

Principle of Least Privilege

Principle of least privilege

- Privilege
 - Ability to access or modify a resource
- Principle of least privilege
 - A system module should only have the minimal privileges needed for intended purposes
- Privilege separation
 - Separate the system into independent modules
 - Each module follows the principle of least privilege
 - Limit interaction between modules

Unix access control

- File has access control list (ACL)
 - Grants permission to user ids
 - Owner, group, other
- Process has user id
 - Inherit from creating process
 - Process can change id
 - Restricted set of options
 - Special “root” id

	File 1	File 2	...
User 1	read	write	-
User 2	write	write	-
User 3	-	-	read
...			
User m	Read	write	write

Unix file access control list

- Each file has owner and group
- Permissions set by owner
 - Read, write, execute
 - Owner, group, other, setuid/setgid
 - Represented by vector of four octal values
- Only owner or root can change permissions
 - This privilege cannot be delegated or shared

- $\underbrace{rwx}_{\text{ownr}} \underbrace{rwx}_{\text{grp}} \underbrace{rwx}_{\text{otr}}$

Privileged Programs

- Privilege management is coarse-grained in today's OS
 - Root can do anything
- Many programs run as root
 - Even though they only need to perform a small number of privileged operations
- What's the problem?
 - Privileged programs are juicy targets for attackers
 - By finding a bug in parts of the program that do not need privilege, attacker can gain root

What Can We Do?

- Drop privilege as soon as possible
- Ex: a network daemon only needs privilege to bind to low port # (<1024) at the beginning
 - Solution?
 - Drop privilege right after binding the port
- What benefit do we gain?
 - Even if attacker finds a bug in later part of the code, can't gain privilege any more
- How to drop privilege?
 - Setuid/setgid programming in UNIX

Effective user id (EUID) in UNIX

- Each process has three Ids
 - Real user ID (RUID)
 - same as the user ID of parent (unless changed)
 - used to determine which user started the process
 - Effective user ID (EUID)
 - from set user ID bit on the file being executed, or sys call
 - determines the permissions for process
 - file access and port binding
 - Saved user ID (SUID)
 - So previous EUID can be restored
- Real group ID, effective group ID, used similarly

Operations on UIDs

- Root
 - ID=0 for superuser root; can access any file
- Fork and Exec
 - Inherit three IDs, except exec of file with setuid bit
- Setuid system calls
 - seteuid(newid) can set EUID to
 - Real ID or saved ID, regardless of current EUID
 - Any ID, if EUID=0
- Details are actually more complicated
 - Several different calls: setuid, seteuid, setreuid

Setuid/setgid/sticky bits on executable Unix file

- Setuid/setgid/sticky bits
 - Setuid – set EUID of process to ID of file owner
 - Setgid – set EGID of process to GID of file
 - Sticky
 - Off: if user has write permission on directory, can rename or remove files, even if not owner
 - On: only file owner, directory owner, and root can rename or remove file in the directory

Setting UIDs

- **setresuid()** sets the real user ID, the effective user ID, and the saved set-user-ID of the calling process.
- **seteuid()** sets the effective user ID of the calling process.
- **setuid()** sets the effective user ID of the calling process. If the effective UID of the caller is root, the real UID and saved set-user-ID are also set.

Setting UIDs - What's Allowed?

Users choose any new UID to pass in to `setuid()`, but the OS checks them against certain rules and will raise an error, for example, if a normal user tries to call

`setuid(0)`.

```
setresuid(newruuid, neweuid, newsuid)  
(euid == 0)  
||  
(newruuid in (ruid, euid, suid) &&  
neweuid in (ruid, euid, suid) &&  
newsuid in (ruid, euid, suid))
```

```
seteuid(neweuid)  
(euid == 0)  
||  
(neweuid in (ruid, euid, suid))
```

```
setuid(newuid)  
(euid == 0)  
||  
(newuid in (ruid, suid))
```

Cases

```
(euid == 0) ⇒ (ruid:=newuid,  
euid:=newuid, suid:=newuid)  
(anything else) ⇒ (euid:=newuid)
```

Note: all policies are for Linux, differs on FreeBSD

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 neweuid in (ruid, euid, suid) &&
 newsuid in (ruid, euid, suid))
```

```
seteuid(neweuid)
(euid == 0)
||
(neweuid in (ruid, euid, suid))
```

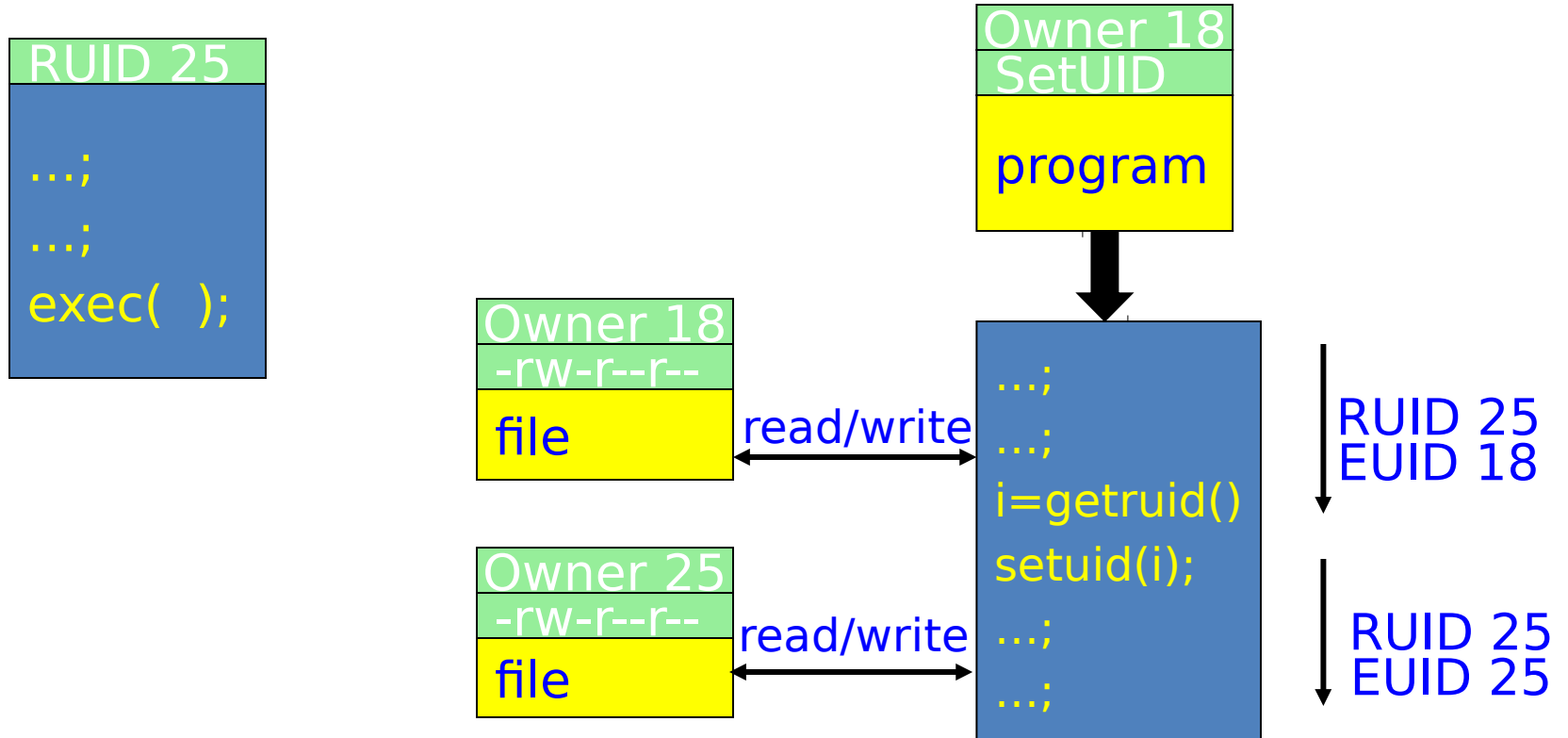
```
setuid(newuid)
(euid == 0)
||
(newuid in (ruid, suid))
```

Cases

