## 61A Lecture 23

## Wednesday, October 30

## Announcements

## Announcements

- Homework 7 due Tuesday 11/5 @ 11:59pm.


## Announcements

- Homework 7 due Tuesday 11/5 @ 11:59pm.
- Project 1 composition revisions due Thursday 11/7 @ 11:59pm.


## Announcements

- Homework 7 due Tuesday 11/5 @ 11:59pm.
- Project 1 composition revisions due Thursday 11/7 @ 11:59pm.
- Midterm 2 is graded.


## Announcements

- Homework 7 due Tuesday 11/5 @ 11:59pm.
- Project 1 composition revisions due Thursday 11/7 @ 11:59pm.
- Midterm 2 is graded.
-(And yes, it was very challenging.)


## Announcements

- Homework 7 due Tuesday 11/5 @ 11:59pm.
- Project 1 composition revisions due Thursday 11/7 @ 11:59pm.
- Midterm 2 is graded.
-(And yes, it was very challenging.)
-Mean: 30


## Announcements

- Homework 7 due Tuesday 11/5 @ 11:59pm.
- Project 1 composition revisions due Thursday 11/7 @ 11:59pm.
- Midterm 2 is graded.
-(And yes, it was very challenging.)
-Mean: 30
-Solutions will be posted and exams distributed soon.

Scheme

Scheme is a Dialect of Lisp

Scheme is a Dialect of Lisp

What are people saying about Lisp?

Scheme is a Dialect of Lisp

What are people saying about Lisp?
-"The greatest single programming language ever designed."
-Alan Kay, co-inventor of Smalltalk and 00P

Scheme is a Dialect of Lisp

What are people saying about Lisp?
-"The greatest single programming language ever designed."
-Alan Kay, co-inventor of Smalltalk and 00P
-"The only computer language that is beautiful." -Neal Stephenson, DeNero's favorite sci-fi author

Scheme is a Dialect of Lisp

What are people saying about Lisp?
-"The greatest single programming language ever designed."
-Alan Kay, co-inventor of Smalltalk and 00P
-"The only computer language that is beautiful." -Neal Stephenson, DeNero's favorite sci-fi author
-"God's programming language."
-Brian Harvey, Berkeley CS instructor extraordinaire

Scheme is a Dialect of Lisp

What are people saying about Lisp?
-"The greatest single programming language ever designed."
-Alan Kay, co-inventor of Smalltalk and 00P
-"The only computer language that is beautiful."
-Neal Stephenson, DeNero's favorite sci-fi author
-"God's programming language."
-Brian Harvey, Berkeley CS instructor extraordinaire

http://imgs.xkcd.com/comics/lisp_cycles.png

Scheme Fundamentals

Scheme Fundamentals

Scheme programs consist of expressions, which can be:

Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,


## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 10 2), (not true)


## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 102 ), (not true)

Numbers are self-evaluating; symbols are bound to values.

## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 10 2), (not true)

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 10 2), (not true)

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

```
> (quotient 10 2)
```

5

## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 10 2), (not true)

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

```
> (quotient 10 2)
5
```

```
"quotient" names Scheme's
built-in integer division
procedure (i.e., function)
```


## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 10 2), (not true)

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

```
> (quotient 10 2)
5
> (quotient (+ 8 7) 5)
3
```

```
"quotient" names Scheme's
built-in integer division
procedure (i.e., function)
```


## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 102 ), (not true), ...

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

```
> (quotient 10 2)
5
> (quotient (+ 8 7) 5)
3
> (+ (* 3
    (+ (* 2 4)
                                    (+ 3 5)))
    (+ (- 10 7)
    6))
```

```
"quotient" names Scheme's
built-in integer division
procedure (i.e., function)
```


## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 102 ), (not true), ...

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

```
> (quotient 10 2)
5
> (quotient (+ 8 7) 5)
3
> (+ (* 3
    (+(* 2 4)
                                    (+ 3 5)))
    (+ (- 10 7)
    6))
```


## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 102 ), (not true), ...

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

```
> (quotient 10 2)
5
> (quotient (+ 8 7) 5)
3
> (+}(*)
    (+ (* 2 4)
                                    (+ 3 5)))
        (+ (- 10 7)
            6))
```


## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 102 ), (not true), ...

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

```
> (quotient 10 2)
5
> (quotient (+ 8 7) 5)
3
>(\Psi}(*)
    6))
```


## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 10 2), (not true), ...

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

```
> (quotient 10 2)
5
> (quotient (+ 8 7) 5)
3
> 㫙溇3
(+H(*)
                                    (+ 3 5)))
    (+ (- 10 7)
    6))
```


## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 10 2), (not true), ...

