Announcements

- •Homework 6 is due Tuesday 10/22 @ 11:59pm
- Includes a mid-semester survey about the course so far
- Project 3 is due Thursday 10/24 @ 11:59pm
- •Midterm 2 is on Monday 10/28 7pm-9pm
- Guerrilla section 3 this weekend
- Object-oriented programming, recursion, and recursive data structures
 2pm-5pm on Saturday and 10am-1pm on Sunday
- Please let us know you are coming by filling out the Piazza poll

61A Lecture 19

Friday, October 18

Comparing orders of growth (n is the problem size)

$\Theta(b^n) eq {\sf Exponential growth!} extrm{Recursive fib takes}$	
	$\Theta(\phi^n) \text{ steps, where } \phi = \frac{1+\sqrt{5}}{2} \approx 1.61828$ Incrementing the problem scales R(n) by a factor.
$\Theta(n^6)$	Incrementing the problem scales $R(n)$ by a factor.
$\Theta(n^2)$	Quadratic growth. E.g., operations on all pairs.
	Incrementing n increases $R(n)$ by the problem size $n_{\scriptscriptstyle \bullet}$
$\Theta(n)$	Linear growth. Resources scale with the problem.
$\Theta(\sqrt{n})$	
$\Theta(\log n)$	Logarithmic growth. These processes scale well.
	Doubling the problem only increments $R(n)$.
$\Theta(1)$ \checkmark Constant. The problem size doesn't matter.	

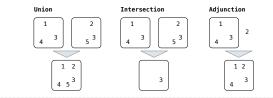
Comparing Orders of Growth

Sets One more built-in Python container type -Set literals are enclosed in braces -Duplicate elements are removed on construction -Sets are unordered, just like dictionary entries

>>> s = {3, 2, 1, 4, 4}
>>> s
{1, 2, 3, 4}
>>> 3 in s
True
>>> len(s)
4
>>> sunion({1, 5})
{1, 2, 3, 4, 5}
>>> s.intersection({6, 5, 4, 3})
{3, 4}

Sets

What we should be able to do with a set: •Membership testing: Is a value an element of a set? •Union: Return a set with all elements in set1 or set2 $\,\cdot\, Intersection\colon Return a set with any elements in set1 and set2$ ·Adjunction: Return a set with all elements in s and a value v



Sets as Unordered Sequences

Sets as Unordered Sequences

Proposal 1: A set is represented by a recursive list that contains no duplicate items.

def set_contains(s, v): if empty(s): return False elif s.first == v: return True else:

(Demo)

Time order of growth For a set operation that takes "linear" time, we say that def adjoin_set(s, v):
 if set_contains(s, v):
 return s $\Theta(n)$ n: size of the set The size of the set else: return Rlist(v, s) R(n): number of steps required to perform the operation def intersect_set(set1, set2): in_set2 = lambda v: set_contains(set2, v) return filter_rlist(set1, in_set2) $R(n) = \Theta(n) \qquad \text{An example f(n)}$ $\Theta(n^2)$ Assume sets are the same size which means that there are positive constants k_1 and k_2 such that def union_set(set1, set2):
 not_in_set2 = lambda v: not set_contains(set2, v)
 set1_not_set2 = filter_rlist(set1, not_in_set2)
 return extend_rlist(set1_not_set2, set2) $k_1 \cdot n \le R(n) \le k_2 \cdot n$ $\Theta(n^2)$ for sufficiently large values of n.

Implementing Sets

Implementing Sets

Sets as Unordered Sequences

Review: Order of Growth

def empty(s):
 return s is Rlist.empty

return set contains(s.rest, v)

(Demo)

Sets as Ordered Sequences

 $\ensuremath{\textbf{Proposal}}\xspace$ 2: A set is represented by a recursive list with unique elements ordered from least to greatest

Sets as Ordered Sequences

def set_contains(s, v):
 if empty(s) or s.first > v:
 return False
 elif s.first == v:
 return True
 else:
 return set_contains(s.rest, v)

Order of growth? $\Theta(n)$

Set Intersection Using Ordered Sequences

This algorithm assumes that elements are in order.

def intersect_set(set1, set2):
 if empty(set1) or empty(set2):
 return Rlist.empty
 else:
 el, e2 = set1.first, set2.first
 if e1 == e2:
 return Rlist(e1, intersect_set(set1.rest, set2.rest))
 elif e1 < e2:
 return intersect_set(set1.rest, set2)
 elif e2 < e1:
 return intersect_set(set1, set2.rest)
</pre>

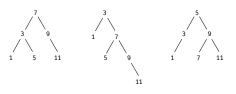
(Demo)

Order of growth? $\Theta(n)$

Sets as Binary Search Trees

Tree Sets

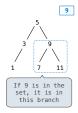
Proposal 3: A set is represented as a Tree. Each entry is: -Larger than all entries in its left branch and -Smaller than all entries in its right branch



Membership in Tree Sets

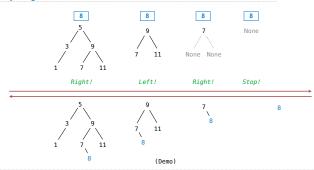
Set membership traverses the tree •The element is either in the left or right sub-branch •By focusing on one branch, we reduce the set by about half

def set_contains(s, v):
 if s is None:
 return False
 elif s.entry == v:
 return True
 elif s.entry < v:
 return set_contains(s.right, v)
 elif s.entry v:
 return set_contains(s.left, v)</pre>



Order of growth?

Adjoining to a Tree Set



More Set Operations

What Did I Leave Out?

Sets as ordered sequences: •Adjoining an element to a set •Union of two sets

Sets as binary trees: •Intersection of two sets •Union of two sets •Balancing a tree

That's all on homework 7!