61A Lecture 19

Friday, October 18

•Homework 6 is due Tuesday 10/22 @ 11:59pm

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 - •Includes a mid-semester survey about the course so far

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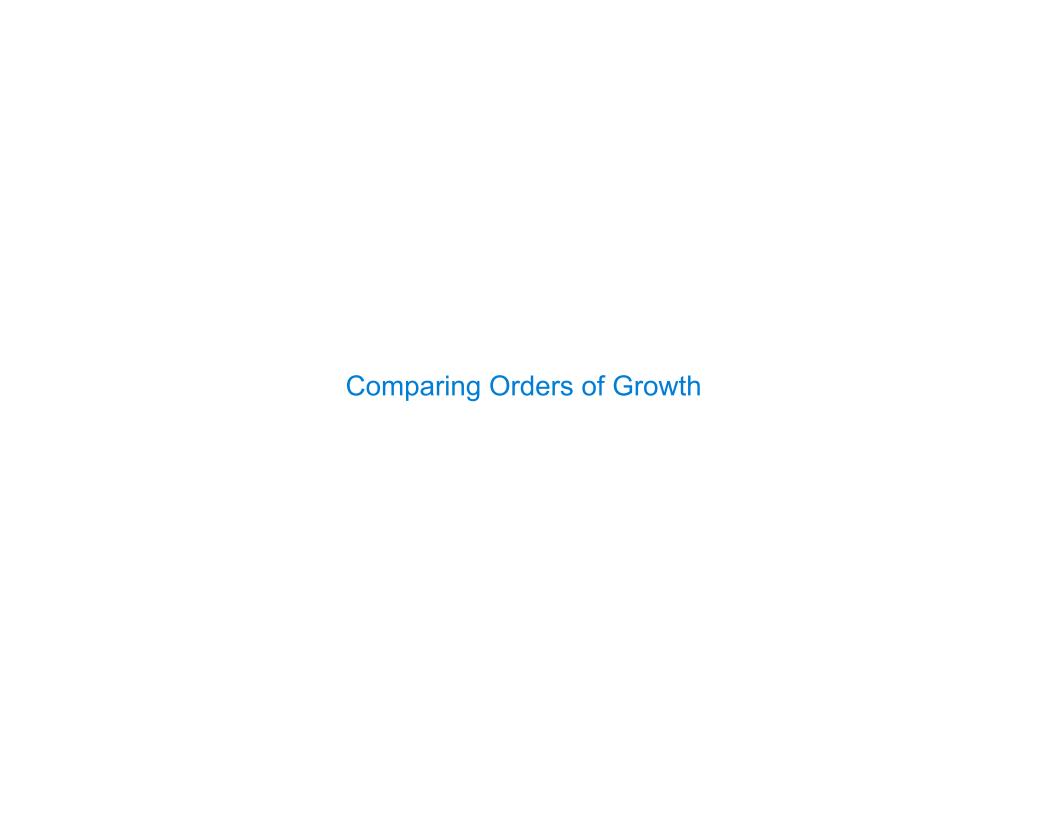
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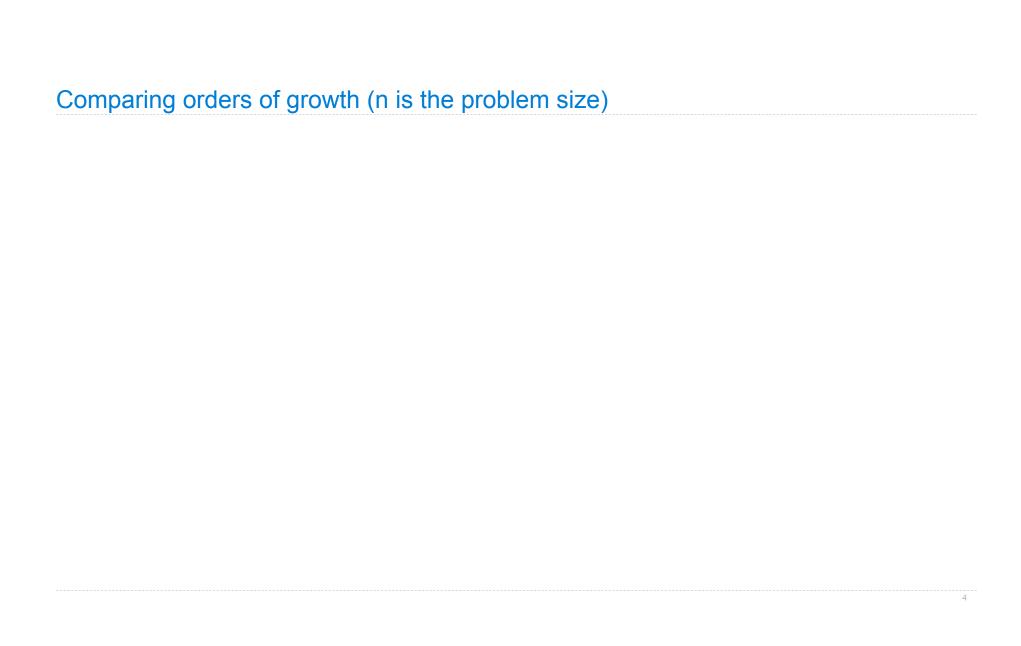
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 - Please let us know you are coming by filling out the Piazza poll





 $\Theta(b^n)$

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 Exponential growth! Recursive fib takes
$$\Theta(\phi^n) \text{ steps, where } \phi = \frac{1+\sqrt{5}}{2} \approx 1.61828$$

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Incrementing n increases R(n) by the problem size n.

