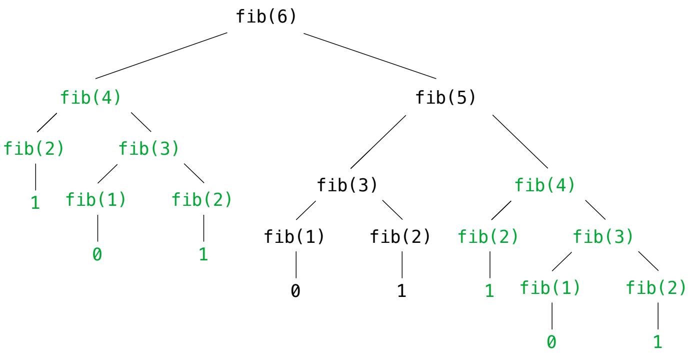
Last day of Midterm 2 Material

61A Lecture 23

Friday, October 19

Trees with Internal Node Values

Trees can have values at their roots as well as their leaves.



Trees with Internal Node Values (Entries)

Trees need not only have values at their leaves.

```
class Tree(object):
    def __init__(self, entry, left=None, right=None):
        self.entry = entry
        self.left = left
        self.right = right

def fib_tree(n):
    if n == 1:
        return Tree(0)
    if n == 2:
        return Tree(1)
    left = fib_tree(n-2)
    right = fib_tree(n-1)
    return Tree(left.entry + right.entry, left, right)
```

Valid if left and right are each either None or a Tree instance

A valid tree cannot be a subtree of itself (no cycles!)

Demo

The Consumption of Time

Implementations of the same functional abstraction can require different amounts of time to compute their result.

`def count_factors(n):` (Demo)

Time (remainders)

```
factors = 0
for k in range(1, n+1):
    if n % k == 0:
        factors += 1
return factors
```

n

```
sqrt_n = sqrt(n)
k, factors = 1, 0
while k < sqrt_n:
    if n % k == 0:
        factors += 2
    k += 1
if k * k == n:
    factors += 1
return factors
```

$\lfloor\sqrt{n}\rfloor$

The Consumption of Space

Which environment frames do we need to keep during evaluation?

Each step of evaluation has a set of **active** environments.

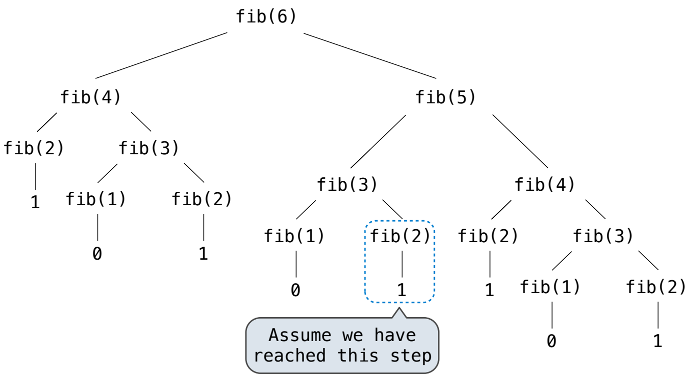
Values and frames in active environments consume memory.

Memory used for other values and frames can be reclaimed.

Active environments:

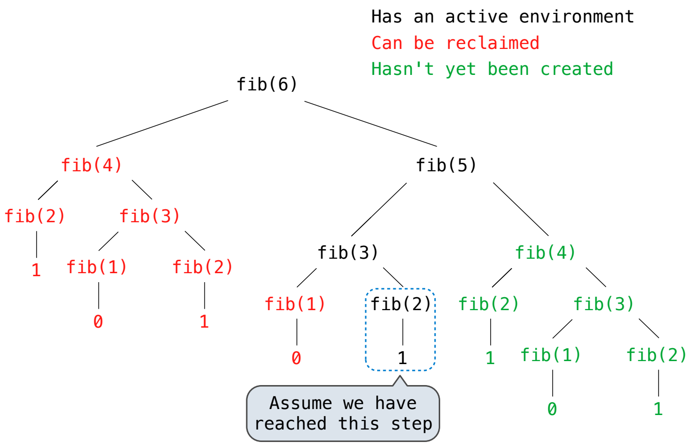
- Environments for any statements currently being executed
- Parent environments of functions named in active environments

