## 1 Readings

Benenti, Casati, and Strini:
Classical circuits and computation Ch.1.2, 2.6
Quantum Gates Ch. 3.2-3.4
Universality Ch. 3.5-3.6

## 2 Unitary Operators

A postulate of quantum physics is that quantum evolution is unitary. That is, if we have some arbitrary quantum system $U$ that takes as input a state $|\phi\rangle$ and outputs a different state $U|\phi\rangle$, then we can describe $U$ as a unitary linear transformation, defined as follows.
If $U$ is any linear transformation, the adjoint of $U$, denoted $U^{\dagger}$, is defined by $(U \vec{v}, \vec{w})=\left(\vec{v}, U^{\dagger} \vec{w}\right)$. In a basis, $U^{\dagger}$ is the conjugate transpose of $U$; for example, for an operator on $\mathscr{C}^{2}$,

$$
U=\left(\begin{array}{ll}
a & b \\
c & d
\end{array}\right) \Rightarrow U^{\dagger}=\left(\begin{array}{cc}
\frac{a}{a} & \bar{c} \\
\bar{b} & \frac{1}{d}
\end{array}\right) .
$$

We say that $U$ is unitary if $U^{\dagger}=U^{-1}$. For example, rotations and reflections are unitary. Also, the composition of two unitary transformations is also unitary (Proof: $U, V$ unitary, then $(U V)^{\dagger}=V^{\dagger} U^{\dagger}=V^{-1} U^{-1}=$ $\left.(U V)^{-1}\right)$.

Some properies of a unitary transformation $U$ :

- The rows of $U$ form an orthonormal basis.
- The colums of $U$ form an orthonormal basis.
- $U$ preserves inner products, i.e. $(\vec{v}, \vec{w})=(U \vec{v}, U \vec{w})$. Indeed, $(U \vec{v}, U \vec{w})=(U|v\rangle)^{\dagger} U|w\rangle=\langle v| U^{\dagger} U|w\rangle=$ $\langle v \mid w\rangle$. Therefore, $U$ preserves norms and angles (up to sign).
- The eigenvalues of $U$ are all of the form $e^{i \theta}$ (since $U$ is length-preserving, i.e., $(\vec{v}, \vec{v})=(U \vec{v}, U \vec{v})$ ).
- $U$ can be diagonalized into the form

$$
\left(\begin{array}{cccc}
e^{i \theta_{1}} & 0 & \cdots & 0 \\
0 & \ddots & \ddots & 0 \\
\vdots & \ddots & \ddots & \vdots \\
0 & \cdots & 0 & e^{\theta_{d}}
\end{array}\right)
$$

## 3 Schrödinger's Equation

Schrödinger's equation is the equation of motion which describes the evolution in time of the quantum state.

$$
i \hbar \frac{d|\psi(t)\rangle}{d t}=H|\psi\rangle .
$$

Here $\hbar$ is a constant (called Planck's constant - we'll usually assume $\hbar=1$ ), and $H$ is a linear Hamiltonian which is Hermitian, $H^{\dagger}=H$. Equivalently, $H$ has an orthonormal set of eigenvectors $\left|\psi_{i}\right\rangle$, all with real eigenvalues $\lambda_{i}: H\left|\phi_{i}\right\rangle=\lambda_{i}\left|\phi_{i}\right\rangle$.
For those of you who are familiar with Schrödinger's equation, the unitarity restriction on quantum gates is simply the time-discrete version of the restriction that the Hamiltonian is Hermitian. This is a particular instance of the general relation between a unitary operator $U$ and a Hermitian operator A

$$
U=e^{i} A
$$

which follows directly from $U U^{\dagger}=1, A^{\dagger}=A$, hence $U^{\dagger}=\exp \left(-i A^{\dagger}\right)=\exp (-i A)$.
We will now prove explicitly that if the system satisfies Schrödinger's equation, then its evolution in discrete time is described by a unitary operator and determine this operator in terms of the eigenvalues of $H$. (We will assume that $H$ is time independent.)
Write $|\psi(t)\rangle$ in the basis of eigenvectors of $H$ :

$$
\begin{gathered}
|\psi(t)\rangle=\sum_{j} a_{i}(t)\left|\phi_{j}\right\rangle \\
\Downarrow \\
i \hbar \frac{d \Sigma a_{j}\left|\phi_{j}\right\rangle}{d t}=H \Sigma a_{j}\left|\phi_{j}\right\rangle=\Sigma a_{j} \lambda_{j}\left|\phi_{j}\right\rangle \\
\Downarrow \\
i \hbar \frac{d a_{j}}{d t}=\lambda_{j} a_{j} \\
\Downarrow \\
a_{j}(t)=e^{-\frac{i}{\hbar} \lambda_{j} t} a_{j}(0) \\
\Downarrow \\
|\psi(t)\rangle=e^{-\frac{i}{\hbar} \lambda_{j} t} a_{j}(0)\left|\phi_{j}\right\rangle
\end{gathered}
$$