```
(euid == 0) ⇒ (ruid:=newuid,
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Note: all policies are for Linux, differs on FreeBSD

Drop Privilege



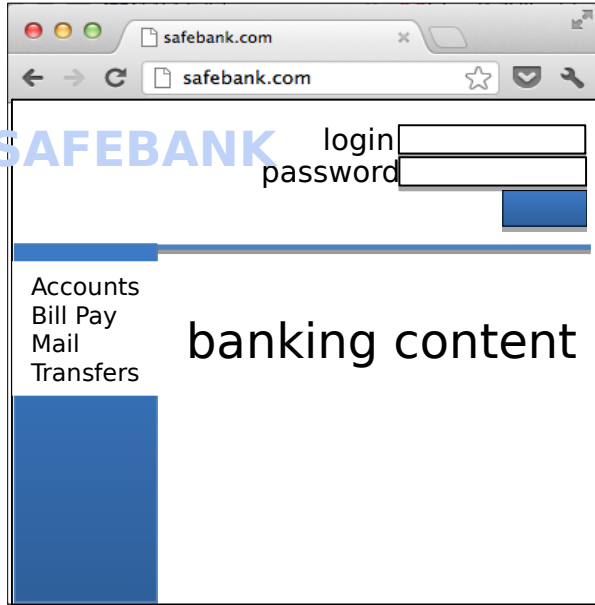
Web Security: Vulnerabilities & Attacks

Slide credit: Anthony Joseph and John Mitchell

Introduction

Web & http

(browser)



CLIENT

HTTP REQUEST:

```
GET /account.html HTTP/1.1  
Host: www.safebank.com
```



HTTP RESPONSE:

```
HTTP/1.0 200 OK  
<HTML> . . . </HTML>
```



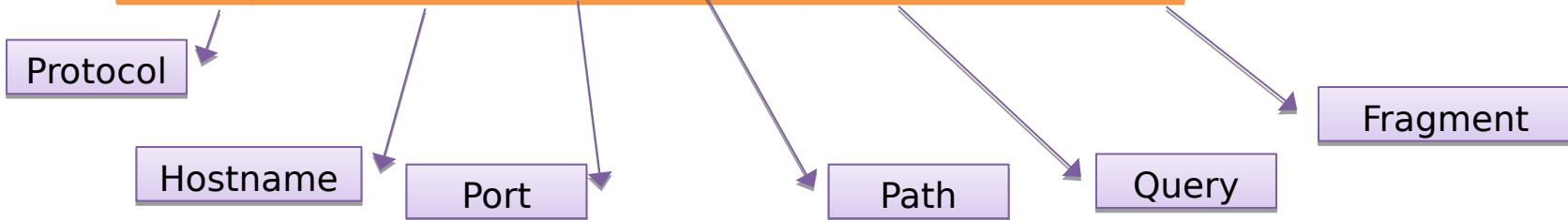
SERVER

URLs

- Global identifiers of network-retrievable documents

- **Example:**

`http://safebank.com:81/account?id=10#statement`



- Special characters are encoded as hex:
 - `%0A` = newline

HTTP Request

HTTP Response

Method File HTTP version Headers

```
GET /index.html HTTP/1.1
Accept: image/gif, image/x-bitmap,
image/jpeg, /*
Accept-Language: en
Connection: Keep-Alive
User-Agent: Chrome/21.0.1180.75 (Macintosh;
Intel Mac OS X 10_7_4)
Host: www.safebank.com
Referer: http://www.google.com?q=dingbats
```

Blank line
Data – none for GET

```
GET : no side effect
POST : possible
```

HTTP version Status code Reason phrase

```
HTTP/1.0 200 OK
Date: Sun, 12 Aug 2012 02:20:42 GMT
Server: Microsoft-Internet-Information-
Server/5.0
Connection: keep-alive
Content-Type: text/html
Last-Modified: Thu, 9 Aug 2012 17:39:05 GMT
Set-Cookie: ...
Content-Length: 2543

<HTML> This is web content formatted using
html </HTML>
```