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.


```
"quotient" names Scheme's
built-in integer division
procedure (i.e., function)
```

Combinations can span multiple lines
(spacing doesn't matter)

## Scheme Fundamentals

Scheme programs consist of expressions, which can be:

- Primitive expressions: 2, 3.3, true, +, quotient,
- Combinations: (quotient 10 2), (not true), ...

Numbers are self-evaluating; symbols are bound to values.
Call expressions include an operator and 0 or more operands in parentheses.

```
> (quotient 10 2)
5
> (quotient (+ 8 7) 5)
3
> (H)
-(+}(*)
(+ 3 5)))
    (\square)(-10 7)
        6))
```

```
"quotient" names Scheme's
built-in integer division
procedure (i.e., function)
```


(Demo)

Special Forms

## Special Forms

## Special Forms

A combination that is not a call expression is a special form:

## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)


## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)

Evaluation:
(1) Evaluate the predicate expression.
(2) Evaluate either
the consequent or alternative.

## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)
- And and or: (and <e $\left.\left.e_{1}\right\rangle \ldots,<e_{n}\right\rangle$ ), (or $\left.\left.<e_{1}\right\rangle \ldots,<e_{n}\right\rangle$ )

Evaluation:
(1) Evaluate the predicate expression.
(2) Evaluate either
the consequent or alternative.

## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)
- And and or: (and <e $e_{1}>\ldots<e_{n}>$ ), (or $<e_{1}>\ldots<e_{n}>$ )
- Binding symbols: (define <symbol> <expression>)


## Evaluation:

(1) Evaluate the predicate expression.
(2) Evaluate either
the consequent or alternative.

## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)
- And and or: (and <e $\left.e_{1}>\ldots<e_{n}\right\rangle$ ), (or $\left.<e_{1}>\ldots .<e_{n}\right\rangle$ )
- Binding symbols: (define <symbol> <expression>)


## Evaluation:

(1) Evaluate the predicate expression.
(2) Evaluate either
the consequent or alternative.

```
> (define pi 3.14)
> (* pi 2)
6.28
```


## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)
- And and or: (and <e $e_{1}>\ldots<e_{n}>$ ), (or $<e_{1}>\ldots<e_{n}>$ )
- Binding symbols: (define <symbol> <expression>)


## Evaluation:

(1) Evaluate the predicate expression.
(2) Evaluate either
the consequent or alternative.

```
> (define pi 3.14)
6.28
```


## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)
- And and or: (and <e $e_{1}>\ldots<e_{n}>$ ), (or <e $e_{1}>\ldots<e_{n}>$ )
- Binding symbols: (define <symbol> <expression>)
- New procedures: (define (<symbol> <formal parameters>) <body>)


## Evaluation:

(1) Evaluate the predicate expression.
(2) Evaluate either
the consequent or alternative.

$$
\begin{aligned}
& >(\text { define pi } 3.14) \\
& >(* \text { pi } 2) \\
& 6.28
\end{aligned}
$$

## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)
- And and or: (and <e $e_{1}>\ldots<e_{n}>$ ), (or $<e_{1}>\ldots<e_{n}>$ )
- Binding symbols: (define <symbol> <expression>)
- New procedures: (define (<symbol> <formal parameters>) <body>)


## Evaluation:

(1) Evaluate the predicate expression.
(2) Evaluate either
the consequent or alternative.

```
>>(define pi 3.14)}\begin{array}{l}{\mathrm{ ( (* pi 2)}}\\{6.28}
> (define (abs x)
        (if (< x 0)
            (- x)
            x))
> (abs -3)
3
```


## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)
- And and or: (and <e $e_{1}>\ldots<e_{n}>$ ), (or $<e_{1}>\ldots<e_{n}>$ )
- Binding symbols: (define <symbol> <expression>)
- New procedures: (define (<symbol> <formal parameters>) <body>)


## Evaluation:

(1) Evaluate the predicate expression.
(2) Evaluate either
the consequent or alternative.

```
>> (define pi 3.14)}\begin{array}{l}{\mathrm{ > (* pi 2)}}\\{6.28}
> (define (abs x)
        (if (< x 0)
            (-x)
                            A procedure is created and bound to the
                                symbol "abs"
>(abs -3)
3
```


## Special Forms

A combination that is not a call expression is a special form:

