# CS3: Introduction to Symbolic Programming 

Lecture 8:<br>The last bit of recursion Miniproject \#2

Fall 2006
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## Schedule

| 7 | Oct 9-13 | Advanced recursion |
| :--- | :--- | :--- |
| 8 | Oct 16-20 | Finishing recursion <br> Miniproject \#2: Number names |
| 9 | Oct 23-27 | Introduction to Higher Order Procedures |
| 10 | Oct 30-Nov 3 | More HOF |
| 11 | Nov 6-10 | Finish HOF <br> Miniproject \#3: Election processing |
| 12 | Nov 13-17 | Lecture: Midterm \#2 <br> Lab: Start on "Lists" |

## Any "notetaker" volunteers?

- A student in the course needs a note taker, which does pay a stipend. If you are taking notes anyway...
- Come and see me after lecture if interested


## The "screwed up" labs

## - This is the order things should have happened:

- First "advanced recursion" Lab: recursions with multiple arguments
-my-equal?, zipping, merging
- Second Lab
- patterns in recursion, no-vowels, sort (using insert), roman-sum-helper
- Last Lab

```
-mad-libs quiz, 1-extra?, fibonacci, thorough-
reversal
```


## Number Spelling Miniproject

- Read Simply Scheme, page 233, which has hints
- Another hint (principle): don't force "everything" into the recursion.
- Special/border cases may be easier to handle before you send yourself into a recursion
- Accumulating recursions are sometimes called "tail" recursions (by TAs, me, etc).
- But, not all recursions that keep track of a number are "tail" recursions.
- A tail recursion has no combiner, so it can end as soon as a base case is reached
- Compilers can do this efficiently
- An embedded recursion needs to combine up all the recursive steps to form the answer
- The poor compiler has to remember everything


## Tail or embedded? (1/3)

(define (length sent)
(if (empty? sent)
0
$(+1$ (length (bf sent)))))

## Embedded!

```
(length '(a b c d)) }
    (+ 1 (length '(b c d)))
    (+ 1 (+ 1 (length '(c d))))
    (+ 1 (+ 1 (+ 1 (length '(d)))))
    (+ 1 (+ 1 (+ 1 (+ 1 (length '())))))
    (+ 1 (+ 1 (+ 1 (+ 1 0))))
    (+ 1 (+ 1 (+ 1 1)))
    (+ 1 (+ 1 2))
    (+ 1 3)
4
```


## Tail or embedded? (2/3)

```
(define (sent-max sent)
    (if (empty? sent)
    '()
    (sent-max-helper (bf sent) (first sent))))
(define (sent-max-helper sent max-so-far)
    (if (empty? sent)
    max-so-far
    (sent-max-helper (bf sent)
    (if (> max-so-far (first sent))
    max-so-far
    (first sent)))))
```


## Tail or embedded? (3/3)

```
(define (find-evens sent)
    (cond ((empty? sent) ;base case
    '()
        )
    ((odd? (first sent)) ;rec case 1
        (find-evens (bf sent)) )
    (else ;rec case 2: even
    (se (first sent)
    (find-evens (bf sent))) )
    ))
```


## $>$ (find-evens '(2 $\left.\begin{array}{lllll}2 & 3 & 4 & 5 & 6\end{array}\right)$

$$
\text { sent }=(23456)
$$

(se 2
sent $=(3456$ )
sent $=(456)$
(se $4 \square$ sent $=(56)$
sent $=(6)$
(se 6
sent $=()$
()
$\rightarrow$ (se 2 (se 4 (se 6 ())
$\Rightarrow\left(\begin{array}{lll}2 & 4 & 6\end{array}\right)$

## Tree recursion: fibonacci

- The fibonacci sequence: $\begin{array}{llllllllll}1 & 1 & 2 & 3 & 5 & 8 & 13 & 21 & 34 & 55\end{array}$

```
(define (fib n)
    (if (<= n 2)
    1
    (+ (fib (- n 1)) ; r recursive case
    (fib (- n 2)))))
```


## Tree recursion: Pascals triangle

|  | columns (C) |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | 0 | 1 | 2 | 3 | 4 | 5 |  |
|  | 0 | 1 |  |  |  |  |  | $\ldots$ |
|  | 1 | 1 | 1 |  |  |  |  | $\ldots$ |
| $\bigcirc$ | 2 | 1 | 2 | 1 |  |  |  | $\ldots$ |
| S | 3 | 1 | 3 | 3 | 1 |  |  | $\ldots$ |
| (R) | 4 | 1 | 4 | 6 | 4 | 1 |  | $\ldots$ |
|  | 5 | 1 | 5 | 10 | 10 | 5 | 1 | $\ldots$ |
|  |  | $\ldots$ | $\ldots$ |  |  | $\ldots$ | $\ldots$ |  |

Pascal's Triangle

- How many ways can you choose $C$ things from $R$ choices?
- Coefficients of the $(x+y)^{\wedge} R$ : look in row $R$
- etc.
(define (pascal C R)
(cond
( ( = C 0) 1) ;base case
( (= C R) 1) ;base case
(else ;tree recurse
(+ (pascal C (- R 1))
(pascal (- C 1) (- R 1)) )
)))


## $>$ (pascal 2 5)

(pascal 2 5)
(+ (pascal 24 )
( +
$\underset{\left(+\begin{array}{lll}\text { (pascal } 2 & 2) & \rightarrow 1\end{array}\right]}{\left(\begin{array}{ll}\text { pascal } & 2\end{array}\right.}$
(pascal 1 2) (+
(pascal 1 3)
(pascal 12)
(pascal 02 2) $\rightarrow 1$
(pascal 1 4)

(pascal 0 3)
$\rightarrow \quad 1$

## pair-all

- Write pair-all, which takes a sentence of prefixes and a sentence of suffixes and returns a sentence pairing all prefixes to all suffixes.

```
- (pair-all '(a b c) '(1 2 3)) >
    (a1 b1 c1 a2 b2 c2 a3 b3 c3)
- (pair-all '(spr s k) '(ite at ing ong)) >
(sprite sprat spring sprong site sat sing
    song kite kat king kong)
```


## binary

- Write binary, a procedure to generate the possible binary numbers given $n$ bits.

```
(binary 1) 
(binary 2) }->(00\quad01 10 11
(binary 3) }->(000 001 010 011 100 101 110 111
```


## roman-sum-helper (from lab)

## Write roman-sum-helper:

```
(define (roman-sum number-sent)
    (if (empty? number-sent)
    0
    (roman-sum-helper (first number-sent)
        (bf number-sent)
        (first number-sent)) ) )
```

Roman-sum-helper takes three arguments:
(define (roman-sum-helper so-far number-list mostrecent) ... )
(roman-sum ' (100 $\left.10 \begin{array}{llll}10 & 1 & 5\end{array}\right)$ ) will recurse with:
(roman-sum-helper 100 '(10 50 1 5) 100)
(roman-sum-helper 110 '(50 1 5) 10)
(roman-sum-helper 140 '(1 5) 50)
(roman-sum-helper 141 '(5) 1)
(roman-sum-helper 156 '( ) 5)

