

## CS61A Lecture 31

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## Announcements

$\square$ HW9 due tonight
$\square$ Ants extra credit due tonight
$\square$ See Piazza for submission instructions
$\square$ Hog revisions out, due Monday

ㅁ HW10 out tonight

Scheme has built-in pairs that use weird names:

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair

A pair is represented by a dot between the elements, enclosed in parentheses

$$
\begin{aligned}
& >(\text { cons } 12) \\
& \text { (1 . 2) } \\
& \text { > (car (cons } 12 \text { )) } \\
& 1 \\
& >(c d r(c o n s 12)) \\
& 2
\end{aligned}
$$

## Recursive Lists

A recursive list can be represented as a pair in which the second element is a recursive list or the empty list

Scheme lists are recursive lists:

- nil is the empty list
- A non-empty Scheme list is a pair in which the second element is nil or a Scheme list

Scheme lists are written as space-separated combinations

```
> (define x (cons 1 (cons 2 (cons 3 (cons 4 nil)))))
> x
(1 2 % 3 4)
> (cdr x)
(2 3 4)
>(cons 1 (cons 2 (cons 34)))
(1 2 3 . 4)

\section*{Symbolic Programming}

Symbols are normally evaluated to produce values; how do we refer to symbols?
```

> (define a 1)
> (define b 2)
> (list a b)
(1 2)

```

Quotation prevents something from being evaluated by Lisp


Quotation can also be applied to combinations to form lists
```

>(car '(a b c))
a
>(cdr '(a b c))
(b c)

```

\section*{Scheme Lists and Quotation}

Dots can be used in a quoted list to specify the second element of the final pair
\[
>\left(c d r\left(c d r \quad\left(\begin{array}{lll}
1 & 2 & 3
\end{array}\right)\right)\right.
\]

However, dots appear in the output only of ill-formed lists
\[
\begin{aligned}
& \text { > '(1 2 . 3) } \\
& \text { (1 } 2.3 \text { ) } \\
& >\text { '(1 2 . (3 4)) } \\
& \text { (1 } 234 \text { 4) } \\
& \text { > '(1 } 2 \text { 3.nil) } \\
& \text { (1 } 23 \text { ) }
\end{aligned}
\]


What is the printed result of evaluating this expression?
\[
\begin{aligned}
& >(c d r \\
& \left(\begin{array}{lll}
3 & 4 & 5
\end{array}\right)
\end{aligned}
\]

\section*{The Let Special Form}

Let expressions introduce a new frame, with the given bindings
(let ((<name> <exp>) ...) <body>)

(define (filter fn s)
(if (null? s)
s
(let ((first (car s))
(rest (filter fn (cdr s))))
(if (fn first)
(cons first rest)
rest)))
> (filter even? '(1 234567\()\) )
(2 4 6)

\section*{Quick Sort}

Quick sort algorithm:
1. Choose a pivot (e.g. first element)
2. Partition into three pieces:
< pivot, = pivot, > pivot
3. Recurse on first and last piece
(define (filter-comp comp pivot s)
(filter (lambda (x) (comp x pivot)) s))
(define (quick-sort s)

(if (<= (length s) 1)
S
(let ((pivot (car s)))
(append (quick-sort (filter-comp < pivot s))
(filter-comp \(=\) pivot s)
(quick-sort (filter-comp > pivot s))))))

\section*{The Begin Special Form}

Begin expressions allow sequencing
\[
\text { (begin <exp }{ }_{1}><\exp _{2}>\ldots<\exp _{n}>\text { ) }
\]
(define (repeat k fn)
(if (> k 0)
(begin (fn) (repeat (-k 1) fn)) 'done))
(define (tri fn)
(repeat 3 (lambda () (fn) (lt 120))))
(define (sier d k)
(tri (lambda () (if (= k 1) (fd d) (leg d k))))
(define (leg d k)
(sier (/ d 2) (- k 1)) (penup) (fd d) (pendown))```