We get that the change after a discrete time difference is unitary:

$$
|\psi(t)\rangle=\left(\begin{array}{cccc}
e^{-\frac{i}{\hbar} \lambda_{1} t} & & & 0 \\
& \cdot & & \\
& & \cdot & \\
0 & & e^{-\frac{i}{\hbar} \lambda_{d} t}
\end{array}\right)\left(\begin{array}{c}
a_{0} \\
\cdot \\
\cdot \\
a_{d}
\end{array}\right)=U(t)|\psi(0)\rangle
$$

In this basis, $U(t)$ is diagonal.

## 4 Quantum Gates

We already had some simple examples of unitary transforms, or "quantum gates". Here are most of the common ones you will encounter.

### 4.1 One-qubit gates:

- Hadamard Gate.

$$
\begin{aligned}
& H=\frac{1}{\sqrt{2}}\left(\begin{array}{cc}
1 & 1 \\
1 & -1
\end{array}\right) \\
& H|0\rangle=\frac{1}{\sqrt{2}}(|0\rangle+|1\rangle)=|+\rangle \\
& H|1\rangle=\frac{1}{\sqrt{2}}(|0\rangle-|1\rangle)=|-\rangle
\end{aligned}
$$

The Hadamard Gate is one of the most important gates. Note that $H^{\dagger}=H-$ since $H$ is real and symmetric - and $H^{2}=I$.
In the complex plane $H$ can be visualized as a reflection around $\pi / 8$, or a rotation around $\pi / 4$ followed by a reflection.
On the Bloch sphere $H$ can also be visualized in several ways. One is a rotation of $\pi / 2$ about the $y$-axis, followed by reflection in the $x-y$ plane (see Nielsen and Chuang, p. ). Another is a rotation of $\pi$ about the axis $(1 / \sqrt{2}, 0,1 / \sqrt{2})$ (Benenti, p. 111).
Note the action of $H$ on larger number of qubits:

$$
\begin{array}{rc}
H \otimes H|00\rangle \equiv H^{\otimes 2}|00\rangle & =\frac{|00\rangle+|01\rangle+|10\rangle+|11\rangle}{\sqrt{2^{2}}} \\
H^{\otimes n}\left|00 \ldots . .0_{n}\right\rangle & =\frac{1}{\sqrt{2^{n}}} \sum_{x=0}^{2^{n}-1}|x\rangle
\end{array}
$$

Thus $H^{\otimes n}$ produces an equal superposition of all computational basis states.

- Rotation Gate. This rotates in the complex plane by $\theta$.

$$
R=\left(\begin{array}{cc}
\cos \theta & -\sin \theta \\
\sin \theta & \cos \theta
\end{array}\right)
$$

- NOT Gate, also known as bit flip gate, or $X$ (Pauli $X$ ). This flips a bit from 0 to 1 and vice versa.

$$
N O T=\left(\begin{array}{ll}
0 & 1 \\
1 & 0
\end{array}\right)
$$

- Phase Flip, also known as $Z$ (Pauli $Z$ ).

$$
Z=\left(\begin{array}{cc}
1 & 0 \\
0 & -1
\end{array}\right)
$$

The phase flip is a NOT gate acting in the $|+\rangle=\frac{1}{\sqrt{2}}(|0\rangle+|1\rangle),|-\rangle=\frac{1}{\sqrt{2}}(|0\rangle-|1\rangle)$ basis. Indeed, $Z|+\rangle=|-\rangle$ and $Z|-\rangle=|+\rangle$.