- If expression: (if <predicate> <consequent> <alternative>)
- And and or: (and $\left.<e_{1}\right\rangle \ldots,<e_{n}>$ ), (or $\left.<e_{1}>\ldots \ldots<e_{n}\right\rangle$ )
- Binding symbols: (define <symbol> <expression>)
- New procedures: (define (<symbol> <formal parameters>) <body>)


## Evaluation:

(1) Evaluate the predicate expression.
(2) Evaluate either
the consequent or alternative.

```
>>(define pi 3.14)}\begin{array}{l}{\mathrm{ > (* pi 2)}}\\{6.28}
> (define (abs x)
        (if (< x 0)
            (-x)
A procedure is created and bound to the
            (-x
                            A pymbol "abs"
> (abs -3)
3
(Demo)
```

Counting Trees

## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.

## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.
a long noun phrase

## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.

## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent. a long noun phrase
a two word modifier

## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.

some trees are balanced

## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


The number of trees over $n$ leaves with $k$ leaves in the left and $n-k$ in the right is:

## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


The number of trees over $n$ leaves with $k$ leaves in the left and $n-k$ in the right is:
(The number of trees with $\mathbf{k}$ leaves) $*$ (The number of trees with $\mathbf{n}-\mathbf{k}$ leaves)

## Example: Counting Binary Trees

The structure of a sentence can be described by a tree. Each sub-tree is a constituent.


The number of trees over $n$ leaves with $k$ leaves in the left and $n-k$ in the right is: (The number of trees with $\mathbf{k}$ leaves) $*$ (The number of trees with $\mathbf{n}-\mathbf{k}$ leaves)

Lambda Expressions

## Lambda Expressions

Lambda expressions evaluate to anonymous procedures.

## Lambda Expressions

Lambda expressions evaluate to anonymous procedures.
(lambda (<formal-parameters>) <body>)

## Lambda Expressions

Lambda expressions evaluate to anonymous procedures.


## Lambda Expressions

Lambda expressions evaluate to anonymous procedures.


## Lambda Expressions

Lambda expressions evaluate to anonymous procedures.


An operator can be a call expression too:

## Lambda Expressions

Lambda expressions evaluate to anonymous procedures.


An operator can be a call expression too:
((lambda (x y z) (+ x y (square z))) 12 3)

## Lambda Expressions

Lambda expressions evaluate to anonymous procedures.


An operator can be a call expression too:


Evaluates to the
add $-\boldsymbol{x}-\delta-\boldsymbol{y}-\delta-\boldsymbol{z}^{2}$ procedure

Pairs and Lists

Pairs and Lists

## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!
$>($ define $x$ (cons 12 ) $)$


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!
$>$ (define $x$ (cons 12 ))
$>\mathrm{X}$


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!
$>$ (define $x$ (cons 12 ))
$>\mathrm{x}$
(1.2)


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!
$>$ (define $x$ (cons 12 ))
$>\mathrm{x}$
(1 . 2)
> (car x)


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!
$>$ (define $x$ (cons 12 ))
$>\mathrm{x}$
(1 . 2)
> ( $\operatorname{car} \mathrm{x}$ )
1


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!
$>$ (define $x$ (cons 12 ))
$>x$
(1 . 2)
$>(\operatorname{car} x)$
1
$>(\operatorname{cdr} x)$


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!

```
    > (define x (cons 1 2))
```

    \(>\mathrm{x}\)
    (1 . 2)
    > ( \(\operatorname{car} \mathrm{x}\) )
    1
    $>(\operatorname{cdr} x)$
2

## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!

```
    > (define x (cons 1 2))
```

    \(>x\)
    (1 . 2)
    > ( \(\operatorname{car} \mathrm{x}\) )
    1
    $>(\operatorname{cdr} x)$
2
> (cons 1 (cons 2 (cons 3 (cons 4 nil))))

## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!

```
> (define x (cons 1 2))
```

$>x$
(1 . 2)
> ( $\operatorname{car} \mathrm{x}$ )
1
$>(\operatorname{cdr} x)$
2
> (cons 1 (cons 2 (cons 3 (cons 4 nil))))
(1 234 )

## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!
$>$ (define $x$ (cons 12 ))
> X
(1 . 2)
$>(\operatorname{car} \mathrm{x})$
Not a well-formed list!
1
$>(\operatorname{cdr} x)$
2
> (cons 1 (cons 2 (cons 3 (cons 4 nil))))
(1 234 )


## Pairs and Lists

In the late 1950s, computer scientists used confusing names.

- cons: Two-argument procedure that creates a pair
- car: Procedure that returns the first element of a pair
- cdr: Procedure that returns the second element of a pair
- nil: The empty list

They also used a non-obvious notation for recursive lists.

