


inst.eecs.berkeley.edu/~cs61c
UCB CS61C : Machine Structures

Lecture 07
Introduction to MIPS : Decisions II

2011-09-12




Lecturer SOE
Dan Garcia

Hello to Dr Mauro Sgarzi from Italy!!

KINECT? YOUR BODY IS ANTENNA!

Researchers at Microsoft and UW are working on a system that uses the fact that your body can act as an antenna and notes how ambient electric fields change to figure out what your position or motion was. The idea is you don't need a camera or Wiimote to interact with it!



www.nytimes.com/2011/09/11/business/using-gestures-to-control-electronic-devices.html

Review

- Memory is byte-addressable, but **lw** and **sw** access one word at a time.
- A pointer (used by **lw** and **sw**) is just a memory address, so we can add to it or subtract from it (using offset).
- A Decision allows us to decide what to execute at run-time rather than compile-time.
- C Decisions are made using conditional statements within **if**, **while**, **do while**, **for**.
- MIPS Decision making instructions are the conditional branches: **beq** and **bne**.
- New Instructions:
 - lw**, **sw**, **beq**, **bne**, **j**

Cal CS61C L07 Introduction to MIPS : Decisions II (2) Garcia, Fall 2011 © UCB


Last time: Loading, Storing bytes 1/2

- In addition to word data transfers (**lw**, **sw**), MIPS has byte data transfers:
 - load byte: **lb**
 - store byte: **sb**
- same format as **lw**, **sw**
- E.g., **lb \$s0, 3(\$s1)**
 - contents of memory location with address = sum of "3" + contents of register **s1** is copied to the low byte position of register **s0**.

Cal CS61C L07 Introduction to MIPS : Decisions II (3) Garcia, Fall 2011 © UCB

Loading, Storing bytes 2/2

- What do with other 24 bits in the 32 bit register?
 - lb: sign extends to fill upper 24 bits



- Normally don't want to sign extend chars
- MIPS instruction that doesn't sign extend when loading bytes:
 - load byte unsigned: **lbu**

Cal CS61C L07 Introduction to MIPS : Decisions II (4) Garcia, Fall 2011 © UCB

Overflow in Arithmetic (1/2)

- Reminder: Overflow occurs when there is a mistake in arithmetic due to the limited precision in computers.
- Example (4-bit unsigned numbers):

15	1111
+ 3	+ 0011
18	10010

 - But we don't have room for 5-bit solution, so the solution would be 0010, which is +2, and wrong.

Cal CS61C L07 Introduction to MIPS : Decisions II (5) Garcia, Fall 2011 © UCB

Overflow in Arithmetic (2/2)

- Some languages detect overflow (Ada), some don't (C)
- MIPS solution is 2 kinds of arithmetic instructs:
 - These cause overflow to be detected
 - add (**add**)
 - add immediate (**addi**)
 - subtract (**sub**)
 - These do not cause overflow detection
 - add unsigned (**addu**)
 - add immediate unsigned (**addiu**)
 - subtract unsigned (**subu**)
- Compiler selects appropriate arithmetic
 - MIPS C compilers produce **addu**, **addiu**, **subu**

Cal CS61C L07 Introduction to MIPS : Decisions II (6) Garcia, Fall 2011 © UCB

Two "Logic" Instructions

- Here are 2 more new instructions
- Shift Left: `sll $s1, $s2, 2 #s1=s2<<2`
 - Store in `$s1` the value from `$s2` shifted 2 bits to the left (they fall off end), inserting 0's on right; `<<` in C.
 - Before: `0000 0002hex`
`0000 0000 0000 0000 0000 0000 0010two`
 - After: `0000 0008hex`
`0000 0000 0000 0000 0000 0000 1000two`
 - What arithmetic effect does shift left have?
- Shift Right: `srl` is opposite shift; `>>`

Cal

CS41C L07 Introduction to MIPS : Decision 2 (7)

Gerds, Fall 2011 © UCS

Loops in C/Assembly (1/3)

- Simple loop in C; `A[]` is an array of `ints`

```
do { g = g + A[i];
    i = i + j;
} while (i != h);
```
- Rewrite this as:

```
Loop: g = g + A[i];
      i = i + j;
      if (i != h) goto Loop;
```
- Use this mapping:

```
g,   h,   i,   j, base of A
$s1, $s2, $s3, $s4, $s5
```

Cal

CS41C L07 Introduction to MIPS : Decision 2 (8)