 $\Theta(n)$ | Linear growth. Resources scale with the problem.

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Sets

One more built-in Python container type

Set literals are enclosed in braces

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>>> s = {3, 2, 1, 4, 4}
>>> s
{1, 2, 3, 4}
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>>> 3 in s
True
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6

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>>> 3 in s
True
>>> len(s)
4
>>> s.union({1, 5})
{1, 2, 3, 4, 5}
```

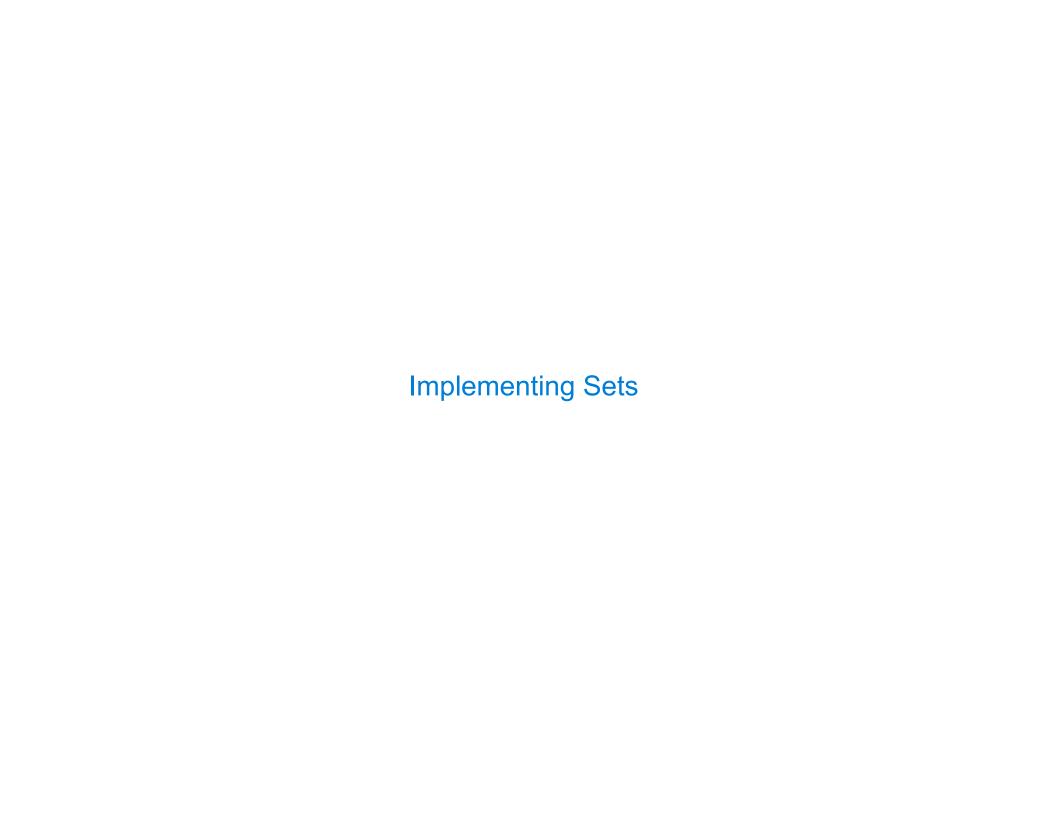
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>>> s
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>>> 3 in s
True
>>> len(s)
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>>> s.union({1, 5})
{1, 2, 3, 4, 5}
>>> s.intersection({6, 5, 4, 3})
{3, 4}
```



|--|

What we should be able to do with a set:

•Membership testing: Is a value an element of a set?

8

What we should be able to do with a set:

- •Membership testing: Is a value an element of a set?
- •Union: Return a set with all elements in set1 or set2

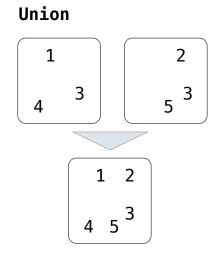
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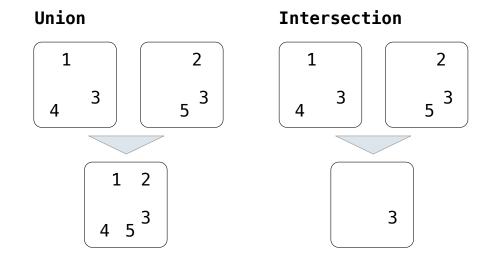
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Union 1 2 4 3 5 1 2 4 5

