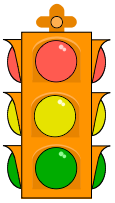
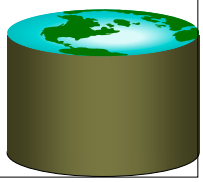
Concurrency Control

R &G - Chapter 19



Smile, it is the key that fits the lock of everybody's heart.

Anthony J. D'Angelo, The College Blue Book



Review

- **DBMSs support concurrency, crash recovery with:**
  - ACID Transactions
  - Log of operations
- **A serial execution of transactions is safe but slow**
  - Try to find schedules equivalent to serial execution
- **One solution for serializable schedules is 2PL**

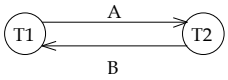
Conflict Serializable Schedules

- **Two schedules are conflict equivalent if:**
  - Involve the same actions of the same transactions
  - Every pair of conflicting actions is ordered the same way
- **Schedule S is conflict serializable if S is conflict equivalent to some serial schedule**

Example

- **A schedule that is not conflict serializable:**

T1:	R(A), W(A),	R(B), W(B)
T2:	R(A), W(A), R(B), W(B)	



Dependency graph

- **The cycle in the graph reveals the problem. The output of T1 depends on T2, and vice-versa.**

Dependency Graph

- ***Dependency graph*: One node per Xact; edge from *Ti* to *Tj* if an operation of Ti conflicts with an operation of Tj and Ti's operation appears earlier in the schedule than the conflicting operation of Tj.**
- **Theorem: Schedule is conflict serializable if and only if its dependency graph is acyclic**

An Aside: View Serializability

- **Schedules S1 and S2 are view equivalent if:**
  - If Ti reads initial value of A in S1, then Ti also reads initial value of A in S2
  - If Ti reads value of A written by Tj in S1, then Ti also reads value of A written by Tj in S2
  - If Ti writes final value of A in S1, then Ti also writes final value of A in S2

T1: R(A)		W(A)
T2:	W(A)	
T3:		W(A)

T1: R(A),W(A)		
T2:	W(A)	
T3:		W(A)

- **View serializability is "weaker" than conflict serializability!**
  - Every conflict serializable schedule is view serializable, but not vice versa!
  - I.e. admits more legal schedules

App-Specific Serializability

- **In some cases, application logic can deal with apparent conflicts**
  - E.g. when all writes commute
    - E.g. increment/decrement (a.k.a. "escrow transactions")

T1:	x=R(A), W(A=x+1),	z=R(A), W(z=z+1)
T2:	y=R(A), W(A=y-1)	

    - Note: doesn't work in some cases for (American) bank accounts
      - Account cannot go below $0.00!!
- **In general, this kind of app logic is not known to DBMS**
  - Only sees encapsulated R/W requests
  - But keep in mind that general serializability is "weaker" than even view serializability

Review: Strict 2PL

Lock Compatibility Matrix

	S	X
S	√	–
X	–	–

- ***Strict Two-phase Locking (Strict 2PL) Protocol*:**
  - Each Xact must obtain a S (*shared*) lock on object before reading, and an X (*exclusive*) lock on object before writing.
  - All locks held by a transaction are released when the transaction completes
  - If an Xact holds an X lock on an object, no other Xact can get a lock (S or X) on that object.
- **Strict 2PL allows only schedules whose precedence graph is acyclic**

Two-Phase Locking (2PL)

- **Two-Phase Locking Protocol**
  - Each Xact must obtain a S (*shared*) lock on object before reading, and an X (*exclusive*) lock on object before writing.
  - A transaction can not request additional locks once it releases any locks.
  - If a Xact holds an X lock on an object, no other Xact can get a lock (S or X) on that object.
- **Can result in Cascading Aborts!**
  - STRICT (!!) 2PL "Avoids Cascading Aborts" (ACA)

Lock Management

- **Lock and unlock requests are handled by the lock manager**
- **Lock table entry:**
  - Number of transactions currently holding a lock
  - Type of lock held (shared or exclusive)
  - Pointer to queue of lock requests
- **Locking and unlocking have to be atomic operations**
  - requires latches ("semaphores"), which ensure that the process is not interrupted while managing lock table entries
  - see CS162 for implementations of semaphores
- **Lock upgrade: transaction that holds a shared lock can be upgraded to hold an exclusive lock**
  - Can cause deadlock problems

