## UCB CS61C : Machine Structures

## Lecture 07 Introduction to MIPS : Decisions II

## 2011-09-12

Hello to Dr Mauro Sgarzi from Italy!!

## KINECT? YOUR BODY IS ANTENNA!

Researchers at Microsoft and UW are working on a system that uses the fact that your body can act as an antenna and notes how ambient electric fields change to figure out what your position or motion was. The idea is you don't need a camera or Wimote to interact with it!


WWW . nytimes.com/2011/09/11/business/using-gestures-to-control-electronic-devices.html

## Review

- Memory is byte-addressable, but 1w and sw access one word at a time.
- A pointer (used by 1 w and sw) is iust a memory address so we can add to it or subtract from it (using offset).
- A Decision allows us to decide what to execute at run-fime rather than compile-fime.
- C Decisions are made using conditional statements Within if, while, do while, for.
- MIPS Decision mqking instructions are the condifional branches: beq and bne.
- New Instructions:
lw, sw, beq, bne, j


## Last time: Loading, Storing bytes 1/2

- In addition to word data transfers
(lw, sw), MIPS has byte data transfers:
- load byte: 1.b
- store byte: sb
- same format as lw, sw
- E.g., lb \$s0, 3 (\$s1)
- contents of memory location with address = sum of " 3 " + contents of register s1 is copied to the low byte position of register so.


## Loading, Storing bytes 2/2

- What do with other 24 bits in the 32 bit register?
- lb: sign extends to fill upper 24 bits


##  ...is copied to "sign-extend" <br> This bit <br> 

- Normally don’† want to sign extend chars
- MIPS instruction that doesn' $\dagger$ sign extend when loading bytes:
- load byte unsigned:


## Overfiow in Arithmetic (1/2)

- Reminder: Overflow occurs when there is a mistake in arithmetic due to the limited precision in computers.
- Example (4-bit unsigned numbers):

| 15 | 1111 |
| ---: | ---: |
| $+\quad 3$ |  |
| 18 | 10011 |

- But we don't have room for 5 -bit solution, so the solution would be 0010 , which is +2 , and wrong.


## Overfiow in Arithmetic (2/2)

- Some languages detect overflow (Ada), some don't (C)
- MIPS solution is 2 kinds of arithmetic instructs:
- These cause overflow to be detected
- add (add)
- add immediate (addi)
- subtract (sub)
- These do not cause overflow detection
- add unsigned (addu)
- add immediate unsigned (addiu)
- subtract unsigned (subu)
- Compiler selects appropriate arithmetic
- MIPS C compilers produce


## Two "Logic" Instructions

- Here are 2 more new instructions
- Shiff Left: sll \$s1, \$s2,2 \#s1=s2<<2
- Store in \$s1 the value from \$s2 shifted 2 bits to the left (they fall off end), inserting 0's on right; << in C.
- Before: $00000002_{\text {hex }}$ $00000000000000000000000000000010_{\text {two }}$
- After: 0000000 B hex $^{\text {a }}$ $00000000000000000000000000001000_{\text {two }}$
- What arithmetic effect does shift left have?
- Shiff Right: srl. is opposite shiff; >>


## Loops in C/Assembly (1/3)

- Simple loop in C; A [ ] is an array of ints

$$
\text { do } \begin{aligned}
&\left\{\begin{array}{l}
g
\end{array}=g+\hat{A}[i] ;\right. \\
& i=i+j ;
\end{aligned}
$$

$$
\text { \} while (i }!=\mathrm{h} \text { ); }
$$

- Rewrite this as:

$$
\begin{aligned}
\text { Loop: } & g=g+\mathbb{a}[i] ; \\
& =1+j ; \\
& \text { if }(i \quad!=h) \text { goto Loop; }
\end{aligned}
$$

- Use this mapping:

$$
\begin{aligned}
& \mathrm{g}, \mathrm{~h}, \mathrm{i}, \\
& \mathrm{\$ s} 1, \mathrm{j}, \mathrm{~b} 2, \\
& \$ \mathrm{~s} 4, \\
& \$ s 5
\end{aligned}
$$