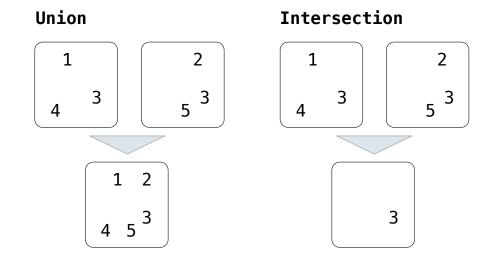
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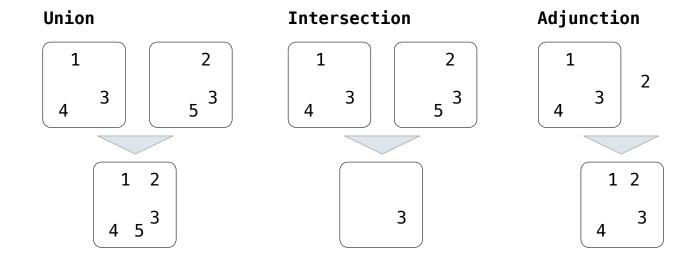
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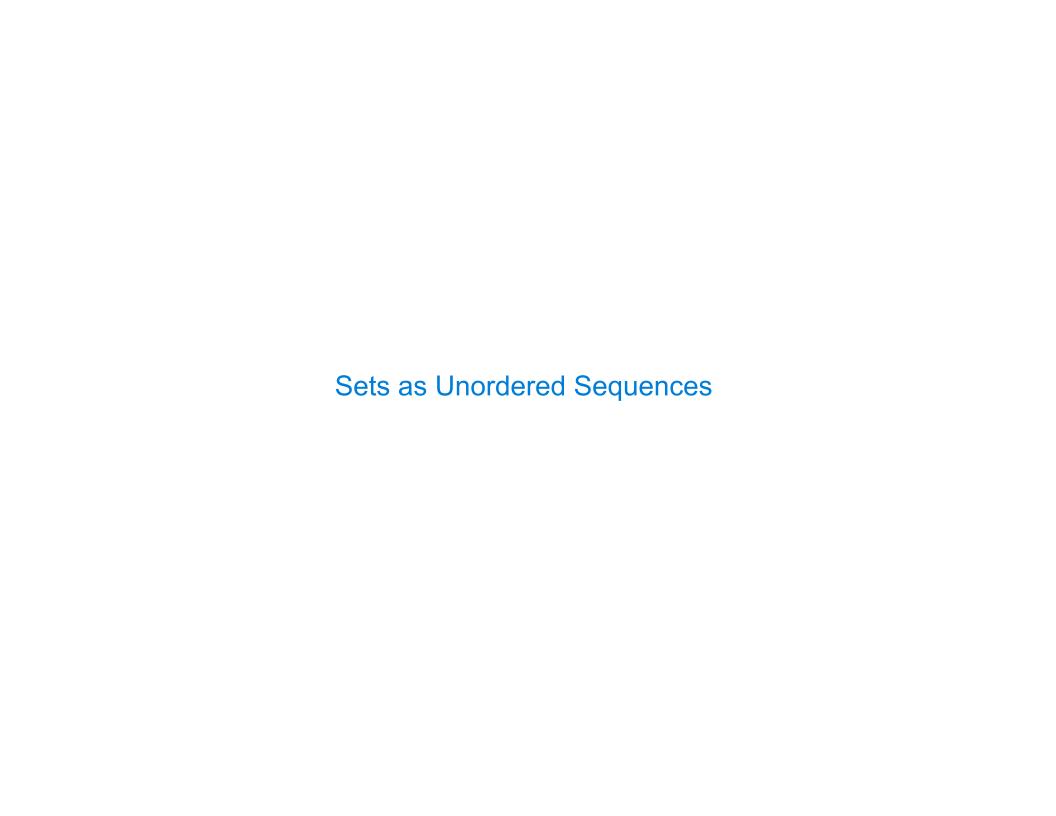


- •Membership testing: Is a value an element of a set?
- •Union: Return a set with all elements in set1 or set2
- •Intersection: Return a set with any elements in set1 and set2
- •Adjunction: Return a set with all elements in s and a value v



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def empty(s):
    return s is Rlist.empty
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def set_contains(s, v):
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def set_contains(s, v):
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def empty(s):
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def set_contains(s, v):
    if empty(s):
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def empty(s):
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def set_contains(s, v):
    if empty(s):
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    elif s.first == v:
        return True
    else:
        return set_contains(s.rest, v)
```

Proposal 1: A set is represented by a recursive list that contains no duplicate items.

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def set_contains(s, v):
    if empty(s):
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(Demo)

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which means that there are positive constants k_1 and k_2 such that

$$k_1 \cdot n \le R(n) \le k_2 \cdot n$$

for sufficiently large values of n.

Sets as Unordered Sequences

```
def adjoin_set(s, v):
```

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```
def adjoin_set(s, v):
    if set_contains(s, v):
```

```
def adjoin_set(s, v):
    if set_contains(s, v):
        return s
```

```
def adjoin_set(s, v):
    if set_contains(s, v):
       return s
    else:
```

```
def adjoin_set(s, v):
    if set_contains(s, v):
        return s
    else:
        return Rlist(v, s)
```