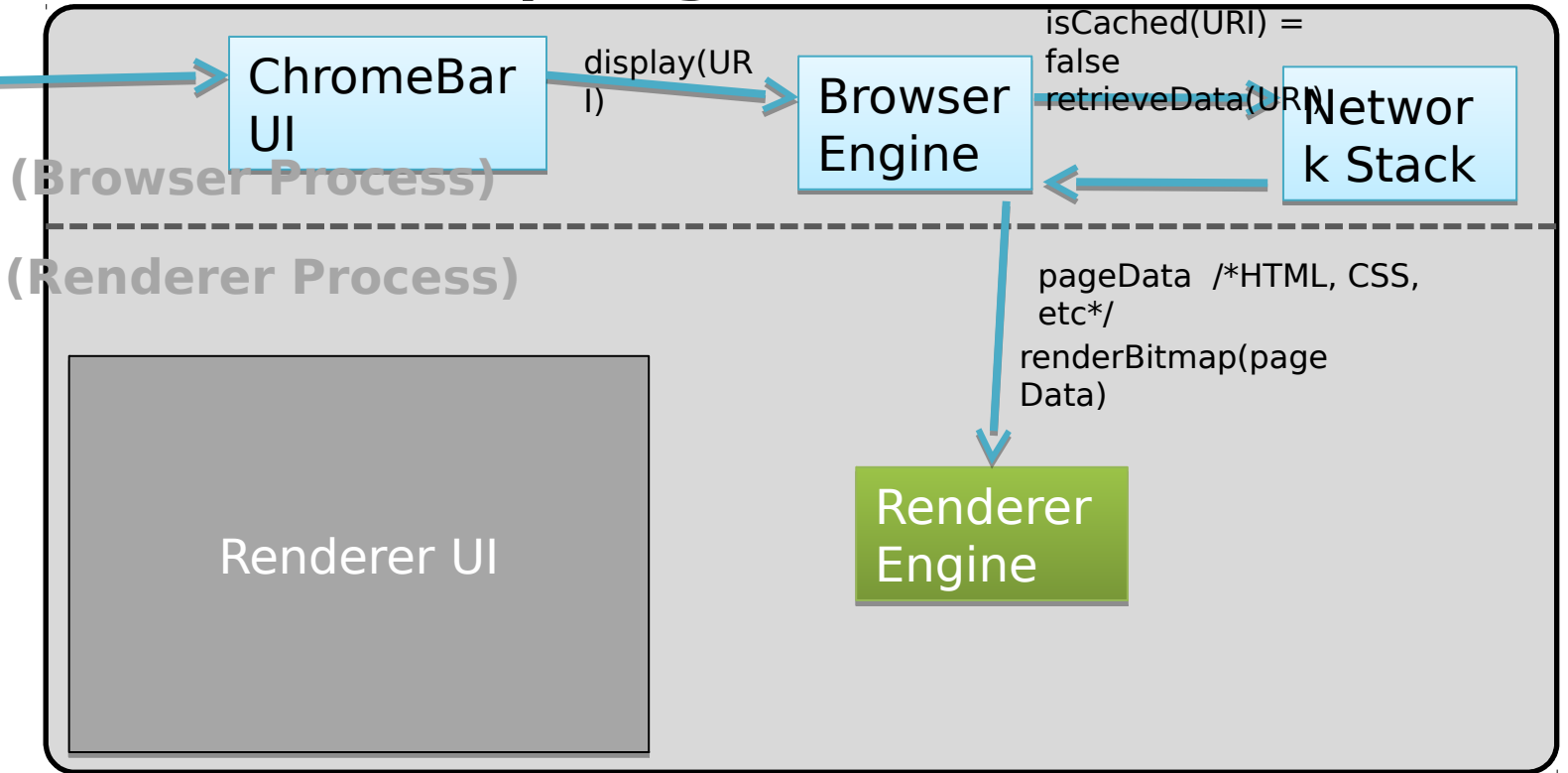
Cookies

How browser renders a page

Suppose you are visiting <http://safebank.com> in a modern web browser.



User enters <http://safebank.com> and presses go.



Rendering and events

- Basic execution model
 - Each browser window or frame
 - Loads content
 - Renders
 - Processes HTML and scripts to display page
 - May involve images, subframes, etc.
 - Responds to events
- Events can be
 - User actions: `OnClick`, `OnMouseover`
 - Rendering: `OnLoad`, `OnBeforeUnload`
 - Timing: `setTimeout()`, `clearTimeout()`

Document Object Model (DOM)

HTML

```
<html>
  <body>
    <div>
      foo
      <a>foo2</a>
    </div>
    <form>
      <input
type="text" />
      <input
type="radio" />
      <input
type="checkbox" />
    </form>
  </body>
</html>
```

Object-oriented interface used to read and write rendered pages

- web page in HTML is structured data
- DOM provides representation of this hierarchy

