61A Lecture 19

Friday, October 18

## Announcements

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"Please let us know you are coming by filling out the Piazza poll

Comparing Orders of Growth

Comparing orders of growth ( n is the problem size)

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| $\Theta\left(b^{n}\right)$ | Exponential growth! Recursive fib takes <br> $\Theta\left(n^{6}\right) \ldots\left(\phi^{n}\right)$ steps, where $\phi=\frac{1+\sqrt{5}}{2} \approx 1.61828$ <br> $\Theta\left(n^{2}\right)$ |
| ---: | :--- |
| $\Theta(n)$ | Incrementing the problem scales $\mathrm{R}(\mathrm{n})$ by a factor. <br> Quadratic growth. E.g., operations on all pairs. <br> Incrementing n increases $\mathrm{R}(\mathrm{n})$ by the problem size n. |
| $\Theta(\log n)$ | Logarithmic growth. These processes scale well. |
| $\Theta(1)$ | Doubling the problem only increments $\mathrm{R}(\mathrm{n})$. |

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Sets

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## Sets

One more built-in Python container type

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>>> s = {3, 2, 1, 4, 4}
>>> s
{1, 2, 3, 4}
>>> 3 in s
True
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>>> s = {3, 2, 1, 4, 4}
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>>> s.union({1, 5})
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{1, 2, 3, 4, 5}
>>> s.intersection({6, 5, 4, 3})
{3, 4}
```

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def empty(s):
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    elif s.first == v:
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which means that there are positive constants $k_{1}$ and $k_{2}$ such that

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for sufficiently large values of $\boldsymbol{n}$.

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        return s
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def adjoin_set(s, v):
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    else:
        return Rlist(v, s)
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## Sets as Unordered Sequences

Time order of growth

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The size of the set

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\Theta(n' 2
Assume sets are
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def union_set(set1, set2):
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    the same size
```

def union_set(set1, set2):
not_in_set2 = lambda v: not set_contains(set2, v)
$\Theta\left(n^{2}\right)$
set1_not_set2 = filter_rlist(set1, not_in_set2)
return extend_rlist(set1_not_set2, set2)
(Demo)

## Sets as Ordered Sequences

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Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

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```
def set_contains(s, v):
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Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

```
def set_contains(s, v):
    if empty(s) or s.first > v:
            return False
```


## Sets as Ordered Sequences

Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

```
def set_contains(s, v):
    if empty(s) or s.first > v:
            return False
    elif s.first == v:
        return True
```


## Sets as Ordered Sequences

Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

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def set_contains(s, v):
    if empty(s) or s.first > v:
            return False
    elif s.first == v:
        return True
    else:
```


## Sets as Ordered Sequences

Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

```
def set_contains(s, v):
    if empty(s) or s.first > v:
            return False
    elif s.first == v:
        return True
    else:
            return set_contains(s.rest, v)
```


## Sets as Ordered Sequences

Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

```
def set_contains(s, v):
    if empty(s) or s.first > v:
            return False
    elif s.first == v:
        return True
    else:
            return set_contains(s.rest, v)
```

Order of growth?

## Sets as Ordered Sequences

Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

```
def set_contains(s, v):
    if empty(s) or s.first > v:
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    else:
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```

                    Order of growth? \(\Theta(n)\)
    
## Set Intersection Using Ordered Sequences

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This algorithm assumes that elements are in order.

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```
    def intersect_set(set1, set2):
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```
    def intersect_set(set1, set2):
        if empty(set1) or empty(set2):
```


## Set Intersection Using Ordered Sequences

This algorithm assumes that elements are in order.

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
```


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```
def intersect_set(set1, set2):
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```
def intersect_set(set1, set2):
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    else:
        e1, e2 = set1.first, set2.first
```


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```


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        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
```


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        e1, e2 = set1.first, set2.first
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            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
```


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        elif e1 < e2:
            return intersect_set(set1.rest, set2)
```


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        elif e2 < e1:
```


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            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
            return intersect_set(set1.rest, set2)
        elif e2 < e1:
            return intersect_set(set1, set2.rest)
```


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    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
            return intersect_set(set1.rest, set2)
        elif e2 < el:
            return intersect_set(set1, set2.rest)
                (Demo)
```


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This algorithm assumes that elements are in order.

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    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
            return intersect_set(set1.rest, set2)
        elif e2 < e1:
            return intersect_set(set1, set2.rest)
                    (Demo)
                    Order of growth?
```


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This algorithm assumes that elements are in order.