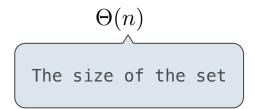
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Time order of growth

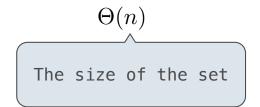
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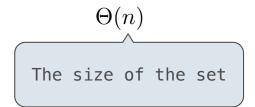
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def adjoin_set(s, v):
    if set_contains(s, v):
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def intersect_set(set1, set2):
```



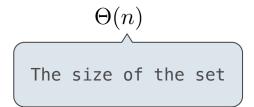
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def adjoin_set(s, v):
    if set_contains(s, v):
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        return Rlist(v, s)

def intersect_set(set1, set2):
    in_set2 = lambda v: set_contains(set2, v)
```



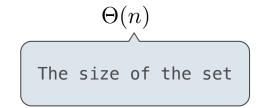
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def adjoin_set(s, v):
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def intersect_set(set1, set2):
    in_set2 = lambda v: set_contains(set2, v)
    return filter_rlist(set1, in_set2)
```



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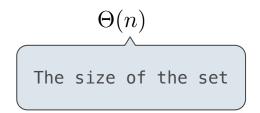
def intersect_set(set1, set2):
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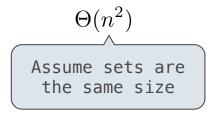


$$\Theta(n^2)$$

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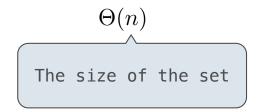


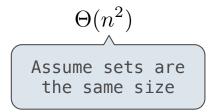


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    return filter_rlist(set1, in_set2)

def union_set(set1, set2):
```

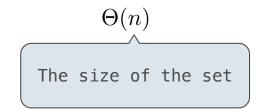


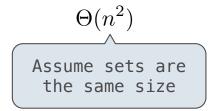


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def adjoin_set(s, v):
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def intersect_set(set1, set2):
    in_set2 = lambda v: set_contains(set2, v)
    return filter_rlist(set1, in_set2)

def union_set(set1, set2):
    not_in_set2 = lambda v: not set_contains(set2, v)
```

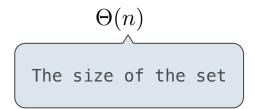


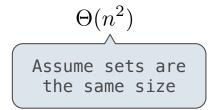


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def intersect_set(set1, set2):
    in_set2 = lambda v: set_contains(set2, v)
    return filter_rlist(set1, in_set2)

def union_set(set1, set2):
    not_in_set2 = lambda v: not set_contains(set2, v)
    set1_not_set2 = filter_rlist(set1, not_in_set2)
```

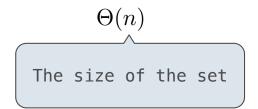


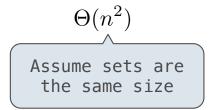


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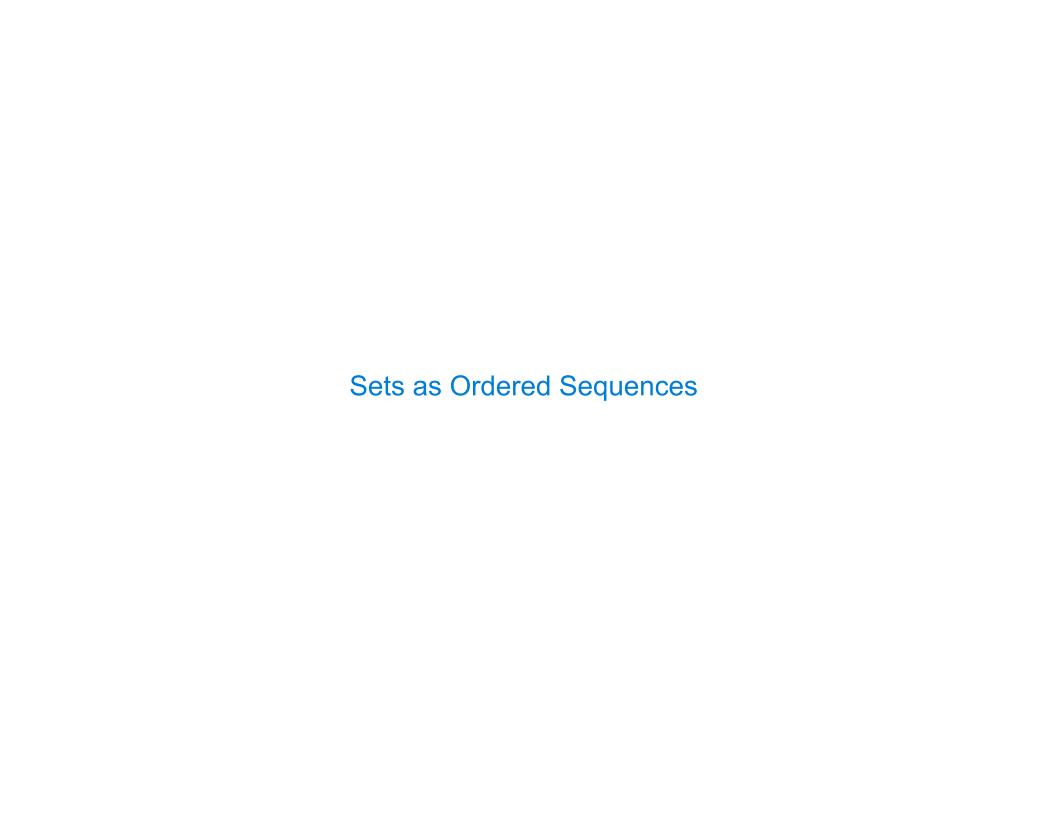
def union_set(set1, set2):
    not_in_set2 = lambda v: not set_contains(set2, v)
    set1_not_set2 = filter_rlist(set1, not_in_set2)
    return extend_rlist(set1_not_set2, set2)
```