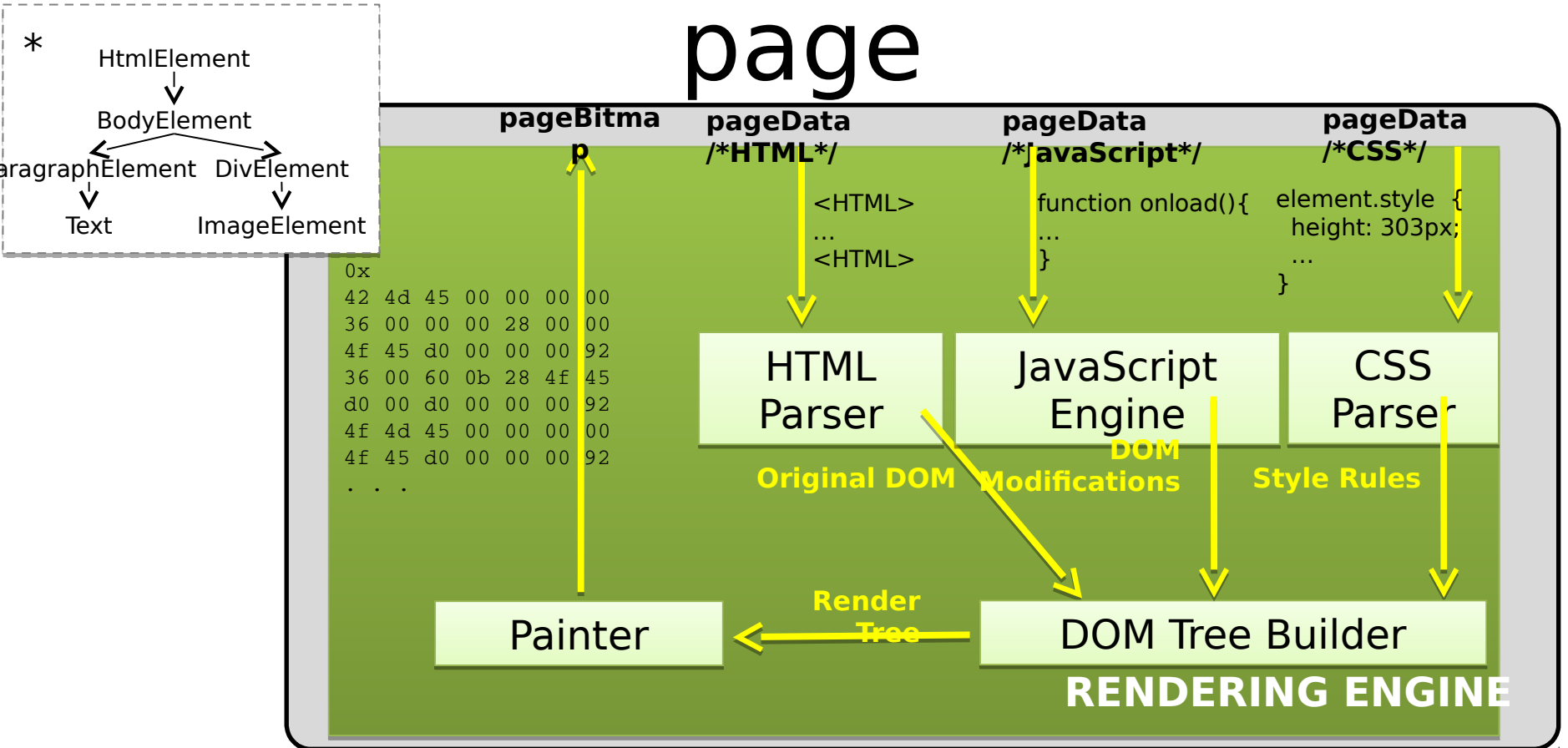
Examples

- **Properties:** document.alinkColor, document.URL, document.forms[], document.links[], document.anchors[]
- **Methods:** document.write(document.referrer)

DOM Tree

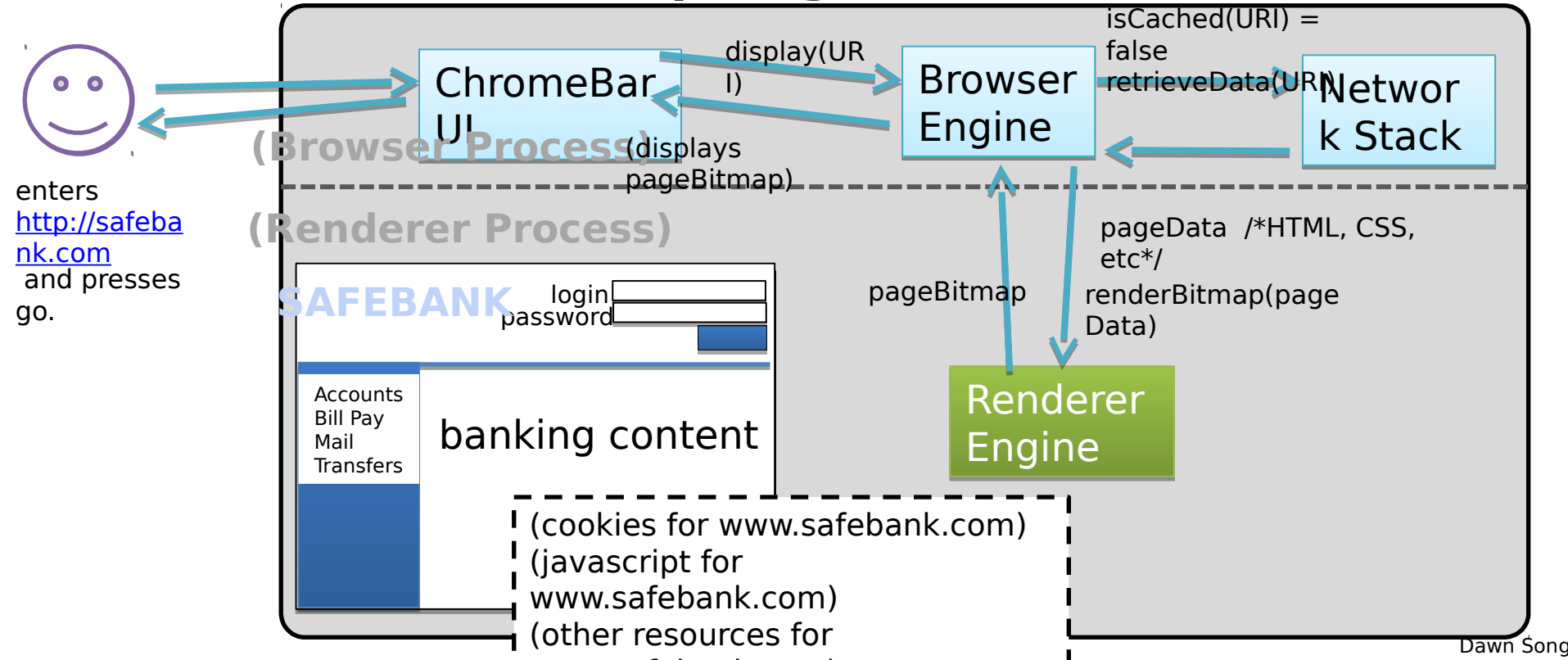
```
|-> Document
  |-> Element
    (<html>)
      |-> Element
        (<body>)
          |-> Element
            (<div>)
              |-> text node
              |-> Anchor
                |-> text node
                |-> Form
                  |-> Text-box
                  |-> Radio
                    Button
                      |-> Check Box
                      |-> Button
```