- A (recursive) list in Scheme is a pair in which the second element is nil or a Scheme list.
- Scheme lists are written as space-separated combinations.
- A dotted list has any value for the second element of the last pair; maybe not a list!

```
> (define x (cons 1 2))
```

> X
(1 . 2)
$>(\operatorname{car} x)$
Not a well-formed list!
1
$>(\operatorname{cdr} x)$
2
> (cons 1 (cons 2 (cons 3 (cons 4 nil))))
(1 234 )

Symbolic Programming

## Symbolic Programming

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols? > (define a 1)

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?
> (define a 1)
> (define b 2)

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?
> (define a 1)
$>$ (define b 2)
> (list a b)

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?
> (define a 1)
$>$ (define b 2)
> (list a b)
(1 2)

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?
> (define a 1)
$>$ (define b 2)
No sign of "a" and "b" in the resulting value
(1 2)

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.
> (list 'a 'b)

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.

```
> (list 'a 'b)
(a b)
```


## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.

```
> (list 'a 'b)
(a b)
> (list 'a b)
```


## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.

```
> (list 'a 'b)
(a b)
> (list 'a b)
(a 2)
```


## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.

(a 2)

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.
a b)
(a 2)

Quotation can also be applied to combinations to form lists.

## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.
> (list 'a b)
(a 2)

Quotation can also be applied to combinations to form lists.

```
> (car '(a b c))
```


## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.
> (list 'a b)
(a 2)

Quotation can also be applied to combinations to form lists.

```
> (car '(a b c))
a
```


## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.
> (list 'a b)
(a 2)

Quotation can also be applied to combinations to form lists.

```
> (car '(a b c))
a
> (cdr '(a b c))
```


## Symbolic Programming

Symbols normally refer to values; how do we refer to symbols?


Quotation is used to refer to symbols directly in Lisp.
> (list 'a b)
(a 2)

Quotation can also be applied to combinations to form lists.

```
> (car '(a b c))
a
>(cdr '(a b c))
(b c)
```

Scheme Lists and Quotation

## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

$$
>(c d r(c d r \quad(12.3)))
$$

## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.
$>(c d r(c d r '(12.3)))$
3

## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.
$>(\operatorname{cdr}(\operatorname{cdr} \quad(12,3)))$
3
However, dots appear in the output only of ill-formed lists.

## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.
$>(\operatorname{cdr}(\operatorname{cdr} \quad(12,3)))$
3
However, dots appear in the output only of ill-formed lists.
$>{ }^{\prime}\left(\begin{array}{ll}1 & 2\end{array}\right)$

## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.
$>(\operatorname{cdr}(\operatorname{cdr} \quad(12,3)))$
3
However, dots appear in the output only of ill-formed lists.
$>$ '(1 2.3 )


## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 , 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2, 3)
```



## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2 . 3)
> '(1 2 . (3 4))
```


## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2 . 3)
> '(1 2 . (3 4))
```



## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2 . 3)
> '(1 2 . (3 4))
```



## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

|  |
| :---: |
|  |  |
|  |  |
|  |  |



## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2 . 3)
>'(1 2 . (3 4))
(1 2 3 4)
> '(1 2 3 . nil)
```



## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2 . 3)
> '(1 2 . (3 4))
(1 2 3 4)
> '(1 2 3 . nil)
```



## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2 . 3)
> (1 2 . (3 4))
(1 2 3 4)
> '(1 2 3 . nil)
(1 2 3)
```



## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2 . 3)
> '(1 2 . (3 4))
(1 2 3 4)
> '(1 2 3 . nil)
(1 2 3)
```



What is the printed result of evaluating this expression?

## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2 . 3)
> '(1 2 . (3 4))
(1 2 3 4)
> '(1 2 3 . nil)
(1 2 3)
(1 2 3)
```



What is the printed result of evaluating this expression?
$>(c d r \quad((12) \cdot(34 .(5))))$

## Scheme Lists and Quotation

Dots can be used in a quoted list to specify the second element of the final pair.

```
> (cdr (cdr '(1 2 . 3)))
3
```

However, dots appear in the output only of ill-formed lists.

```
> '(1 2 . 3)
(1 2 . 3)
> '(1 2 . (3 4))
(1 2 3 4)
> '(1 2 3 . nil)
(1 2 3)
(1 2 3)
```



What is the printed result of evaluating this expression?

```
>(cdr '((1 2) . (3 4 . (5))))
(3 4 5)
```

Sierpinski's Triangle

