## 61A Lecture 8

## Wednesday, September 18

## Announcements

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-Homework 3 is due in two weeks: Tuesday 10/1 @ 11:59pm
= It contains lots of recursion problems, for practice!
- Optional Hog strategy contest ends Thursday 10/3 @ 11:59pm


## Hog Contest Rules

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## Fall 2011 Winners

## Keegan Mann,

Yan Duan \& Ziming Li, Brian Prike \& Zhenghao Qian, Parker Schuh \& Robert Chatham

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Chenyang Yuan, Joseph Hui

Fall 2013 Winners

YOUR NAME COULD BE HERE... FOREVER!

## Order of Recursive Calls

## The Cascade Function

(Demo)

## The Cascade Function

|  |  |
| :---: | :---: |
| 1 | def cascade $(n):$ |
| 2 | if $n<10:$ |
| 3 | $\operatorname{print}(n)$ |
| 4 | else: |
| 5 | $\operatorname{print}(n)$ |
| 6 | $\operatorname{cascade}(n / / 10)$ |
| $\Rightarrow 7$ | $\operatorname{print}(n)$ |
| $\Rightarrow 8$ |  |
| 9 | cascade $(123)$ |

## (Demo)



## The Cascade Function

(Demo)

cascade
1

| $\begin{array}{c}\text { Return } \\ \text { value }\end{array}$ | None |
| :--- | :--- |

## The Cascade Function

## Program output:

123
12
1
12

```
def cascade(n):
    if n < 10:
        print(n)
    else:
        print(n)
        cascade(n//10)
        print(n)
    g cascade(123)
```

(Demo)

```
```

Global frame }\longrightarrow\mathrm{ func cascade(n)

```
```

Global frame }\longrightarrow\mathrm{ func cascade(n)
cascade
cascade
cascade
cascade
n 123
n 123
cascade
cascade
n 12
n 12

- Each cascade frame is from a
- Each cascade frame is from a
different call to cascade.

```
```

    different call to cascade.
    ```
```

Return None
value None
cascade

| n | 1 |
| ---: | :--- |
| $\begin{aligned} \text { Return } \\ \text { value }\end{aligned}$ | None |

## The Cascade Function

## Program output:

123
12
1
12

```
def cascade(n):
    if n < 10:
        print(n)
    else:
        print(n)
        cascade(n//10)
        print(n)
    cascade(123)
```

(Demo)

```
```

Global frame }\longrightarrow\mathrm{ func cascade(n)

```
```

Global frame }\longrightarrow\mathrm{ func cascade(n)
cascade
cascade
cascade
cascade
n 123
n 123
cascade
cascade
n 12
n 12
Return None
Return None
-Until the Return value appears,
-Until the Return value appears,
that call has not completed.

```
```

    that call has not completed.
    ```
```

- Each cascade frame is from a different call to cascade.
- Until the Return value appears, that call has not completed.

```
cascade
    n 1
Return None
```


## The Cascade Function

```
def cascade(n):
    if n < 10:
        print(n)
    else:
        print(n)
        cascade(n//10)
        print(n)
    cascade(123)
```


## Program output:

123
12
1
12
(Demo)

```
Global frame }\longrightarrow\mathrm{ func cascade(n)
    cascade
cascade
            n 123
cascade
    n 12
Return None
    value None
cascade
```

- Each cascade frame is from a different call to cascade.
- Until the Return value appears, that call has not completed.
- Any statement can appear before or after the recursive call.
n 1

| $\begin{array}{c}\text { Return } \\ \text { value }\end{array}$ | None |
| :---: | :---: |

## The Cascade Function

| 1 | def cascade $(n):$ |
| :---: | :---: |
| 2 | if $n<10:$ |
| 3 | $\operatorname{print}(n)$ |
| 4 | else: |
| 5 | $\operatorname{print}(n)$ |
| 6 | $\operatorname{cascade}(n / / 10)$ |
| $\Rightarrow 7$ | $\operatorname{print}(n)$ |
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## Program output:

123
12
1
12
(Demo)

```
Global frame }\longrightarrow\mathrm{ func cascade(n)
    cascade
cascade
            n 123
cascade
    n 12
Return None
    value None
cascade
```

- Each cascade frame is from a different call to cascade.
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| n | 1 |
| :--- | :--- |

| $\begin{array}{c}\text { Return } \\ \text { value }\end{array}$ | None |
| :--- | :--- |

## The Cascade Function

```
def cascade(n):
    if n < 10:
        print(n)
    else:
        print(n)
        cascade(n//10)
        print(n)
    cascade(123)
```


## Program output:

123
12
1
12
(Demo)

```
Global frame 
    cascade
cascade
            n 123
cascade
    n 12
Return None
    value
cascade
```