Gerds, Fall 2011 © UCS

Loops in C/Assembly (2/3)

- Final compiled MIPS code:

```
Loop: sll $t1,$s3,2 # $t1= 4*I
      addu $t1,$t1,$s5 # $t1=addr A+4i
      lw  $t1,0($t1) # $t1=A[i]
      addu $s1,$s1,$t1 # g=g+A[i]
      addu $s3,$s3,$s4 # i=i+j
      bne $s3,$s2,Loop # goto Loop
                        # if i!=h
```
- Original code:

```
Loop: g = g + A[i];
      i = i + j;
      if (i != h) goto Loop;
```

Cal

CS41C L07 Introduction to MIPS : Decision 2 (9)

Gerds, Fall 2011 © UCS

Loops in C/Assembly (3/3)

- There are three types of loops in C:
 - `while`
 - `do ... while`
 - `for`
- Each can be rewritten as either of the other two, so the method used in the previous example can be applied to these loops as well.
- Key Concept: Though there are multiple ways of writing a loop in MIPS, the key to decision-making is conditional branch

Cal

CS41C L07 Introduction to MIPS : Decision 2 (10)

Gerds, Fall 2011 © UCS

Administrivia

- The schedule through week 7 has been determined
 - Midterm 7-9pm on 2011-10-06
- Other administrivia?

Cal

CS41C L07 Introduction to MIPS : Decision 2 (11)

Gerds, Fall 2011 © UCS

Inequalities in MIPS (1/4)

- Until now, we've only tested equalities (`==` and `!=` in C). General programs need to test `<` and `>` as well.
 - Introduce MIPS Inequality Instruction:
 - "Set on Less Than"
 - Syntax: `slt reg1, reg2, reg3`
 - Meaning: `reg1 = (reg2 < reg3);`
- ```
if (reg2 < reg3)
 reg1 = 1;
else reg1 = 0;
```
- ← Same thing...
- "set" means "change to 1",  
"reset" means "change to 0".

Cal

CS41C L07 Introduction to MIPS : Decision 2 (12)

Gerds, Fall 2011 © UCS

## Inequalities in MIPS (2/4)

- How do we use this? Compile by hand:  
`if (g < h) goto Less; #g:$s0,h:$s1`
- Answer: compiled MIPS code...  

```
slt $t0,$s0,$s1 # $t0 = 1 if g<h
bne $t0,$0,Less # goto Less
 # if $t0!=0
 # (if (g<h)) Less:
```
- Register \$0 always contains the value 0, so `bne` and `beq` often use it for comparison after an `slt` instruction.
- A `slt` → `bne` pair means `if(... < ...) goto...`

Cal

CS41C L07 Introduction to MIPS : Decisions II (2)

Gerds, Fall 2011 © UCS

## Inequalities in MIPS (3/4)

- Now we can implement `<`, but how do we implement `>`, `≤` and `≥`?
- We could add 3 more instructions, but:
  - MIPS goal: Simpler is Better
- Can we implement `≤` in one or more instructions using just `slt` and branches?
  - What about `>`?
  - What about `≥`?

Cal

CS41C L07 Introduction to MIPS : Decisions II (3)

Gerds, Fall 2011 © UCS

## Inequalities in MIPS (4/4)

```
a:$s0, b:$s1
slt $t0,$s0,$s1 # $t0 = 1 if a<b
beq $t0,$0,skip # skip if a >= b
<stuff> # do if a<b
skip:
```

Two independent variations possible:

Use `slt $t0,$s1,$s0` instead of

`slt $t0,$s0,$s1`

Use `bne` instead of `beq`

Cal

CS41C L07 Introduction to MIPS : Decisions II (4)

Gerds, Fall 2011 © UCS

## Immediates in Inequalities

- There is also an immediate version of `slt` to test against constants: `slti`
  - Helpful in `for` loops

```
C if (g >= 1) goto Loop
```

```
M Loop: . . .
```

```
I slti $t0,$s0,1 # $t0 = 1 if
```

```
P # $s0<1 (g<1)
```

```
S beq $t0,$0,Loop # goto Loop
```

```
 # if $t0==0
```

```
 # (if (g>=1))
```

Cal

An `slt` → `beq` pair means `if(... ≥ ...) goto...`

CS41C L07 Introduction to MIPS : Decisions II (5)