```
\Theta(n)
def adjoin set(s, v):
    if set contains(s, v):
        return s
                                                              The size of the set
    else:
        return Rlist(v, s)
                                                                     \Theta(n^2)
def intersect set(set1, set2):
    in set2 = lambda v: set contains(set2, v)
    return filter rlist(set1, in set2)
                                                                Assume sets are
                                                                  the same size
def union set(set1, set2):
                                                                     \Theta(n^2)
    not in set2 = lambda v: not set contains(set2, v)
    set1 not set2 = filter rlist(set1, not in set2)
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    not in set2 = lambda v: not set contains(set2, v)
    set1 not set2 = filter rlist(set1, not in set2)
    return extend rlist(set1 not set2, set2)
                                          (Demo)
```



Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

def set_contains(s, v):

```
def set_contains(s, v):
    if empty(s) or s.first > v:
        return False
```

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def set_contains(s, v):
    if empty(s) or s.first > v:
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    elif s.first == v:
        return True
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    else:
```

```
def set_contains(s, v):
    if empty(s) or s.first > v:
        return False
    elif s.first == v:
        return True
    else:
        return set_contains(s.rest, v)
```

Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

```
def set_contains(s, v):
    if empty(s) or s.first > v:
        return False
    elif s.first == v:
        return True
    else:
        return set_contains(s.rest, v)
```

Order of growth?

Proposal 2: A set is represented by a recursive list with unique elements ordered from least to greatest

```
def set_contains(s, v):
    if empty(s) or s.first > v:
        return False
    elif s.first == v:
        return True
    else:
        return set contains(s.rest, v)
```

Order of growth? $\Theta(n)$

```
def intersect_set(set1, set2):
```

This algorithm assumes that elements are in order.

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
```

1

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
```

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
```

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
```

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
```

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
```

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:</pre>
```

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
            return intersect set(set1.rest, set2)</pre>
```

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
            return intersect_set(set1.rest, set2)
        elif e2 < e1:</pre>
```

This algorithm assumes that elements are in order.

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
            return intersect_set(set1.rest, set2)
        elif e2 < e1:
            return intersect set(set1, set2.rest)</pre>
```

This algorithm assumes that elements are in order.

```
def intersect_set(set1, set2):
    if empty(set1) or empty(set2):
        return Rlist.empty
    else:
        e1, e2 = set1.first, set2.first
        if e1 == e2:
            return Rlist(e1, intersect_set(set1.rest, set2.rest))
        elif e1 < e2:
            return intersect_set(set1.rest, set2)
        elif e2 < e1:
            return intersect_set(set1, set2.rest)</pre>
```

This algorithm assumes that elements are in order.

Order of growth?

This algorithm assumes that elements are in order.

Sets as Binary Search Trees

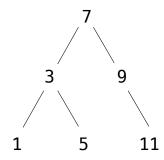
_		
Iroo	Sets	

Proposal 3: A set is represented as a Tree. Each entry is:

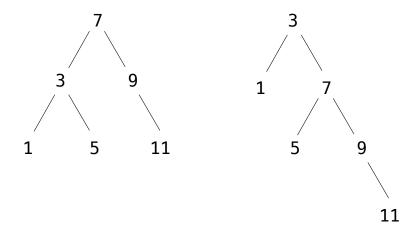
Larger than all entries in its left branch and

- Larger than all entries in its left branch and
- Smaller than all entries in its right branch

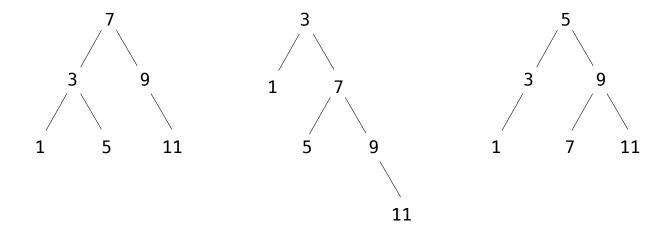
- Larger than all entries in its left branch and
- Smaller than all entries in its right branch



- Larger than all entries in its left branch and
- Smaller than all entries in its right branch



- Larger than all entries in its left branch and
- Smaller than all entries in its right branch



Set membership traverses the tree

• The element is either in the left or right sub-branch

- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

- The element is either in the left or right sub-branch
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```
def set_contains(s, v):
```

- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
    if s is None:
```

- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
    if s is None:
        return False
```

- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
```

- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
```

- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:</pre>
```

- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)</pre>
```

- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

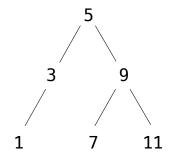
```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
```

- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

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def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
        return set_contains(s.left, v)
```

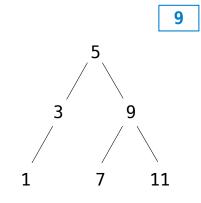
- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
        return set contains(s.left, v)
```



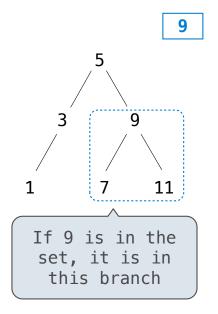
- The element is either in the left or right sub-branch
- •By focusing on one branch, we reduce the set by about half

```
def set_contains(s, v):
    if s is None:
        return False
    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
        return set contains(s.left, v)
```



- The element is either in the left or right sub-branch
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```
def set_contains(s, v):
    if s is None:
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    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
        return set contains(s.left, v)
```



Set membership traverses the tree

- The element is either in the left or right sub-branch
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```
def set_contains(s, v):
    if s is None:
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    elif s.entry == v:
        return True
    elif s.entry < v:
        return set_contains(s.right, v)
    elif s.entry > v:
        return set_contains(s.left, v)
```

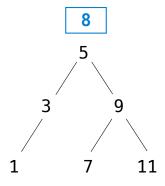
3 9

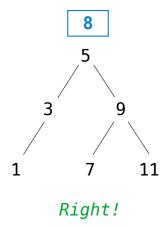
1 7 11

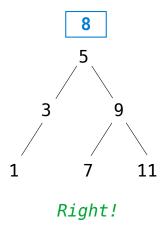
If 9 is in the set, it is in this branch

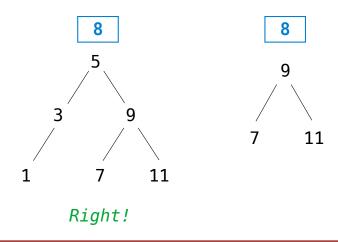
Order of growth?

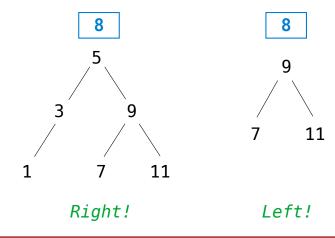
|--|

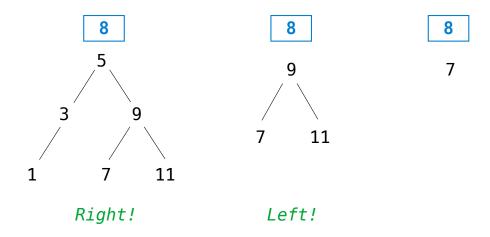


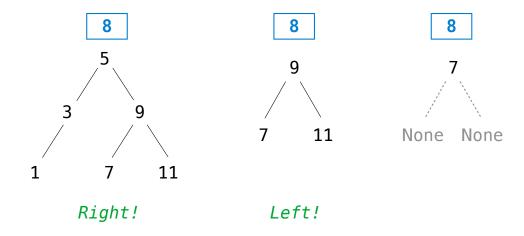


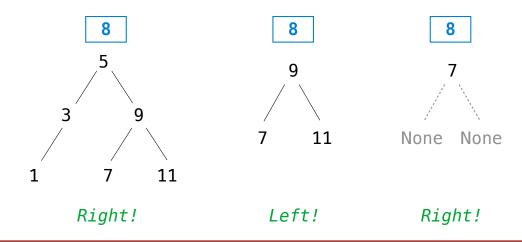