- Each cascade frame is from a different call to cascade.
- Until the Return value appears, that call has not completed.
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| $n$ | 1 |
| :--- | :--- |

| Return |
| ---: | ---: |
| value | None


| $\begin{array}{r}\text { Return } \\ \text { value }\end{array}$ | None |
| ---: | :--- |

## The Cascade Function

```
1 def cascade(n):
    if n < 10:
        print(n)
    else:
        print(n)
        cascade(n//10)
        print(n)
    cascade(123)
```


## Program output:


(Demo)


## cascade

n 12
Return
value None value
cascade

| $n$ | 1 |
| ---: | :--- |
| $\begin{array}{r}\text { Return } \\ \text { value }\end{array}$ | None |

- Each cascade frame is from a different call to cascade.
- Until the Return value appears, that call has not completed.
- Any statement can appear before or after the recursive call.


## The Cascade Function



Two Definitions of Cascade

Two Definitions of Cascade
(Demo)

```
def cascade(n):
    if n < 10:
        print(n)
    else:
        print(n)
        cascade(n//10)
        print(n)
```


## Two Definitions of Cascade

```
        (Demo)
    def cascade(n):
        print(n)
        if n >= 10:
        cascade(n//10)
        print(n)
```

```
```

def cascade(n):

```
```

def cascade(n):

```
```

def cascade(n):
if n < 10:
if n < 10:
if n < 10:
print(n)
print(n)
print(n)
else:
else:
else:
print(n)
print(n)
print(n)
cascade(n//10)
cascade(n//10)
cascade(n//10)
print(n)

```
```

        print(n)
    ```
```

        print(n)
    ```
```


## Two Definitions of Cascade

(Demo)

```
def cascade(n):
    if n < 10:
        print(n)
    else:
        print(n)
        cascade(n//10)
        print(n)
```

def cascade(n):
print(n)
if n >= 10:
cascade(n//10)
print(n)

- If two implementations are equally clear, then shorter is usually better.
- In this case, the longer implementation is more clear (at least to me).


## Two Definitions of Cascade

(Demo)
def cascade(n):
print(n)
if n >= 10:
cascade(n//10)
print(n)

```
def cascade(n):
```

def cascade(n):
if n < 10:
if n < 10:
print(n)
print(n)
else:
else:
print(n)
print(n)
cascade(n//10)
cascade(n//10)
print(n)

```
        print(n)
```

- If two implementations are equally clear, then shorter is usually better.
- In this case, the longer implementation is more clear (at least to me).
- When learning to write recursive functions, put the base cases first.


## Two Definitions of Cascade

(Demo)

```
def cascade(n):
    if n < 10:
        print(n)
    else:
        print(n)
        cascade(n//10)
        print(n)
```

def cascade(n):
print(n)
if $\mathrm{n}>=10$ :
cascade(n//10)
print(n)

- If two implementations are equally clear, then shorter is usually better.
- In this case, the longer implementation is more clear (at least to me).
- When learning to write recursive functions, put the base cases first.
- Both are recursive functions, even though only the first has typical structure.

Tree Recursion

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Tree-shaped processes arise whenever executing the body of a recursive function makes more than one call to that function.

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$$
\text { n: } 1,2,3,4,5,6,7,8,9,
$$



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Tree-shaped processes arise whenever executing the body of a recursive function makes more than one call to that function.

n: 1, 2, 3, 4, 5, 6, 7, 8, 9,<br>fib(n): 0, 1, 1, 2, 3, 5, 8, 13, 21,



## Tree Recursion

Tree-shaped processes arise whenever executing the body of a recursive function makes more than one call to that function.

$$
\begin{aligned}
& n: 1,2,3,4,5,6,7,8,9, 35 \\
& \text { fib(n): } 0,1,1,2,3,5,8,13,21,
\end{aligned}
$$



## Tree Recursion

Tree-shaped processes arise whenever executing the body of a recursive function makes more than one call to that function.

$$
\begin{array}{rrrr}
n: 1,2,3,4,5,6,7,8,9, & \ldots, & 35 \\
\text { fib(n): } 0,1,1,2,3,5,8,13,21, & \ldots, 702,887
\end{array}
$$



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$\mathrm{n}: 1,2,3,4,5,6,7,8,9, \quad 35$
fib(n): 0, 1, 1, 2, 3, 5, 8, 13, 21,..
def fib(n):
def fib(n):
if n == 1:
if n == 1:
return 0
return 0


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$\mathrm{n}: 1,2,3,4,5,6,7,8,9, \quad 35$
fib(n): 0, 1, 1, 2, 3, 5, 8, 13, 21,..
def fib(n):
def fib(n):
if n == 1:
if n == 1:
return 0
return 0
elif n == 2:
elif n == 2:


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Tree-shaped processes arise whenever executing the body of a recursive function makes more than one call to that function.

```
        n: 1, 2, 3, 4, 5, 6, 7, 8, 9, ... , 35
        fib(n): 0, 1, 1, 2, 3, 5, 8, 13, 21, ..., 5,702,887
def fib(n):
    if n == 1:
        return 0
    elif n == 2:
        return 1
```



## Tree Recursion

Tree-shaped processes arise whenever executing the body of a recursive function makes more than one call to that function.

```
            n: 1, 2, 3, 4, 5, 6, 7, 8, 9, !., 35
fib(n): 0, 1, 1, 2, 3, 5, 8, 13, 21, ..., 5,702,887
```

```
def fib(n):
```

def fib(n):
if n == 1:
if n == 1:
return 0
return 0
elif n == 2:
elif n == 2:
return 1
return 1
else:

```
    else:
```



## Tree Recursion

Tree-shaped processes arise whenever executing the body of a recursive function makes more than one call to that function.

```
            n: 1, 2, 3, 4, 5, 6, 7, 8, 9, #., 35
fib(n): 0, 1, 1, 2, 3, 5, 8, 13, 21, ..., 5,702,887
```

```
def fib(n):
```

def fib(n):
if n == 1:
if n == 1:
return 0
return 0
elif n == 2:
elif n == 2:
return 1
return 1
else:
else:
return fib(n-2) + fib(n-1)

```
        return fib(n-2) + fib(n-1)
```



## A Tree-Recursive Process

The computational process of fib evolves into a tree structure

## A Tree-Recursive Process

The computational process of fib evolves into a tree structure
fib(6)

## A Tree-Recursive Process

The computational process of fib evolves into a tree structure


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## Repetition in Tree-Recursive Computation

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This process is highly repetitive; fib is called on the same argument multiple times.

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We can speed up this computation dramatically in a few weeks by remembering results.

## Example: Counting Partitions

## Counting Partitions

The number of partitions of a positive integer $\mathbf{n}$, using parts up to size $m$, is the number of ways in which $n$ can be expressed as the sum of positive integer parts up to $m$ in increasing order.

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partition(6, 4)

## Counting Partitions

The number of partitions of a positive integer $\mathbf{n}$, using parts up to size $m$, is the number of ways in which $n$ can be expressed as the sum of positive integer parts up to $m$ in increasing order.
partition(6, 4)

$$
\begin{aligned}
& 2+4=6 \\
& 1+1+4=6 \\
& 3+3=6 \\
& 1+2+3=6 \\
& 1+1+1+3=6 \\
& 2+2+2=6 \\
& 1+1+2+2=6 \\
& 1+1+1+1+2=6 \\
& 1+1+1+1+1+1=6
\end{aligned}
$$

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partition(6, 4)

$$
\begin{aligned}
& 2+4=6 \\
& 1+1+4=6 \\
& 3+3=6 \\
& 1+2+3=6 \\
& 1+1+1+3=6 \\
& 2+2+2=6 \\
& 1+1+2+2=6 \\
& 1+1+1+1+2=6 \\
& 1+1+1+1+1+1=6
\end{aligned}
$$

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$$
\begin{aligned}
& 2+4=6 \\
& 1+1+4=6 \\
& 3+3=6 \\
& 1+2+3=6 \\
& 1+1+1+3=6 \\
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& 1+1+1+1+2=6 \\
& 1+1+1+1+1+1=6
\end{aligned}
$$



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$$
\begin{aligned}
& 2+4=6 \\
& 1+1+4=6 \\
& 3+3=6 \\
& 1+2+3=6 \\
& 1+1+1+3=6 \\
& 2+2+2=6 \\
& 1+1+2+2=6 \\
& 1+1+1+1+2=6 \\
& 1+1+1+1+1+1=6
\end{aligned}
$$



## Counting Partitions

The number of partitions of a positive integer $\mathbf{n}$, using parts up to size $m$, is the number of ways in which $n$ can be expressed as the sum of positive integer parts up to $m$ in increasing order.
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- Recursive decomposition: finding simpler instances of the problem.



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partition(6, 4)

- Recursive decomposition: finding simpler instances of the problem.

Explore two possibilities:


## Counting Partitions

The number of partitions of a positive integer $\mathbf{n}$, using parts up to size $m$, is the number of ways in which $n$ can be expressed as the sum of positive integer parts up to m in increasing order.
partition(6, 4)

- Recursive decomposition: finding simpler instances of the problem.