## Loops in C/Assembly (2/3)

- Final compiled MIPS code:

Loop:
sll \$t1,\$s3,2
\# \$t1= $4 * I$
addu \$t1,\$t1,\$s5
\# \$t1=addr A+4i
lw \$t1,0(\$t1)
\# \$t1=A[i]
addu \$s1,\$s1,\$t1
\# $g=g+A[i]$
addu \$s3,\$s3,\$s4
\# $i=i+j$
bne \$s3,\$s2,Loop
\# goto Loop
\# if i!=h

- Original code:

$$
\begin{aligned}
\text { Loop: } & g=g+[i] ; \\
& \text { i }=i+j ; \\
& \text { if }(\quad!=h) \text { goto Loop ; }
\end{aligned}
$$

## Loops in C/Assembly (3/3)

- There are three types of loops in C:
- while
- do ... while
- for
- Each can be rewritten as either of the other two, so the method used in the previous example can be applied to these loops as well.
- Key Concept: Though there are multiple ways of writing a loop in MIPS, the key to decisionmaking is conditional branch


## Administrivia

- The schedule through week 7 has been determined
- Midterm 7-9pm on 2011-10-06
- Other administrivia?


## Inequalifies in MIPS (1/4)

- Until now, we've only tested equalities (== and $!=$ in C). General programs need to test < and > as well.
- Introduce MIPS Inequality Instruction:
- "Set on Less Than"
- Syntax: slt reg1,reg2,reg3
- Meaning: reg1 = (reg2 < reg3) ;

"reset" means "change to 0 ".


## Inequalifies in MPS (2/4)

- How do we use this? Compile by hand:

$$
\text { if ( } \mathrm{g}<\mathrm{h} \text { ) goto Less; \#g:\$s0, h:\$s1 }
$$

- Answer: compiled MIPS code...

$$
\begin{aligned}
\text { slt } \$ t 0, \$ s 0, \$ s 1 & \# \text { \$t0 }=1 \text { if } g<h \\
\text { bne } \$ t 0, \$ 0, \text { Less } & \text { \# goto Less } \\
& \# \text { if } \$ t 0!=0 \\
& \# \text { (if }(g<h)) \text { Less: }
\end{aligned}
$$

- Register $\$ 0$ always contains the value 0 , so bne and beq often use it for comparison after an sle instruction.
A slt $\Rightarrow$ bne pair means if (... < ...) goto...


## Inequalifies in MPS (3/4)

- Now we can implement <, but how do we implement $>, \leq$ and $\geq$ ?
- We could add 3 more instructions, but:
- MIPS goal: Simpler is Better
- Can we implement $\leq$ in one or more instructions using just slt and branches?
- What about >?
- What about $\geq$ ?


## Inequalifies in MIPS (4/4)

$$
\begin{aligned}
& \# \text { a:\$s0, } b: \$ s 1 \\
& \text { slt } \$ t 0, \$ s 0, \$ s 1 \# \text { \$t0 }=1 \text { if } a<b \\
& \text { beq } \$ t 0, \$ 0, \text { skip } \# \text { skip if } a>=b \\
& \quad<\text { stuff> } \\
& \text { skip: }
\end{aligned}
$$

Two independent variations possible:
Use slt \$t0, \$s1,\$s0 instead of
slt \$t0,\$s0, \$s1
Use bne instead of bea

## Immediates in Inequallities

- There is also an immediate version of slt to test against constants: slti
- Helpful in for loops


## C if (g >= 1) goto Loop

M
I slti $\$$ to, $\$ \mathbf{s} 0,1 \quad \#$ \$to $=1$ if
P $\# \$ s 0<1 \quad(g<1)$
S beq \$t0,\$0,Loop \# goto Loop

## What about unsigned numbers?

- Also unsigned inequality instructions:
sltu, sltiu
... which sets result to 1 or 0 depending on unsigned comparisons
- What is value of $\$ \mathrm{t0}$, $\$ \mathrm{t}$ ?