Gerds, Fall 2011 © UCS

## What about unsigned numbers?

- Also unsigned inequality instructions:  
`sltu, sltiu`  
... which sets result to 1 or 0 depending on unsigned comparisons
- What is value of `$t0, $t1`?  

```
($s0 = FFFF FFFAhex, $s1 = 0000 FFFAhex)
 slt $t0, $s0, $s1
 sltu $t1, $s0, $s1
```

Cal

CS41C L07 Introduction to MIPS : Decisions II (7)

Gerds, Fall 2011 © UCS

## MIPS Signed vs. Unsigned – diff meanings!

- MIPS terms Signed/Unsigned “overloaded”:
  - Do/Don't sign extend
    - (`lb, lbu`)
  - Do/Don't overflow
    - (`add, addi, sub, mult, div`)
    - (`addu, addiu, subu, multu, divu`)
  - Do signed/unsigned compare
    - (`slt, slti/sltu, sltiu`)

Cal

CS41C L07 Introduction to MIPS : Decisions II (8)

Gerds, Fall 2011 © UCS

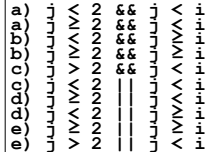
## Peer Instruction

```
Loop:addi $s0,$s0,-1 # i = i - 1
 slti $t0,$s1,2 # $t0 = (j < 2)
 beq $t0,$0,Loop # goto Loop if $t0 == 0
 slt $t0,$s1,$s0 # $t0 = (j < i)
 bne $t0,$0,Loop # goto Loop if $t0 != 0
```

(\$s0=i, \$s1=j)

What C code properly fills in the blank in loop below?

do {i--;} while(\_\_\_\_);



Cal

CS41C L07 Introduction to MIPS : Decisions II (17)

Gerds, Fall 2011 © UCSC

## "And in conclusion..."

- To help the conditional branches make decisions concerning inequalities, we introduce: "Set on Less Than" called **slt, slti, sltu, sltiu**
- One can store and load (signed and unsigned) bytes as well as words with **lb, lbu**
- Unsigned add/sub don't cause overflow
- New MIPS Instructions:

```
sll, srl, lb, lbu
slt, slti, sltu, sltiu
addu, addiu, subu
```

Cal

CS41C L07 Introduction to MIPS : Decisions II (20)

Gerds, Fall 2011 © UCSC

## Bonus Slides

Cal

CS41C L07 Introduction to MIPS : Decisions II (21)

Gerds, Fall 2011 © UCSC

## Example: The C Switch Statement (1/3)

- Choose among four alternatives depending on whether *k* has the value 0, 1, 2 or 3. Compile this C code:

```
switch (k) {
 case 0: f=i+j; break; /* k=0 */
 case 1: f=g+h; break; /* k=1 */
 case 2: f=g-h; break; /* k=2 */
 case 3: f=i-j; break; /* k=3 */
}
```

Cal

CS41C L07 Introduction to MIPS : Decisions II (22)

Gerds, Fall 2011 © UCSC

## Example: The C Switch Statement (2/3)

- This is complicated, so simplify.
- Rewrite it as a chain of if-else statements, which we already know how to compile:

```
if(k==0) f=i+j;
 else if(k==1) f=g+h;
 else if(k==2) f=g-h;
 else if(k==3) f=i-j;
```

- Use this mapping:

```
f:$s0, g:$s1, h:$s2,
i:$s3, j:$s4, k:$s5
```

Cal

CS41C L07 Introduction to MIPS : Decisions II (23)

Gerds, Fall 2011 © UCSC

## Example: The C Switch Statement (3/3)

- Final compiled MIPS code:

```
bne $s5,$0,L1 # branch k!=0
add $s0,$s3,$s4 #k==0 so f=i+j
j Exit # end of case so Exit
L1: addi $t0,$s5,-1 # $t0=k-1
 bne $t0,$0,L2 # branch k!=1
 add $s0,$s1,$s2 #k==1 so f=g+h
 j Exit # end of case so Exit
L2: addi $t0,$s5,-2 # $t0=k-2
 bne $t0,$0,L3 # branch k!=2
 sub $s0,$s1,$s2 #k==2 so f=g-h
 j Exit # end of case so Exit
L3: addi $t0,$s5,-3 # $t0=k-3
 bne $t0,$0,Exit # branch k!=3
 sub $s0,$s3,$s4 # k==3 so f=i-j
Exit:
```

Cal

CS41C L07 Introduction to MIPS : Decisions II (24)

Gerds, Fall 2011 © UCSC