Fibonacci Memory Consumption



Assume we have reached this step

Fibonacci Memory Consumption



Has an active environment

Can be reclaimed

Hasn't yet been created

Assume we have reached this step

Order of Growth

A method for bounding the resources used by a function as the "size" of a problem increases

n: size of the problem

R(n): Measurement of some resource used (time or space)

$$R(n) = \Theta(f(n))$$

means that there are positive constants k_1 and k_2 such that

$$k_1 \cdot f(n) \leq R(n) \leq k_2 \cdot f(n)$$

for sufficiently large values of ***n***.

Iteration vs Memoized Tree Recursion

Iterative and memoized implementations are not the same.

```
def fib_iter(n):
    prev, curr = 1, 0
    for _ in range(n-1):
         prev, curr = curr, prev + curr
    return curr
```

Time	Space
$\Theta(n)$	$\Theta(1)$

```
@memo
def fib(n):
    if n == 1:
        return 0
    if n == 2:
        return 1
    return fib(n-2) + fib(n-1)
```

Time	Space
$\Theta(n)$	$\Theta(n)$

The Consumption of Time

Implementations of the same functional abstraction can require different amounts of time.

```
def count_factors(n):


    factors = 0
    for k in range(1, n+1):
        if n % k == 0:
            factors += 1
    return factors
```

Time	Space
$\Theta(n)$	$\Theta(1)$

```
    sqrt_n = sqrt(n)
    k, factors = 1, 0
    while k < sqrt_n:
        if n % k == 0:
            factors += 2
        k += 1
    if k * k == n:
        factors += 1
    return factors
```

Time	Space
$\Theta(\sqrt{n})$	$\Theta(1)$

Exponentiation

Goal: one more multiplication lets us double the problem size.

```
def exp(b, n):
    if n == 0:
        return 1
    return b * exp(b, n-1)
```

$$b^n = \begin{cases} 1 & \text{if } n = 0 \\ b \cdot b^{n-1} & \text{otherwise} \end{cases}$$

```
def square(x):
    return x*x

def fast_exp(b, n):
    if n == 0:
        return 1
    if n % 2 == 0:
        return square(fast_exp(b, n//2))
    else:
        return b * fast_exp(b, n-1)
```

$$b^n = \begin{cases} 1 & \text{if } n = 0 \\ (b^{\frac{1}{2}n})^2 & \text{if } n \text{ is even} \\ b \cdot b^{n-1} & \text{if } n \text{ is odd} \end{cases}$$

Exponentiation

Goal: one more multiplication lets us double the problem size.

```
def exp(b, n):
    if n == 0:
        return 1
    return b * exp(b, n-1)
```

Time	Space
$\Theta(n)$	$\Theta(n)$

```
def square(x):
    return x*x

def fast_exp(b, n):
    if n == 0:
        return 1
    if n % 2 == 0:
        return square(fast_exp(b, n//2))
    else:
        return b * fast_exp(b, n-1)
```

Time	Space
$\Theta(\log n)$	$\Theta(\log n)$

Comparing orders of growth (n is the problem size)

$\Theta(b^n)$ Exponential growth! Recursive fib takes

$\Theta(\phi^n)$ steps, where $\phi = \frac{1+\sqrt{5}}{2} \approx 1.61828$

$\Theta(n^6)$ ⇢

Incrementing the problem scales R(n) by a factor.

$\Theta(n^2)$ Quadratic growth. E.g., operations on all pairs.

Incrementing n increases R(n) by the problem size n.

$\Theta(n)$ Linear growth. Resources scale with the problem.

$\Theta(\sqrt{n})$ ⇢

$\Theta(\log n)$ Logarithmic growth. These processes scale well.

Doubling the problem only increments R(n).

$\Theta(1)$ Constant. The problem size doesn't matter.