How browser renders a page



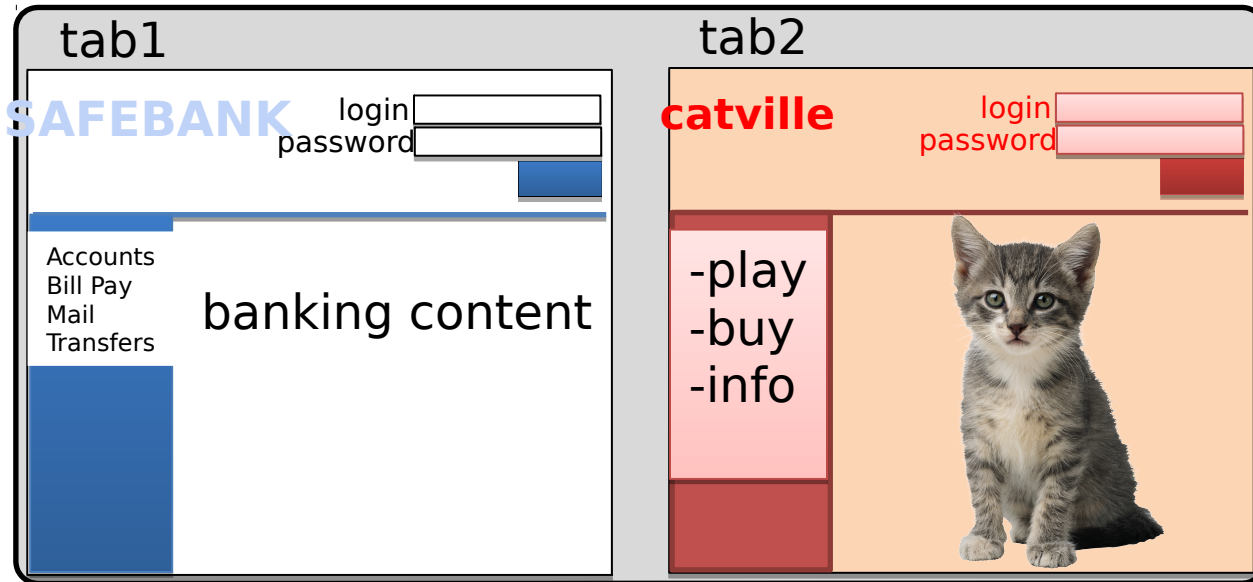
How browser renders a page

Suppose you are visiting <http://safebank.com> in a modern web browser.



Web Security Goals & Threat Model

Web Browser Security Goals



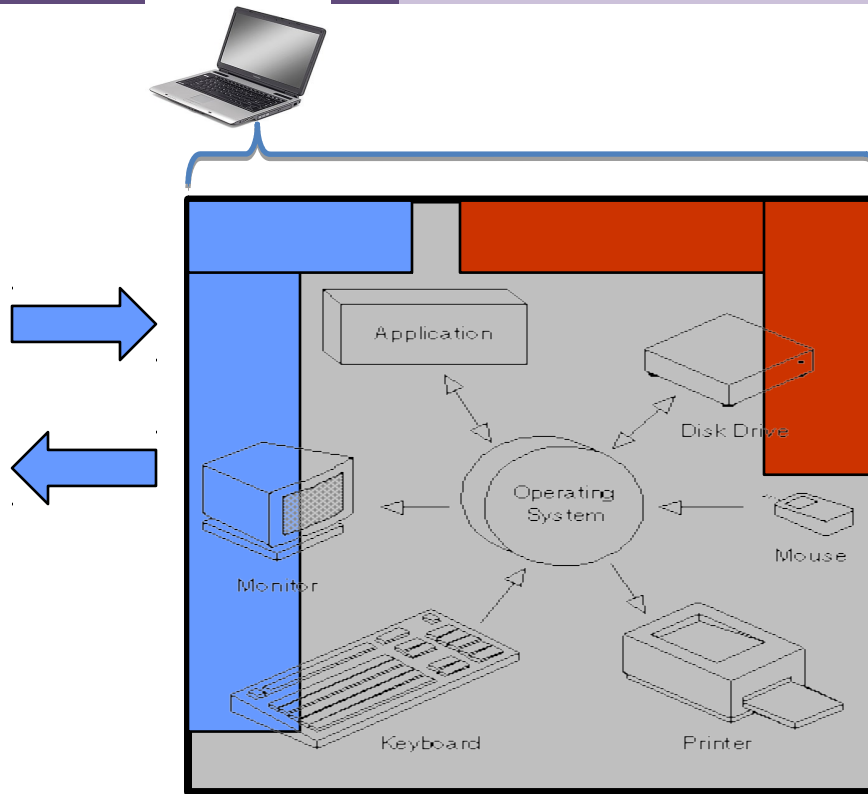
Security Goals

- tab 2 cannot compromise the user's computer or data
- tab2 cannot steal information from tab1 (without user permission)
- tab 2 cannot compromise the session in tab 1

cookies for www.safebank.com) (cookies for www.catville.com)
javascript for www.safebank.com)(javascript for www.catville.com)
other resources for www.safebank.com) resources for www.catville.com)