- General Phase Gate, $R_{z}(\boldsymbol{\delta})$.

$$
R_{z}(\boldsymbol{\delta})=\left(\begin{array}{cc}
1 & 0 \\
0 & e^{i \delta}
\end{array}\right)
$$

Clearly $Z=R_{z}(\pi)$. There are several other special phase gates that are commonly used: $S=R_{z}(\pi / 2)$, $T=_{z}(\pi / 4)$. The latter is sometimes referred to as the $\pi / 8$ gate.

$$
S=\left(\begin{array}{cc}
1 & 0 \\
0 & i
\end{array}\right) \quad T \equiv \pi / 8=\left(\begin{array}{cc}
0 & 1 \\
1 & e^{i \pi / 4}
\end{array}\right)=e^{i \pi / 8}\left(\begin{array}{cc}
e^{-i \pi / 8} & 0 \\
0 & e^{i \pi / 8}
\end{array}\right)
$$

- Phaseflips and bitflips are related by conjugation

Conjugation of $X$ by $H$ means premultiplying $X$ by $H^{-1}$ and postmultiplying it by $H$. But $H=H^{-1}$.
Claim: $H X H=Z$. See Figure 1 .
We can prove this by multiplying out the matrices, or by making use of the decomposition of $H$ into an $X$ and a $Z$ gate:

$$
H=\frac{1}{\sqrt{2}}\left[\begin{array}{rr}
1 & 1 \\
1 & -1
\end{array}\right]=\frac{1}{\sqrt{2}}\left[\left[\begin{array}{ll}
0 & 1 \\
1 & 0
\end{array}\right]+\left[\begin{array}{rr}
1 & 0 \\
0 & -1
\end{array}\right]\right]=\frac{X+Z}{\sqrt{2}}
$$

Then $\quad\left(\frac{x+z}{\sqrt{2}}\right) X\left(\frac{x+z}{\sqrt{2}}\right)=$

$$
\begin{aligned}
& {\left[\frac{X+Z}{\sqrt{2}}\right]\left[\frac{X^{2}+X Z}{\sqrt{2}}\right]=} \\
& {\left[\frac{X+Z}{\sqrt{2}}\right]\left[\frac{I+X Z}{\sqrt{2}}\right]=} \\
& \frac{X I+X X Z+Z I+Z X Z}{2}= \\
& \frac{X+Z+Z+-X}{2}= \\
& \frac{2 Z}{2}=Z
\end{aligned}
$$

Conversely, $H Z H=X$ (Figure 2). Prove this for yourself.

- Any unitary operation on a single qubit can be constructed with various combinations of gates:
$H, R_{z}(\delta)$, e.g.,

$$
R_{z}(\pi / 2+\phi) H R_{z}(\theta) H|0\rangle=e^{i \theta / 2}\left(\cos \theta / 2|0\rangle+e^{i \phi} \sin \theta / 2|1\rangle\right)
$$

$H, X, T=R_{z}(\pi / 4)$
$X, Y, Z$ (Euler rotations)

### 4.2 Two-qubit gates:

- Any one-qubit gate can be tensored with itself or another gate to make a two-qubit gate, as done above for $H \otimes H$. Such tensor products of one-qubit gates have no ability to generate entanglement and are referred to as 'local' gates.
- Controlled Not (CNOT).

$$
\mathrm{CNOT}=\left(\begin{array}{llll}
1 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 0 & 1 \\
0 & 0 & 1 & 0
\end{array}\right)
$$

The first bit of a CNOT gate is the "control bit;" the second is the "target bit." The control bit never changes, while the target bit flips if and only if the control bit is 1 .
The CNOT gate is usually drawn as follows, with the control bit on top and the target bit on the bottom:


Note that $(C N O T)^{2}=1$, i.e., $C N O T^{-1}=C N O T$.

- SWAP


$$
\mathrm{SWAP}=\left(\begin{array}{cccc}
1 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 \\
0 & 1 & 0 & 0 \\
0 & 0 & 0 & 1
\end{array}\right)
$$

4.3 n-qubit gates:

- local n-qubit gates formed as tensor products of one-qubit gates, e.g., $H^{\otimes n}$
- Toffoli gate

This is a 3-qubit generalization of the CNOT gate. The third, target, qubit is flipped iff both the first and second qubits are in state $1 . T O F F^{2}=1$.


The Toffoli gate can be decomposed into a combination of one-qubit and two-qubit gates. See Figures 3 and 4.

