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## Problem Set \#2: Programmable Logic and HDLs Assigned 9/6/2007, Due 9/14/2007 at 2 PM

Problem 1. You are given a single 2 -input MUX. How many of the 2 -input boolean functions can you obtain by wiring the three inputs of this device to some combination of $\mathrm{A}, \mathrm{B}$, Logical 0 , and Logical 1 . You do not have the complements of A and B . (Show how.) How many more can you get with 2 MUXes? 3 Muxes?

| $\ln 0$ | In1 | Sel | Out |
| ---: | ---: | ---: | ---: |
| 0 | X | 0 | 0 |
| 1 | X | 0 | 1 |
| A | X | 0 | A |
| B | X | 0 | B |
| 1 | 0 | A | $\sim \mathrm{~A}$ |
| 1 | 0 | B | $\sim \mathrm{~B}$ |
| 0 | A | B | $\mathrm{~A} \& \mathrm{~B}$ |
| 1 | A | B | $\mathrm{~A} \mid \sim \mathrm{B}$ |
| A | 1 | B | $\mathrm{~A} \mid \mathrm{B}$ |
| A | 0 | B | $\mathrm{~A} \& \sim \mathrm{~B}$ |
| 1 | B | A | $\sim \mathrm{~A} \mid \mathrm{B}$ |
| B | 0 | A | $\sim \mathrm{~A} \mathrm{\&} \mathrm{B}$ |



With a second MUX you can add NOR, NAND, XOR, XNOR, so you get them all. With a third, you can build a 2-LUT.

Problem 2. Based on the information in lecture and the data sheet on the Xilinx Vertex family, how does the number of User I/O pins grow with the number of LUTs? How do the BlockRAM Bits scale with LUTs?

BlockRam Bits scale about linearly with LUTs
User I/O pins grows as the log of the number of LUTs
Problem 3: You are synthesizing logic (by hand) for and idealized FPGA consisting of many CLBs as shown in the following figure and interconnect that ties them together. Show how to use a collection of these blocks to implement the 4-bit Parallel-to-Serial converter discussed in lecture (slide 10 of Lecture 3). For each block, specify the contents of the 4-LUT and the latch. Show how the blocks are wired together.


Problem 4: In lab 0 you implemented an XOR using 4 nand gates. Write the structural verilog module corresponding to the schematic below.


Write the same behavior as a behavioral module using continuous assignment.

```
module xor ( out, a, b );
    input a, b;
    output out;
    wire nAB, Ao, Bo;
    nand_gate (nab, a, b);
    nand_gate (Ao, a, nAB);
    nand_gate (Bo, b, nAB);
    nand_gate (out, Ao, Bo);
endmodule
```

Write the same behavior as a behavioral module using Blocking assignment.

```
module xor ( out, a, b );
    input \(a, b\);
    output out;
    assign out \(=\mathbf{a}^{\wedge} \mathbf{b}\);
endmodule
```

Write the same behavior as a behavioral module using non-Blocking assignment.
module xor ( out, a, b );
input $a, b$;
output out;
reg out
always (a OR b)
out $<=\mathbf{a}^{\wedge} \mathbf{b}$;
endmodule

Problem 5. Some aspects of digital design are deeply rooted in the mathematical elegance of Boolean algebra. Others are motivated by clean geometrical structure that yields simple physical layouts. One example of the latter in the Johnson counter shown in schematic below. With $n$ bits it is possible to represent $2^{n}$ values. This counter only uses 2 n of those possible values, so it has a few more flip-flops, but almost no logic and no carry propagation, so it is fast and easy to route. The clr input resets all the flip-flops to zero.


Generate the truth table how this counter "increments". This has four input columns (clr, c 2 , c 1 , and c 0 ) and three output columns (the new c 2 , c 1 , and c 0 ).

| Inputs | NextState |  |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| c1r | C2 | C1 | C0 | C2 | C1 | C0 |
| 1 | X | X | X | 0 | 0 | 0 |
| 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 0 | 1 | 1 | 0 |
| 0 | 1 | 1 | 0 | 1 | 1 | 1 |
| 0 | 1 | 1 | 1 | 0 | 1 | 1 |
| 0 | 0 | 1 | 1 | 0 | 0 | 1 |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 | 1 |
| 0 | 1 | 0 | 1 | 0 | 1 | 0 |

Show how this logical would be implemented on a 4-input, 3-output PLA.


Write a behavioral Verilog module for this counter using a bit vector for the counter value.

```
module JohnsonCounter( clk, clr);
    input clk, clr;
    wire [2:0] nextstate, current;
    always@(posedge clk) begin
        if(clr) current <= 3'b000;
    else current <= nextstate;
end
    always @(posedge clk) begin
    nextstate[0] <= current[1];
    nextstate[1] <= current[2];
    nextstate[2]<= not current[0];
    end
endmodule
```

Problem 6: Draw the a schematic that would implement the behavior described be each of the follow verilog modules.
module sifter 1 (in, A,B,C, D, clk);
input in, clk;
output $\mathrm{A}, \mathrm{B}, \mathrm{C}$;
$\operatorname{reg} \mathrm{A}, \mathrm{B}, \mathrm{C}$;
always @ (posedge clk) begin
A <= in;
B $<=\mathbf{A}$;
D $<=$ C
C $<=\mathbf{B}$;
end
endmodule

module sifter2 (in, A,B,C, D, clk);
input in, clk;
output A,B,C;
$\operatorname{reg} \mathrm{A}, \mathrm{B}, \mathrm{C}$;
always@ (posedge clk) begin
A = in;
$\mathbf{B}=\mathbf{A}$;

module sifter2 (in, A,B,C, D, clk);
input in, clk;
output A,B,C;
$\operatorname{reg} \mathrm{A}, \mathrm{B}, \mathrm{C}$;
always@ (posedge clk OR reset) begin

$$
\begin{aligned}
& \mathbf{A}=\mathbf{i n} ; \\
& \mathbf{B}=\mathbf{A} ;
\end{aligned}
$$