Explore two possibilities:

- Use at least one 4



## Counting Partitions

The number of partitions of a positive integer $\mathbf{n}$, using parts up to size $m$, is the number of ways in which $n$ can be expressed as the sum of positive integer parts up to m in increasing order.
partition(6, 4)

- Recursive decomposition: finding simpler instances of the problem.
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- Solve two simpler problems:
- partition(2, 4)



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partition(6, 4)

- Recursive decomposition: finding simpler instances of the problem.

Explore two possibilities:

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- Don't use any 4
- Solve two simpler problems:
- partition(2, 4)
- partition(6, 3)



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```

    simpler instances of the problem.
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- partition(2, 4)
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```
def count_partitions(n, m):
```

else:
with_m $=$ count_partitions( $n-m, m)$
without_m $=$ count_partitions( $n, m-1$ )

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else:
with_m $=$ count_partitions ( $n-m, m)$
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return with_m + without_m


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```



```
    -partition(6, 3) without_m = count_partitions(n, m-1)
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    exploring different choices.
```


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```
-Recursive decomposition: finding
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Explore two possibilities:
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-Don't use any 4
-Solve two simpler problems:
    partition(2,4) =-= =-=-=-=-=- else:
    manm, with_m = count_partitions(n-m, m)
```



```
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    return with_m + without_m
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```


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```
-Recursive decomposition: finding
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```

```
def count_partitions(n, m):
```

def count_partitions(n, m):
if n == 0:
if n == 0:
Explore two possibilities:
-Use at least one 4
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- Solve two simpler problems:

```

```

    m, m=-=-=-=> with_m = count_partitions(n-m, m)
    ```

```

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    return with_m + without_m
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    ```

\section*{Counting Partitions}

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```

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Explore two possibilities:

```
```

def count_partitions(n, m):

```
def count_partitions(n, m):
    if n == 0:
    if n == 0:
        return 1
        return 1
    -Use at least one 4
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-Solve two simpler problems:
    partition(2,4) =-=-=-=-=-=-=-- else:
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```



```
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- Solve two simpler problems:
    partition(2,4)-==-=-=-=-=-=--- else:
    m
```



```
    -Tree recursion often involves
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```

```
def count_partitions(n, m):
```

def count_partitions(n, m):
if n == 0:
if n == 0:
return 1
return 1
elif n < 0:
elif n < 0:
Use at least one 4

- Don't use any 4
- Solve two simpler problems:

```


```

-Tree recursion often involves return with_m + without_m exploring different choices.

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    elif n < 0:
        return 0
        return 0
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```


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    partition(2, 4) =- =- =- =- = = = =- =- - - else:
    manm, with_m = count_partitions(n-m, m)
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Winning Hog

\section*{How to Win at Hog}

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What is the chance that I'll score at least \(\mathbf{k}\) points rolling \(\mathbf{n}\) six-sided dice?

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Number of possible rolls

The number of possible rolls is pow(6, n).

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What is the chance that I'll score at least k points rolling n six-sided dice?
\(\frac{\text { Number of ways to score at least } \mathbf{k}}{\text { Number of possible rolls }}\)

The number of possible rolls is pow(6, n).
The number of ways to score at least \(\mathbf{k}\) in \(\mathbf{n}\) rolls can be computed using tree recursion!

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What is the chance that I'll score at least k points rolling n six-sided dice?

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The number of possible rolls is pow(6, n).
The number of ways to score at least \(\mathbf{k}\) in \(\mathbf{n}\) rolls can be computed using tree recursion!
Sum over each possible dice outcome d that does not pig out:
the number of ways to score at least \(\mathbf{k}\) - \(\mathbf{d}\) points using \(\mathbf{n}\) - \(\mathbf{1}\) rolls.

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What is the chance that I'll score at least k points rolling n six-sided dice?

> Number of ways to score at least k

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The number of ways to score at least \(\mathbf{k}\) in \(\mathbf{n}\) rolls can be computed using tree recursion!
Sum over each possible dice outcome d that does not pig out:
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Base case: The number of ways to score at least 0 is pow(5, n).

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What is the chance that I'll score at least k points rolling n six-sided dice?
```

    Number of ways to score at least k
    ```

Number of possible rolls

The number of possible rolls is pow(6, n).
The number of ways to score at least \(\mathbf{k}\) in \(\mathbf{n}\) rolls can be computed using tree recursion!
Sum over each possible dice outcome d that does not pig out:
the number of ways to score at least \(\mathbf{k}\) - \(\mathbf{d}\) points using \(\mathbf{n} \mathbf{- 1}\) rolls.
Base case: The number of ways to score at least 0 is pow(5, n).
Base case: The number of ways to score positive points in 0 rolls is 0 .```