User



System

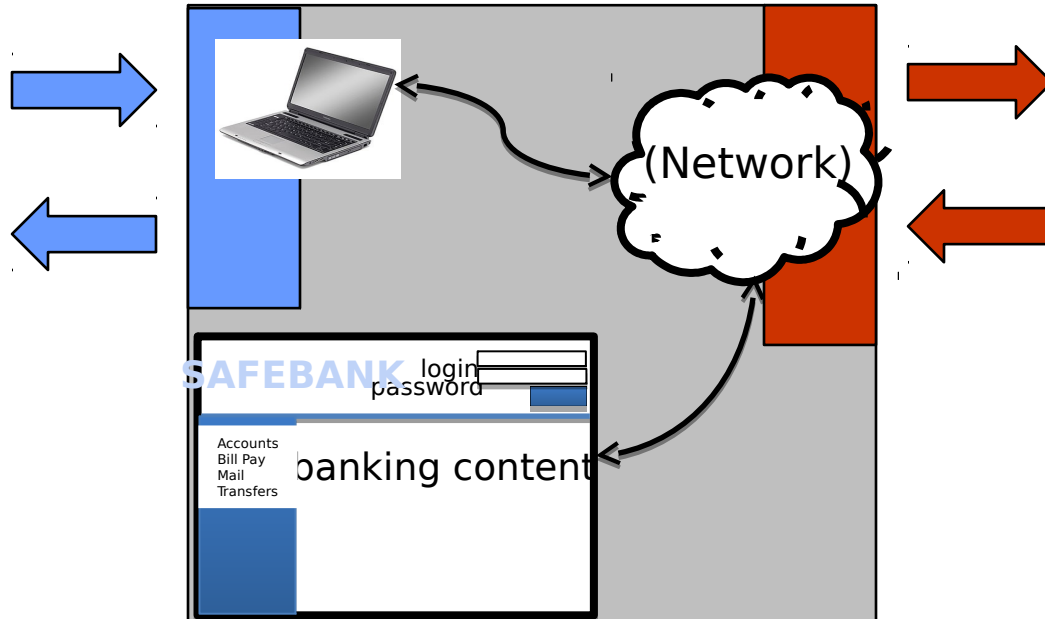


**OS/Malware
Attacker**

May control
malicious files and
applications on host



User

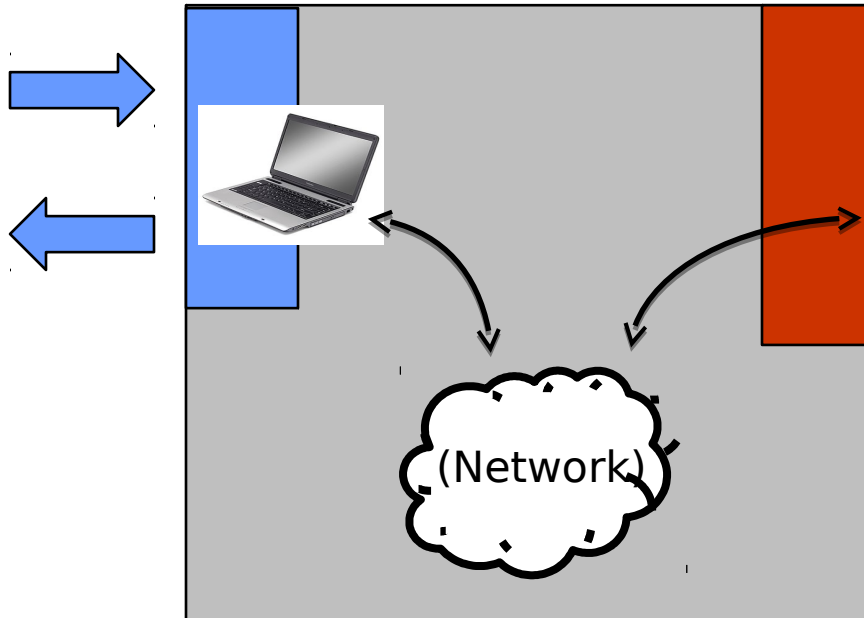


**Network
Attacker**

Intercepts and
controls network
communication



User



Web Attacker

Sets up malicious
site visited by
victim; no control
of network

Web Threat Models

Web attacker

- Control malicious site, which we may call “attacker.com”
- Can obtain SSL/TLS certificate for attacker.com
- User visits attacker.com
 - Or: runs attacker’s Facebook app, site with attack ad, .



Network attacker

- Passive: Wireless eavesdropper
- Active: Evil router, DNS poisoning



OS/Malware attacker

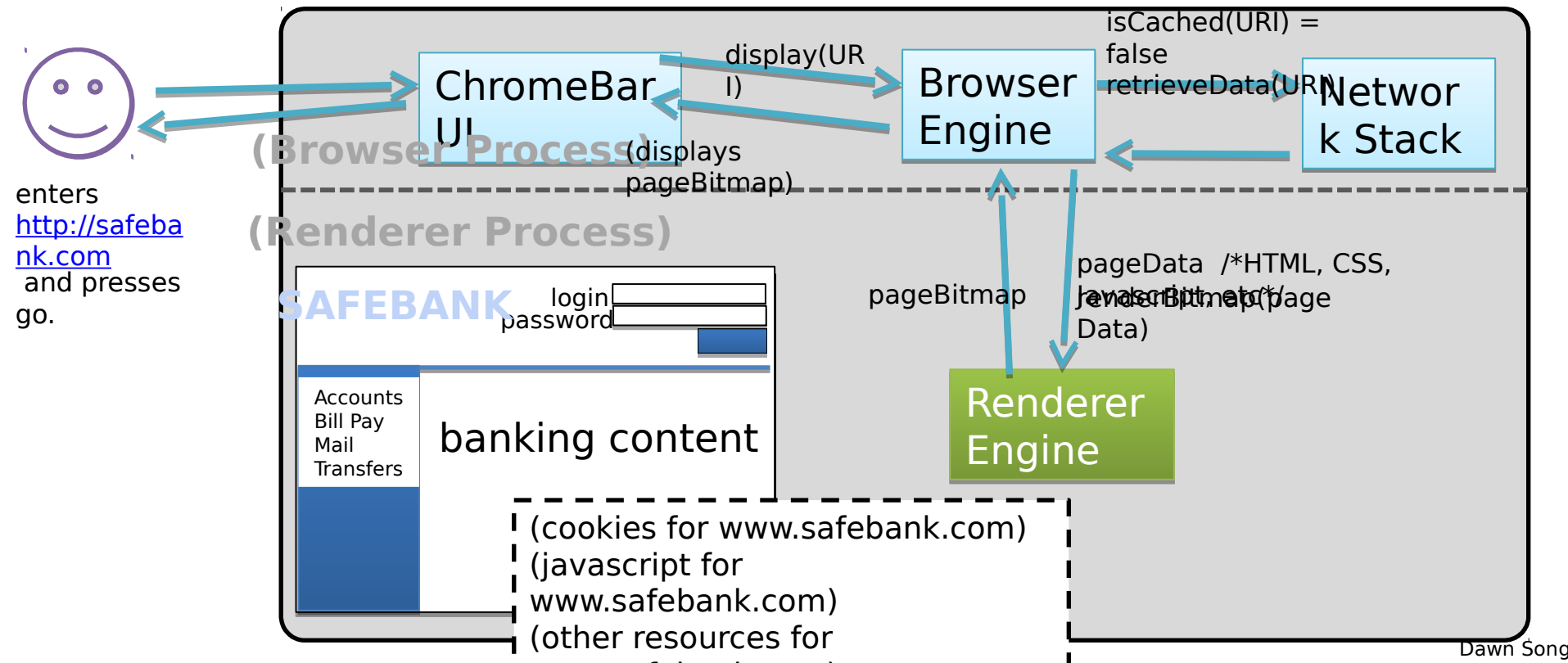
- Attackers may compromise host and install malware on host