$$
\begin{gathered}
\left(\$ s 0=\text { FFFF }_{\text {FFFA }}^{\text {hex }} \text {, } \$ s 1=0000 \mathrm{FFFA}_{\text {hex }}\right) \\
\text { slt } \$ t 0, \$ s 0, \$ \mathrm{~s} 1 \\
\text { sltu } \$ \mathrm{t} 1, \$ s 0, \$ s 1
\end{gathered}
$$

## MIPS Signed vs. Unsigned - diff meanings!

- MIPS terms Signed/Unsigned "overloaded":
- Do/Don't sign extend
" (lb, lbu)
- Do/Don't overflow
" (add, addi, sub, mult, div)
" (addu, addiu, subu, multu, divu)
- Do signed/unsigned compare
" (slt, slti/sltu, sltiu)


## Peer Instruction

Ioop: addi \$s0,\$s0,-1 slti \$t0,\$s1,2 beq \$t0,\$0, Loop slt $\$ t 0, \$ s 1, \$ s 0$ bne $\$ t 0, \$ 0$, Loop \# goto Loop if \$t0 != 0
(\$s0=i, \$s1=j)
What C code properly fills in the blank in loop below? do \{i--; \} while(__);
$\# i=i-1$
$\# \$ t 0=(j<2)$
\# goto Loop if \$t0 == 0
$\# \$ t 0=(j<i)$

## "And in conclusion..."

- To help the conditional branches make decisions concerning inequalities, we introduce: "Set on Less Than" called
slt, slti, sltu, sltiu
- One can store and load (signed and unsigned) bytes as well as words with 1.b, 1.bu
- Unsigned add/sub don’t cause overflow
- New MIPS Instructions:

$$
\begin{aligned}
& \text { sll, srl, lb, lbu } \\
& \text { slt, slti, sltu, sltiu } \\
& \text { addu, addiu, subu }
\end{aligned}
$$

## Bonus Slides



## Example: The C Switch Statement (1/3)

- Choose among four alternatives depending on whether k has the value $0,1,2$ or 3 . Compile this C code:

```
switch (k) {
    case 0: f=i+j; break; /* k=0 */
    case 1: f=g+h; break; /* k=1 */
    case 2: f=g-h; break; /* k=2 */
    case 3: f=i-j; break; /* k=3 */
}
```


## Example: The C Switch Statement (2/3)

- This is complicated, so simplify.
- Rewrite it as a chain of if-else statements, which we already know how to compile:

$$
\begin{aligned}
& \text { if (k==0) f=i+j; } \\
& \text { else if(k==1) f=g+h; } \\
& \quad \text { else if(k==2) } f=g-h ; \\
& \quad \text { else if(k==3) } f=i-j ;
\end{aligned}
$$

- Use this mapping:

$$
\begin{gathered}
f: \$ s 0, ~ g: \$ s 1, \quad h: \$ s 2, \\
i: \$ s 3, j: \$ s 4, k: \$ s 5
\end{gathered}
$$

## Example: The C Switch Statement (3/3)

- Final compiled MIPS code:


## bne $\$ \mathbf{s} 5, \$ 0$, <br> add $\$ \mathbf{s} 0, \$ \mathrm{~s} 3, \$ \mathrm{~s} 4$ <br> j

L1: addi \$t0,\$s5,-1
bne \$t0,\$0,L2 \# branch k!=1
add $\$ \mathbf{s} 0, \$ \mathrm{~s} 1, \$ \mathrm{~s} 2$
j
L2: addi \$t0,\$s5,-2
bne \$t0,\$0,L3
sub \$s0,\$s1,\$s2
j
addi \$t0,\$s5,-3
bne \$t0,\$0,
sub \$s0,\$s3,\$s4
\# branch k!=0
\#k==0 so $f=i+j$
\# end of case so Exit
\# $\$ t 0=k-1$
\#k==1 so $f=g+h$
\# end of case so Exit
\# \$t0=k-2
\# branch k!=2
\#k==2 so $f=g-h$
\# end of case so Exit $\$ t 0=k-3$