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
            return intersect_set(set1.rest, set2)
        elif e2 < e1:
            return intersect_set(set1, set2.rest)
                    (Demo)
                    Order of growth? }\Theta(n
```


## Sets as Binary Search Trees

Tree Sets

## Tree Sets

Proposal 3: A set is represented as a Tree. Each entry is:

## Tree Sets

Proposal 3: A set is represented as a Tree. Each entry is:

- Larger than all entries in its left branch and


## Tree Sets

Proposal 3: A set is represented as a Tree. Each entry is:

- Larger than all entries in its left branch and
- Smaller than all entries in its right branch


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11

## Tree Sets

Proposal 3: A set is represented as a Tree. Each entry is:

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- Smaller than all entries in its right branch


11

Membership in Tree Sets

## Membership in Tree Sets

Set membership traverses the tree

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- The element is either in the left or right sub-branch


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- By focusing on one branch, we reduce the set by about half


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def set_contains(s, v):
```


## Membership in Tree Sets

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- The element is either in the left or right sub-branch
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```
def set_contains(s, v):
        if s is None:
```


## Membership in Tree Sets

Set membership traverses the tree

- The element is either in the left or right sub-branch
- By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
        if s is None:
            return False
```


## Membership in Tree Sets

Set membership traverses the tree

- The element is either in the left or right sub-branch
- By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
        if s is None:
        return False
    elif s.entry == v:
```


## Membership in Tree Sets

Set membership traverses the tree

- The element is either in the left or right sub-branch
- By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
        if s is None:
        return False
    elif s.entry == v:
        return True
```


## Membership in Tree Sets

Set membership traverses the tree

- The element is either in the left or right sub-branch
- By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
        if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
```


## Membership in Tree Sets

Set membership traverses the tree

- The element is either in the left or right sub-branch
- By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
        if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
```


## Membership in Tree Sets

Set membership traverses the tree

- The element is either in the left or right sub-branch
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```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
```


## Membership in Tree Sets

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```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
        return set_contains(s.left, v)
```


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Set membership traverses the tree

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```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
        return set_contains(s.left, v)
```



## Membership in Tree Sets

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- The element is either in the left or right sub-branch
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```
def set_contains(s, v):
        if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
        return set_contains(s.left, v)
```



## Membership in Tree Sets

Set membership traverses the tree

- The element is either in the left or right sub-branch
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```
```

def set_contains(s, v):

```
```

def set_contains(s, v):
if s is None:
if s is None:
return False
return False
elif s.entry == v:
elif s.entry == v:
return True
return True
elif s.entry < v:
elif s.entry < v:
return set_contains(s.right, v)
return set_contains(s.right, v)
elif s.entry > v:
elif s.entry > v:
return set_contains(s.left, v)

```
```

        return set_contains(s.left, v)
    ```
```



```
If 9 is in the
set, it is in
    this branch
```


## Membership in Tree Sets

Set membership traverses the tree

- The element is either in the left or right sub-branch
- By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
        return set_contains(s.left, v)
```



```
If 9 is in the
set, it is in
    this branch
```

Order of growth?

Adjoining to a Tree Set

Adjoining to a Tree Set


Adjoining to a Tree Set


Adjoining to a Tree Set


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Adjoining to a Tree Set


Adjoining to a Tree Set


Adjoining to a Tree Set


Adjoining to a Tree Set


Adjoining to a Tree Set


Adjoining to a Tree Set


Left!
Right!
$\square$

None

None None

Stop!


Adjoining to a Tree Set


8
8
None

Stop!


More Set Operations

What Did I Leave Out?

## What Did I Leave Out?

Sets as ordered sequences:

## What Did I Leave Out?

Sets as ordered sequences:

- Adjoining an element to a set


## What Did I Leave Out?

Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets


## What Did I Leave Out?

Sets as ordered sequences:

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Sets as binary trees:

## What Did I Leave Out?

Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets

Sets as binary trees:

- Intersection of two sets


## What Did I Leave Out?

Sets as ordered sequences:
-Adjoining an element to a set
-Union of two sets
Sets as binary trees:

- Intersection of two sets
-Union of two sets


## What Did I Leave Out?

Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets

Sets as binary trees:

- Intersection of two sets
- Union of two sets
- Balancing a tree


## What Did I Leave Out?

Sets as ordered sequences:
-Adjoining an element to a set
-Union of two sets
Sets as binary trees:

- Intersection of two sets
- Union of two sets
- Balancing a tree

That's all on homework 7!