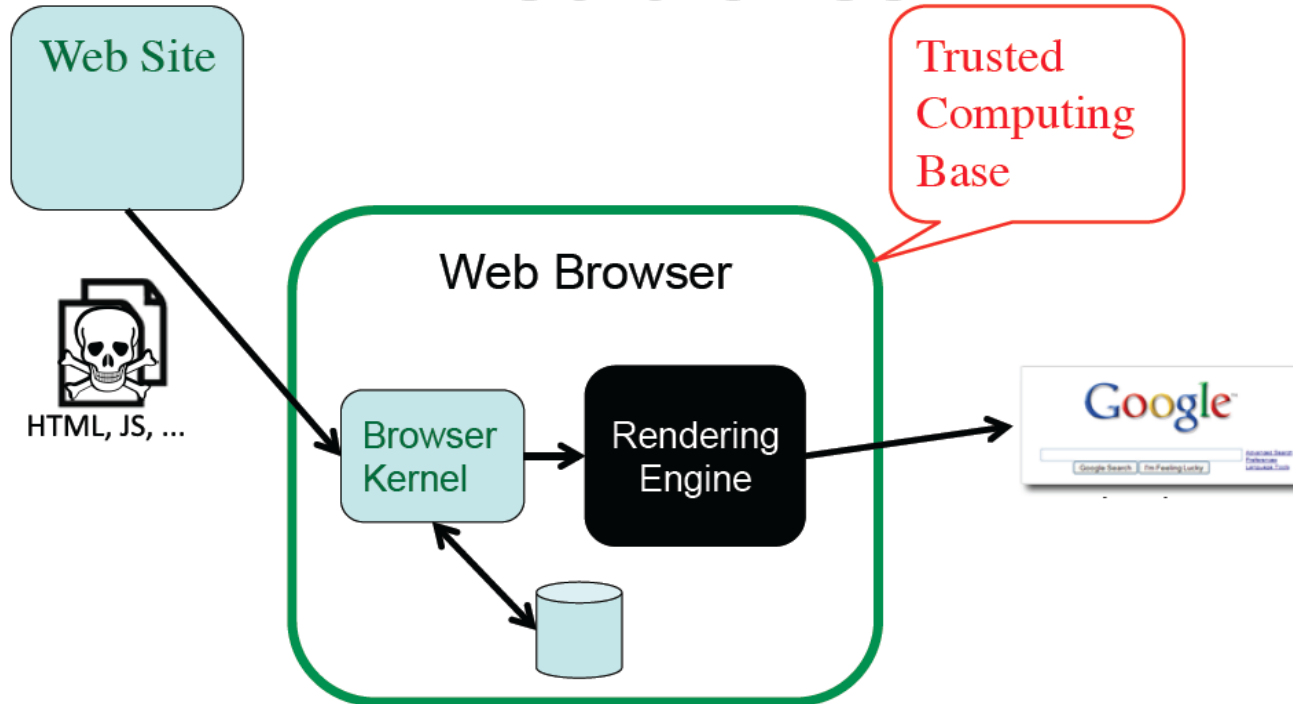
Isolation

How Browser Renders a Page

Suppose you are visiting <http://safebank.com> in a modern web browser.