Deadlocks

- **Deadlock: Cycle of transactions waiting for locks to be released by each other.**
- **Two ways of dealing with deadlocks:**
  - Deadlock prevention
  - Deadlock detection

Deadlock Prevention

- **Assign priorities based on timestamps. Assume Ti wants a lock that Tj holds. Two policies are possible:**
  - Wait-Die: If Ti has higher priority, Ti waits for Tj; otherwise Ti aborts
  - Wound-wait: If Ti has higher priority, Tj aborts; otherwise Ti waits
- **If a transaction re-starts, make sure it gets its original timestamp**
  - Why?

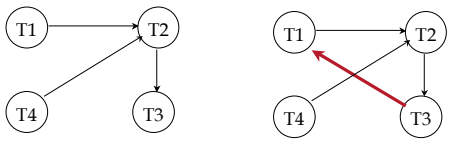
Deadlock Detection

- **Create a waits-for graph:**
  - Nodes are transactions
  - There is an edge from Ti to Tj if Ti is waiting for Tj to release a lock
- **Periodically check for cycles in the waits-for graph**

Deadlock Detection (Continued)

Example:

T1: S(A), S(D), S(B)
T2: X(B) X(C)
T3: S(D), S(C), X(A)
T4: X(B)



Deadlock Detection (cont.)

- **In practice, most systems do detection**
  - Experiments show that most waits-for cycles are length 2 or 3
  - Hence few transactions need to be aborted
  - Implementations can vary
    - Can construct the graph and periodically look for cycles
    - Can do a "time-out" scheme: if you've been waiting on a lock for a long time, assume you're deadlock and abort

Summary

- **Correctness criterion for isolation is "serializability".**
  - In practice, we use "conflict serializability", which is somewhat more restrictive but easy to enforce.
- **There are several lock-based concurrency control schemes (Strict 2PL, 2PL). Locks directly implement the notions of conflict.**
  - The lock manager keeps track of the locks issued. Deadlocks can either be prevented or detected.

Things We're Glossing Over

- **What should we lock?**
  - We assume tuples here, but that can be expensive!
  - If we do table locks, that's too conservative
  - *Multi-granularity* locking
- **Locking in indexes**
  - don't want to lock a B-tree root for a whole transaction!
  - actually do non-2PL "latches" in B-trees
- **CC w/out locking**
  - "optimistic" concurrency control
  - "timestamp" and multi-version concurrency control
  - locking usually better, though
- **App-specific tricks**
  - e.g. increment/decrement ("escrow transactions")

In case we have time

- **The following is an interesting problem**
- **We will not discuss how to solve it, though!**

Dynamic Databases – The "Phantom" Problem

- **If we relax the assumption that the DB is a fixed collection of objects, even Strict 2PL (on individual items) will not assure serializability:**
- **Consider T1 – "Find oldest sailor for rating 1"**
  - T1 locks all pages containing sailor records with *rating* = 1, and finds oldest sailor (say, *age* = 71).
    - May find these pages via a clustered index on rating, and never touch (or lock) any other pages
  - Next, T2 inserts a new sailor; *rating* = 1, *age* = 96.
  - T2 commits.
  - T1 issues another query to find the oldest sailor for rating 1.
    - A phantom sailor appears! (and she's 96 years old!)
- **No serial execution where T1's result could happen!**

The Problem

- **T1 implicitly assumes that it has locked the set of all sailor records with *rating* = 1.**
  - Assumption only holds if no sailor records are added while T1 is executing!
  - Need some mechanism to enforce this assumption. (Index locking and predicate locking.)
- **Example shows that conflict serializability guarantees serializability only if the set of objects is fixed!**
  - e.g. table locks

Predicate Locking

- **Grant lock on all records that satisfy some logical predicate, e.g. *age > 2*salary*.**
- **Index locking is a special case of predicate locking for which an index supports efficient implementation of the predicate lock.**
  - What is the predicate in the sailor example?
- **In general, predicate locking has a lot of locking overhead.**
  - too expensive!
  - Fancier index locking tricks are used in practice