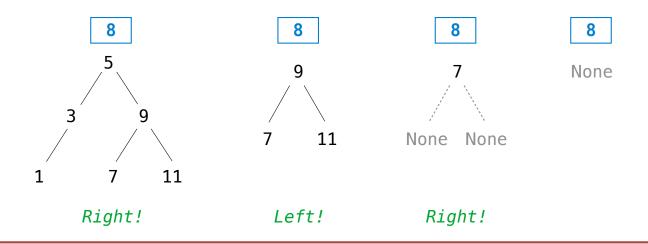


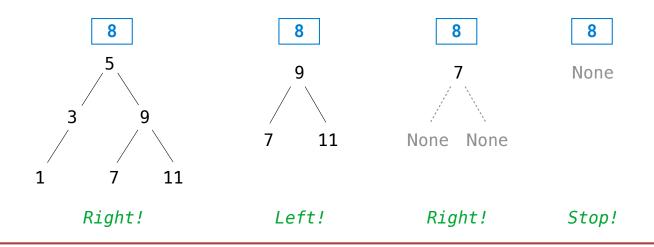


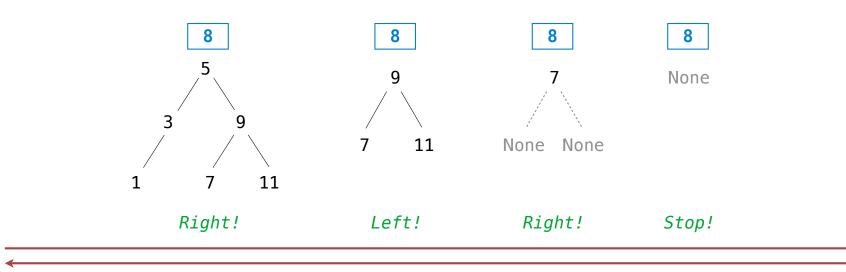


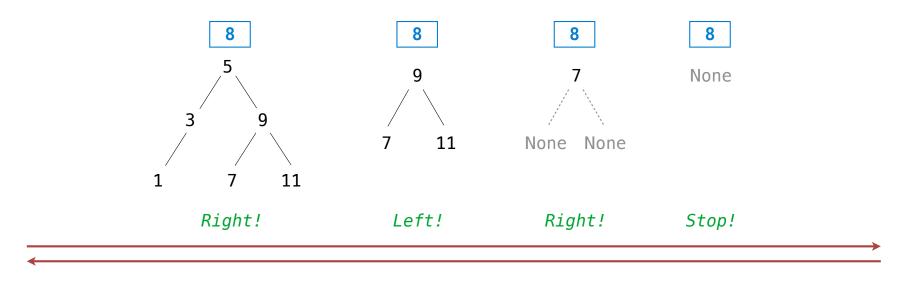




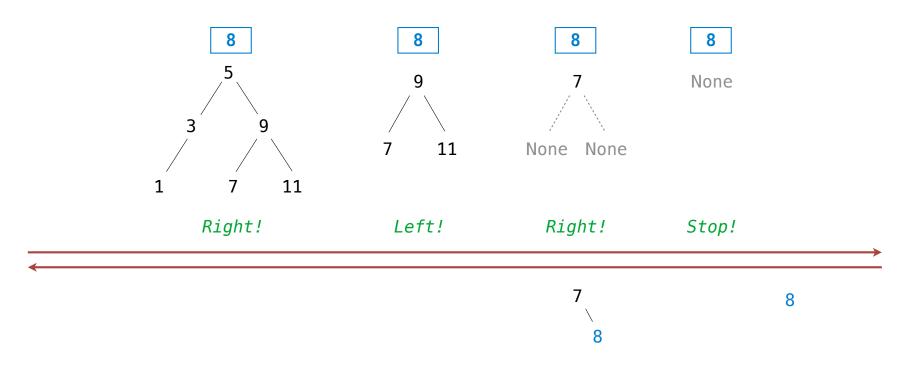


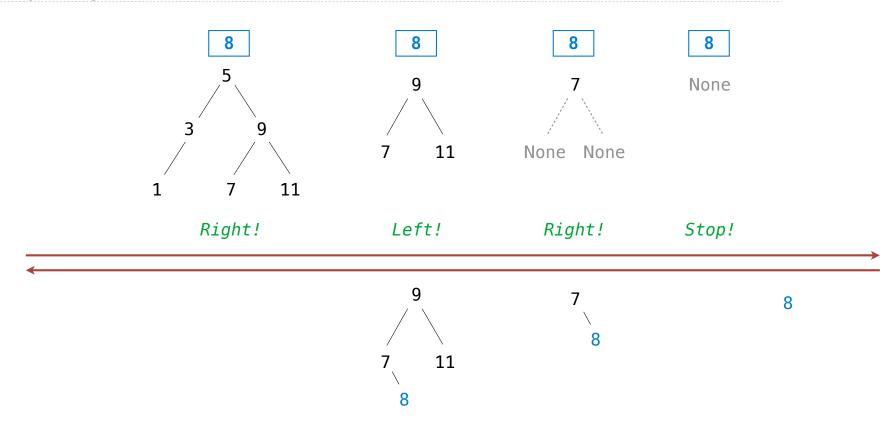


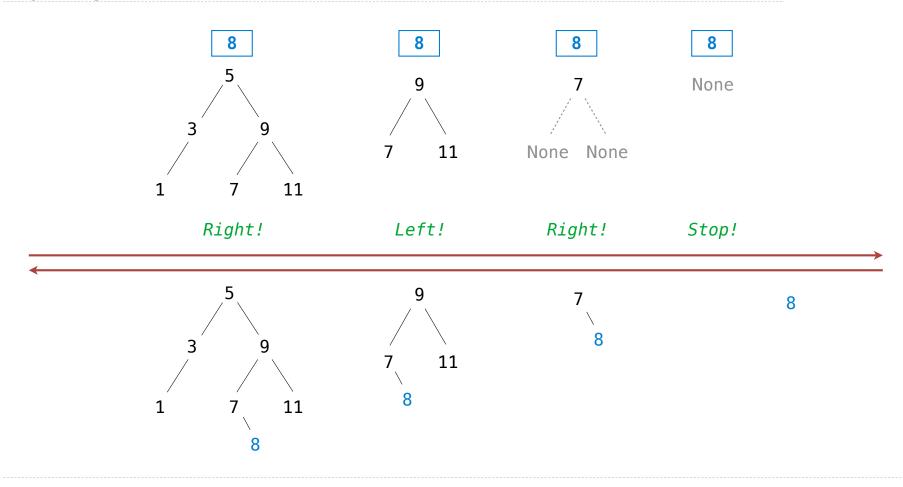


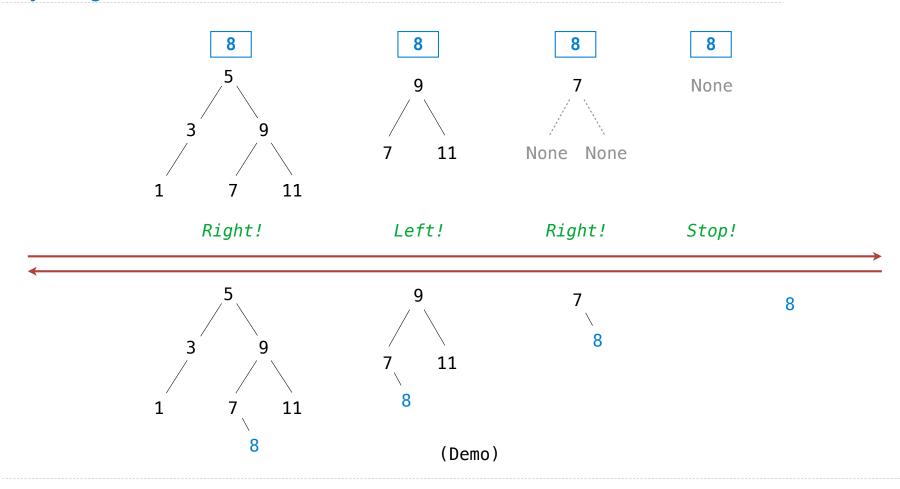


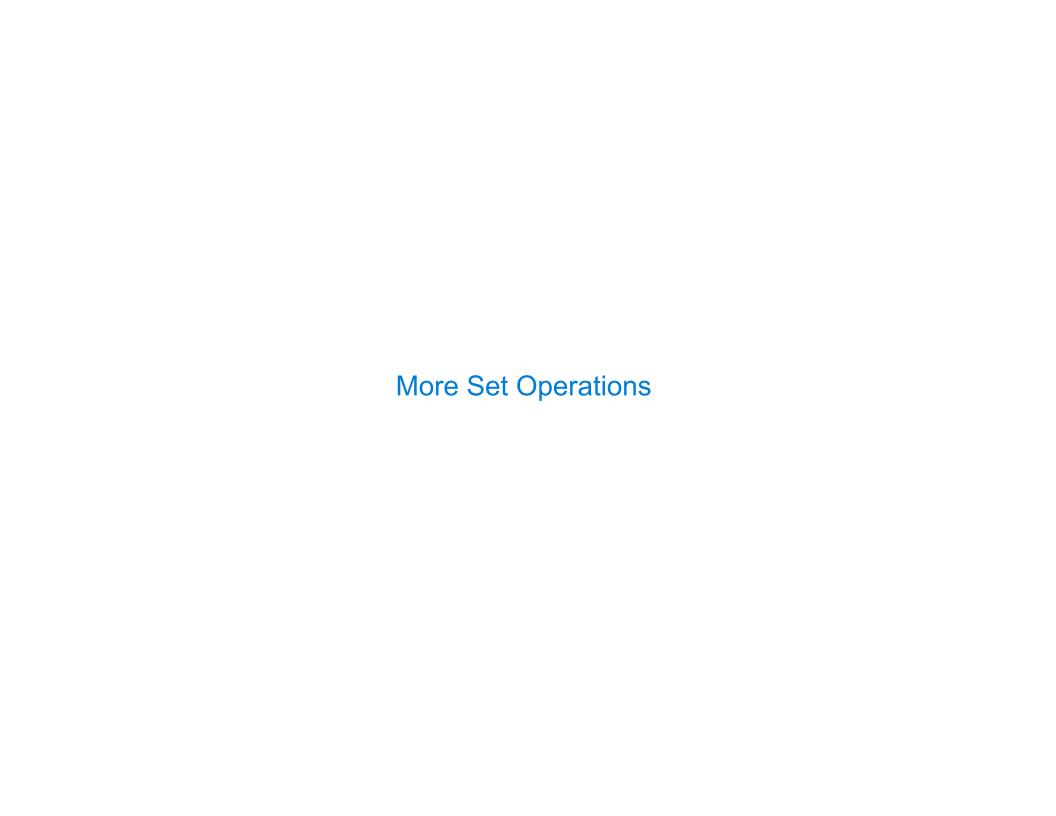
8











Sets as ordered sequences:

Sets as ordered sequences:

Adjoining an element to a set

Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets

Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets

Sets as binary trees:

Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets

Sets as binary trees:

Intersection of two sets

Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets

Sets as binary trees:

- Intersection of two sets
- Union of two sets

Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets

Sets as binary trees:

- Intersection of two sets
- Union of two sets
- Balancing a tree

Sets as ordered sequences:

- Adjoining an element to a set
- Union of two sets

Sets as binary trees:

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That's all on homework 7!