### 4.4 Useful gate equivalences

- SWAP equals $3 \times C N O T$

See Figure 5.
Suppose we have two qubits in state $\left|y_{2}, y_{1}\right\rangle$ :

$$
\begin{aligned}
& {\left[\begin{array}{l}
a \\
b
\end{array}\right] \otimes\left[\begin{array}{l}
c \\
d
\end{array}\right]} \\
& =a c|00\rangle+a d|01\rangle+b c|10\rangle+b d|11\rangle
\end{aligned}
$$

Apply the first CNOT:

$$
a c|00\rangle+a d|01\rangle+b d|10\rangle+b c|11\rangle
$$

Apply the second CNOT:

$$
a c|00\rangle+b c|01\rangle+b d|10\rangle+a d|11\rangle
$$

Apply the third CNot:

$$
\begin{aligned}
& a c|00\rangle+b c|01\rangle+a d|10\rangle+b d|11\rangle \\
& =c a|00\rangle+c b|01\rangle+d a|10\rangle+d b|11\rangle \\
& =\left[\begin{array}{l}
c \\
d
\end{array}\right] \otimes\left[\begin{array}{l}
a \\
b
\end{array}\right]
\end{aligned}
$$

The resulting state is $\left|y_{1}, y_{2}\right\rangle$, i.e., the states of the two qubits have been swapped.

- Control and target of CNOT can be swapped by conjugating both qubits with $H$

See Figure 6.
Proof: see homework 2.

## 5 Universality of Gate Sets

### 5.1 Classical

The NAND gate is universal for classical computation. The NAND gate is the result of applying NOT to $a A N D b=a \wedge b=a \uparrow b$. See Figure 7.
For any boolean function $\{0,1\}^{n} \longrightarrow\{0,1\}$, there is a circuit built of NAND gates (possibly with FANOUT= copy) for that function. Note that neither of these gates are reversible.
In general, the circuit may require an exponential number $2^{n}$ of gates. Functions which can be efficiently evaluated require only a polynomial number $n^{c}$ gates. Complexity theory categorizes the scaling of the resources, esp. the number of gates, with the number of bits $n$. Provided the gate set is universal, the distinction between functions which require exponentially large circuits and those which can be computed with polynomial-size circuits does not depend on the chosen set of gates.

### 5.2 Quantum

A set $G$ of quantum gates is called universal if for any $\varepsilon>0$ and any unitary matrix $U$ on $n$ qubits, there is a sequence of gates $g_{1}, \ldots, g_{l}$ from $G$ such that $\left\|U-U_{g_{l}} \cdots U_{g_{2}} U_{g_{1}}\right\| \leq \varepsilon$.

Here $U_{g}$ is $V \otimes I$, where $V$ is the unitary transformation on $k$ qubits operated on by the quantum gate $g$, and $I$ is the identity acting on the remaining $n-k$ qubits. The operator norm is defined by $\left\|U-U^{\prime}\right\|=$ $\max _{|v\rangle \text { unit vector }} \|\left(U-U^{\prime}\right)|v\rangle \|$. . Recall that for a vector $w,\|w\|=\sqrt{\langle w \mid w\rangle}$. .)
Examples of universal gate sets include

- CNOT and all single qubit gates
- CNOT, Hadamard, and suitable phase flips
- CNOT, Hadamard, $X$ and $T(\pi / 8)$
- Toffoli and Hadamard


Figure 1: An $X$ gate conjugated by $H$ gates is a $Z$ gate.


Figure 2: A $Z$ gate conjugated by $H$ gates is an $X$ gate.

| Inputs |  |  | Outputs |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $a$ | $b$ | $c$ | $a^{\prime}$ | $b^{\prime}$ | $c^{\prime}$ |
| 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 0 | 0 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 | 1 |
| 1 | 0 | 0 | 1 | 0 | 0 |
| 1 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 1 | 1 | 1 |
| 1 | 1 | 1 | 1 | 1 | 0 |



Figure 3: Toffoli gate, a 3-qubit double controlled NOT gate (bit c is flipped iff both a and bare 1.


Figure 4: A Toffoli gate can be decomposed into a circuit of 1- and 2-qubit gates. Here $V=\left[\begin{array}{cc}1 & 0 \\ 0 & i\end{array}\right]=$ $R_{z}(\pi / 2)$.


Figure 5: A $S W A P$ gate is three back to back $C N O T$ gates with control and target qubits alternating.


Figure 6: Control and target qubits of CNOT can be exchanged by conjugating with $H$ on both qubits.

| $a$ | $b$ | $a \uparrow b$ |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 0 | 1 | 1 |
| 1 | 0 | 1 |
| 1 | 1 | 0 |



Figure 7: Classical NAND gate and its truth table.

