 8 falls off, leaving 0.

- The truncated bit is in the 2^8 place, so throwing it away gives an equal number modulo 2^8 . All bits to the left of it are also divisible by 2^8 .
- On unsigned types (**char**), arithmetic is the same, but we choose to represent only non-negative numbers modulo 2^{16} :

$$\begin{array}{r|l} 1 & 0000000000000001_2 \\ + 2^{16} - 1 & 1111111111111111_2 \\ \hline = 2^{16} + 0 & 1|0000000000000000_2 \end{array}$$

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Conversion

- In general Java will silently convert from one type to another if this makes sense and no information is lost from value.
- Otherwise, cast explicitly, as in (byte) x.
- Hence, given

```
byte aByte; char aChar; short aShort; int anInt; long aLong;
```

```
// OK:
```

```
aShort = aByte; anInt = aByte; anInt = aShort; anInt = aChar;
aLong = anInt;
```

```
// Not OK, might lose information:
```

```
anInt = aLong; aByte = anInt; aChar = anInt; aShort = anInt;
aShort = aChar; aChar = aShort; aChar = aByte;
```

```
// OK by special dispensation:
```

```
aByte = 13; // 13 is compile-time constant
aByte = 12+100 // 112 is compile-time constant
```

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Promotion

- Arithmetic operations (+, *, ...) *promote* operands as needed.
- Promotion is just implicit conversion.
- For integer operations,
 - if any operand is **long**, promote both to **long**.
 - otherwise promote both to **int**.
- So,

```
aByte + 3 == (int) aByte + 3 // Type int
aLong + 3 == aLong + (long) 3 // Type long
'A' + 2 == (int) 'A' + 2 // Type int
aByte = aByte + 1 // ILLEGAL (why?)
```

- But fortunately,

```
aByte += 1; // Defined as aByte = (byte) (aByte+1)
```

- Common example:

```
// Assume aChar is an upper-case letter
char lowerCaseChar = (char) ('a' + aChar - 'A'); // why cast?
```

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Bit twiddling

- Java (and C, C++) allow for handling integer types as sequences of bits. No "conversion to bits" needed: they already are.
- Operations and their uses:

Mask	Set	Flip	Flip all
00101100	00101100	00101100	
& 10100111	10100111	^ 10100111	~ 10100111
00100100	10101111	10001011	01011000

- Shifting:

Left	Arithmetic Right	Logical Right
10101101 << 3	10101101 >> 3	10101100 >>> 3
01101000	11110101	00010101

```
(-1) >>> 29?
```

- What is:
 - $x \ll n$?
 - $x \gg n$?

```
(x >>> 3) & ((1<<5)-1)?
```

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