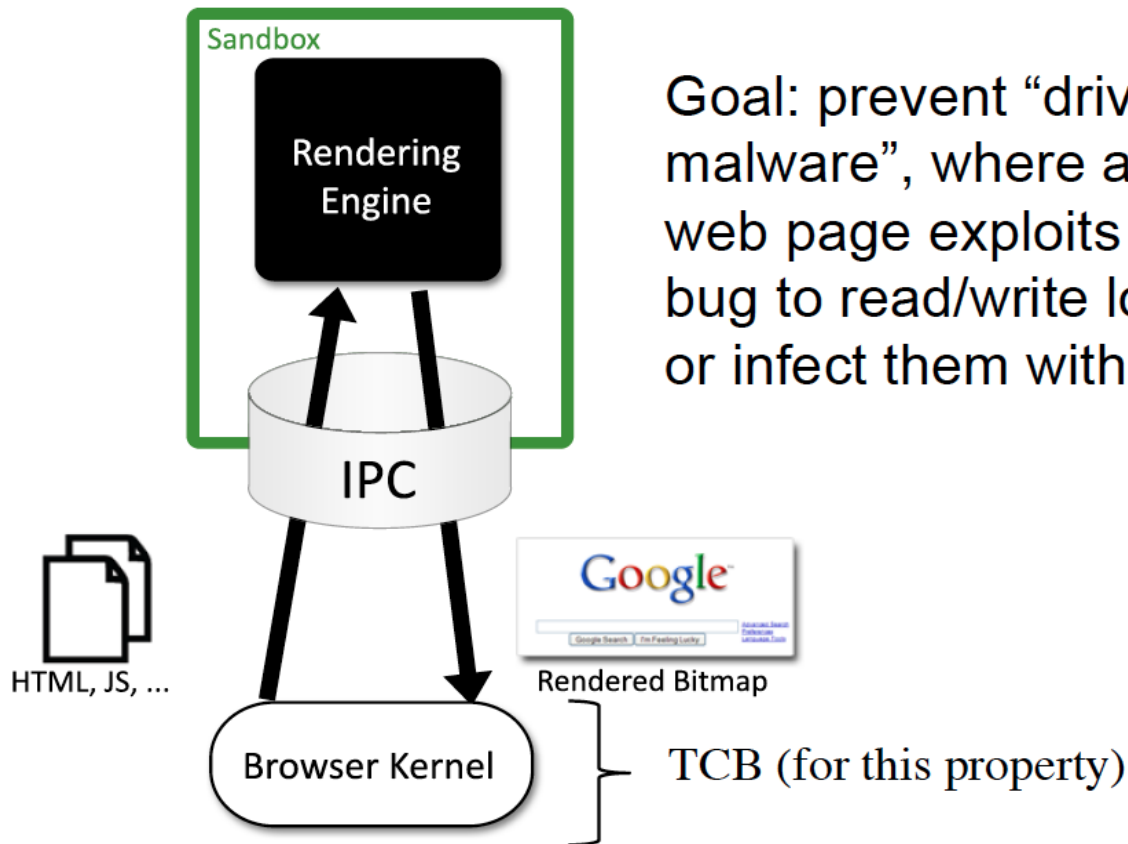


Web browser



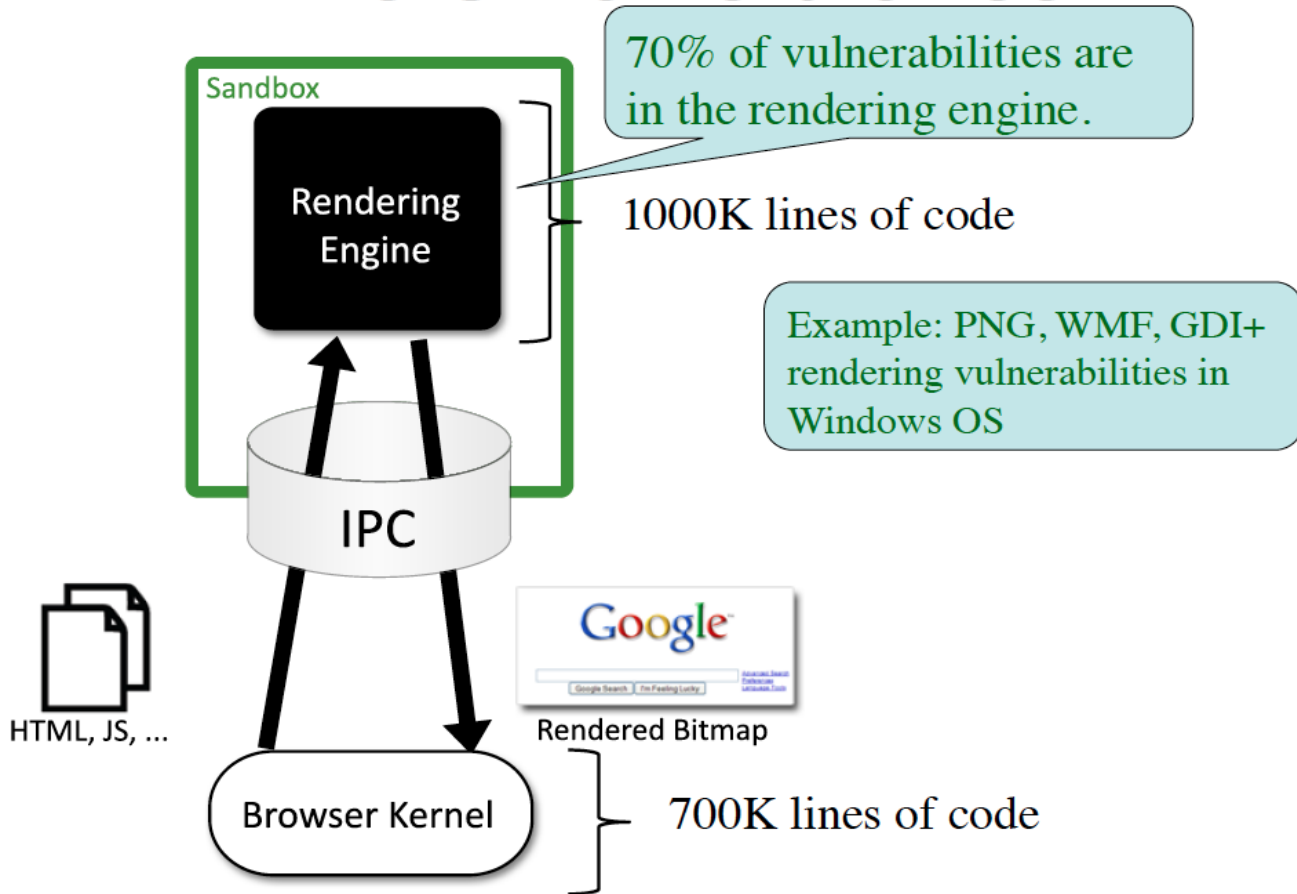
“Drive-by malware”: malicious web page exploits a browser bug to read/write local files or infect them with a virus

The Chrome browser



Goal: prevent “drive-by malware”, where a malicious web page exploits a browser bug to read/write local files or infect them with a virus

The Chrome browser



Same Origin Policy

Operating system



Web Browser



Primitives:

- Processes
- System calls
- File system

- Frames/iFrames
- Content (including JavaScript, ...)
- Document object model, cookies, localStorage

Principals:

Users

- Discretionary access control

“Origins”

- Mandatory access control

Vulnerabilities:

- Buffer overflow
- Root exploit

- Cross-site scripting
- Cross-site request forgery
- Cache history attacks

Origin of Browser Primitives

Cookies

Setting Cookies:

Default origin is **domain** and **path** of setting URL

Javascript

Imported **in** a page or frame:

Has the **same origin** as **that** page or frame

Embedded **in** a page or frame:

Has the **same origin** as **that** page or frame

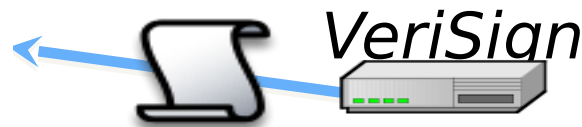
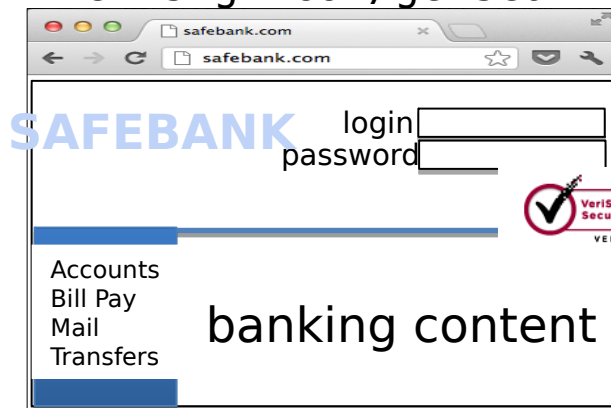
DOM

Each frame of a page:

Origin is
protocol://host:port

Library import

```
<script  
  src=https://seal.verisign.com/getseal?host_name=safebank.com>  
</script>
```



- Script has privileges of the importing page, NOT source server.



Same-origin policy (for Javascript and DOM)

Two documents have the same origin if:

Same protocol (https, http, ftp, etc)

Same domain (safebank.com, etc)

Same port (80, 23, 8080, etc)

Results of same-origin checks against
"http://cards.safebank.com/c1/info.htm"

Same origin:

"http://cards.safebank.com/c2/edit.html"
"http://cards.safebank.com/"

Different origin:

"http://www.cards.safebank.com"
"http://catville.com"
"http://cards.safebank.com"
"http://cards.safebank:8080"