

Recreation

Prove that for every acute angle $\alpha > 0$,

$$\tan \alpha + \cot \alpha \geq 2$$

CS61B Lecture #5: Simple Pointer Manipulation

Announcement

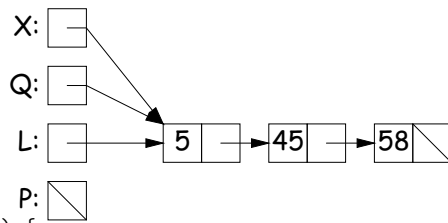
- **Today:** More pointer hacking.
- **Handing in labs and homework:** We'll be lenient about accepting late homework and labs for the first few. Just get it done: part of the point is getting to understand the tools involved. We will *not* accept submissions by email.

Destructive Incrementing

Destructive solutions may modify objects in the original list to save time or space:

```
/** List of all items in P incremented by n. May destroy original. */
static IntList dincrList(IntList P, int n) {
    if (P == null)
        return null;
    else {
        P.head += n;
        P.tail = dincrList(P.tail, n);
        return P;
    }
}
```

```
/** List L destructively incremented
 * by n. */
static IntList dincrList(IntList L, int n) {
    // 'for' can do more than count!
    for (IntList p = L; p != null; p = p.tail)
        p.head += n;
    return L;
}
```



Another Example: Non-destructive List Deletion

If L is the list [2, 1, 2, 9, 2], we want `removeAll(L, 2)` to be the new list [1, 9].

```
/** The list resulting from removing all instances of X from L
 * non-destructively. */
static IntList removeAll(IntList L, int x) {
    if (L == null)
        return null;
    else if (L.head == x)
        return removeAll(L.tail, x);
    else
        return new IntList(L.head, removeAll(L.tail, x));
}
```

Aside: How to Write a Loop (in Theory)

- Try to give a description of how things look on *any arbitrary iteration* of the loop.
- This description is known as a *loop invariant*, because it is true from one iteration to the next.
- The loop body then must
 - Start from any situation consistent with the invariant;
 - Make progress in such a way as to make the invariant true again.

```

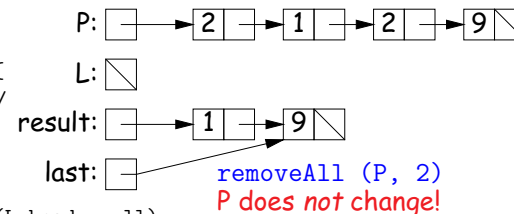
while (condition) {
    // Invariant true here
    loop body
    // Invariant again true here
}
// Invariant true and condition false.
            
```
- So if (*invariant* and not *condition*) is enough to insure we've got the answer, we're done!

Iterative Non-destructive List Deletion

Same as before, but use front-to-back iteration rather than recursion.

```

/** The list resulting from removing all instances of X from L
 * non-destructively. */
static IntList removeAll(IntList L, int x) {
    IntList result, last;
    result = last = null;
    for ( ; L != null; L = L.tail) {
        /* L != null and  $\mathcal{I}$  is true. */
        if (x == L.head)
            continue;
        else if (last == null)
            result = last = new IntList(L.head, null);
        else
            last = last.tail = new IntList(L.head, null);
    }
    return result;
}
            
```

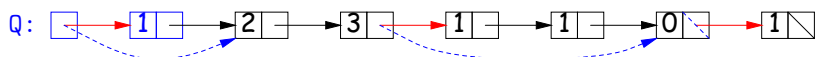


Here, \mathcal{I} is the *loop invariant*:

Result is all elements of L_0 not equal to x up to and not including L , and $last$ points to the last element of $result$, if any. We use L_0 here to mean "the original sequence of int values in L ."

Destructive Deletion

→ : Original : after $Q = \text{dremoveAll}(Q, 1)$



/** The list resulting from removing all instances of X from L.

* The original list may be destroyed. */

```

static IntList dremoveAll(IntList L, int x) {
    if (L == null)
        return null;
    else if (L.head == x)
        return dremoveAll(L.tail, x);
    else {
        L.tail = dremoveAll(L.tail, x);
        return L;
    }
}
            
```

Iterative Destructive Deletion

/** The list resulting from removing all instances of X from L.

* Original contents of L may be destroyed. */

```

static IntList dremoveAll(IntList L, int x) {
    IntList result, last;
    result = last = null;
    while (L != null) {
        IntList next = L.tail;
        if (x != L.head) {
            if (last == null)
                result = last = L;
            else
                last = last.tail = L;
            L.tail = null;
        }
        L = next;
    }
    return result;
}
            
```

